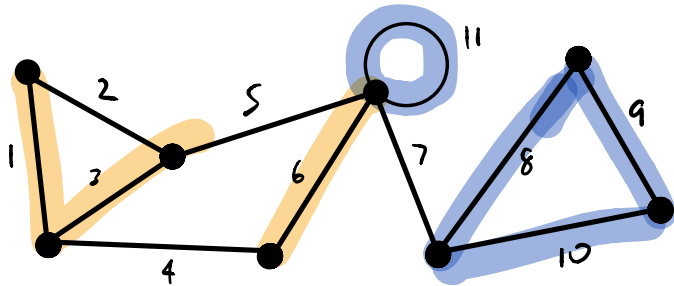


What IS a Matroid?

A matroid is a "groundset" E with a set of "independent" subsets $I \subseteq 2^E$ that satisfy some "nice" criteria.

Let's look at Examples first.

Definition: On a graph, call a set of edges independent if it contains no cycles.



What are some independent sets in the graph above?

$\{1\}$, $\{6\}$, $\{8\}$, $\{1, 3, 6\}$

$\{1, 2, 5, 6, 7, 10, 9\}$

Dependent Set $\{8, 9, 10\}$

$\{11\}$

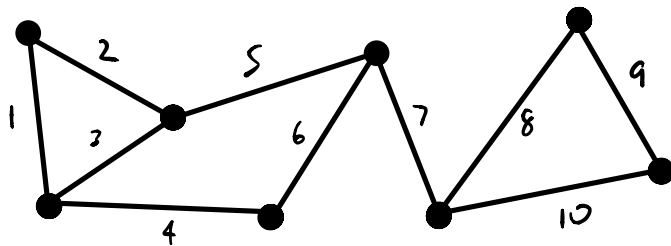
$\{5, 12\}$

$$G = (V, E) \rightarrow M = (E, \mathcal{I})$$

What can we say about $\mathcal{I} = \{A \subseteq E \mid A \text{ is cycle-free}\}$?

- 1) \emptyset is always cycle-free
- 2) Any subset of a cycle-free set is cycle-free.

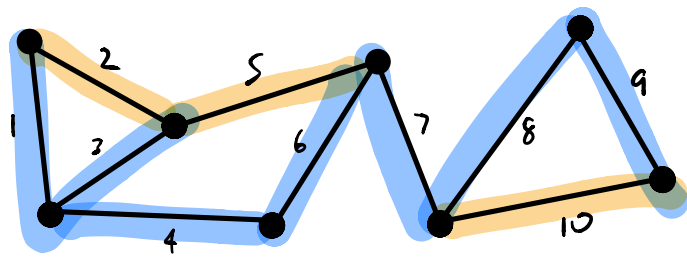
This is good and all, but cycles have a little more structure that can be captured.



What other properties does \mathcal{I} have?

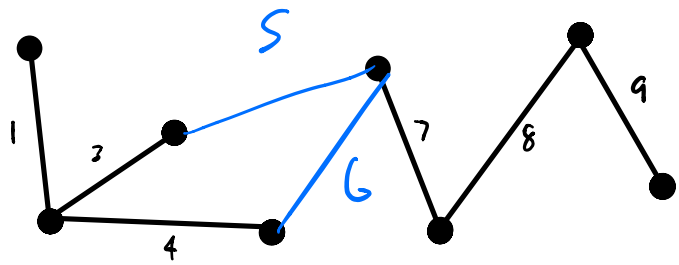
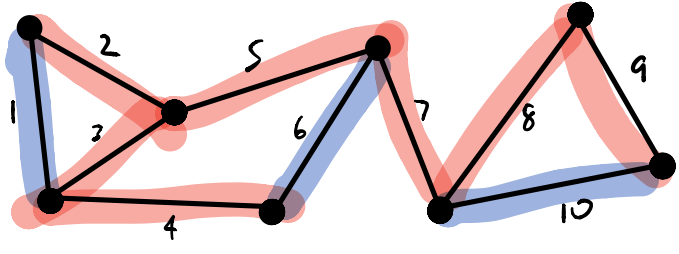
- 1) Closed under intersection
- 2) Spanning trees are maximal,
all independent subsets are subsets of trees
- 3) $\mathcal{I}/e = \{A \subseteq E \setminus e \mid A \cup \{e\} \in \mathcal{I}\}$
- 4) 7 can be added to any independent set

T_1

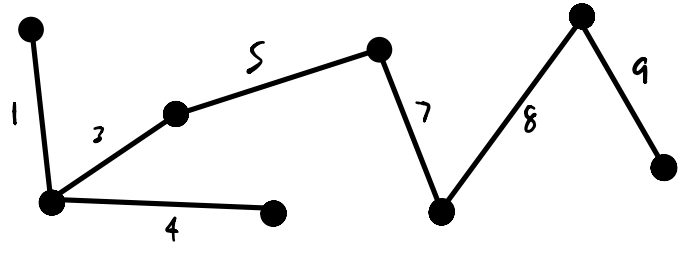


$6 \in T_1$
 $6 \notin T_2$

T_2



$5 \in T_2$
 $5 \notin T_1$



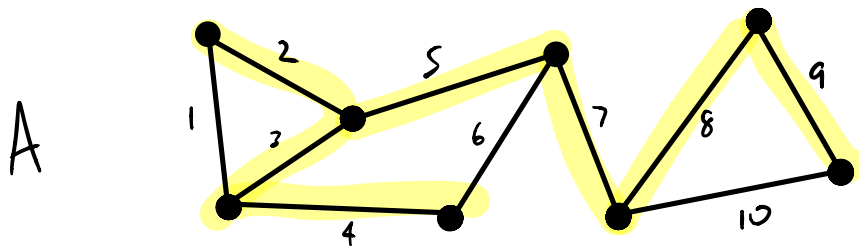
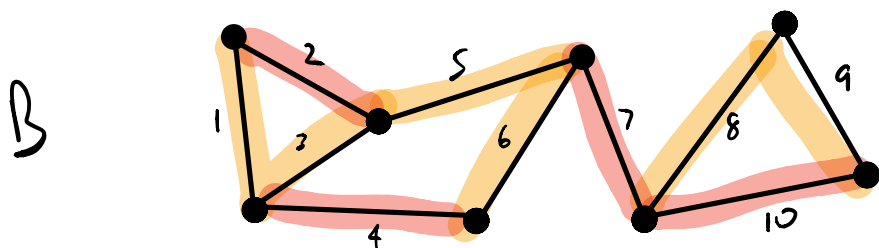
Definition of Matroid (Basis Edition)

A matroid $M = (E, \mathcal{B})$ is a set of "edges" E with a set \mathcal{B} of "bases" $T \subseteq E$ such that...

1) \mathcal{B} is nonempty

2) If $T_1, T_2 \in \mathcal{B}$, then for any $a \in T_1 \setminus T_2$, there exists $b \in T_2 \setminus T_1$ such that $(T_1 \setminus a) \cup b \in \mathcal{B}$

If you have two trees T_1 and T_2 , you can replace any differing edge in T_1 with an edge in T_2 .



$|A| > |B| \geq \#$ of edges that create cycles

Definition of Matroid (Independence Edition)

A matroid $M = (E, \mathcal{I})$ is a set of edges E with a set of Independent subsets \mathcal{I} such that...

1) $\emptyset \in \mathcal{I}$

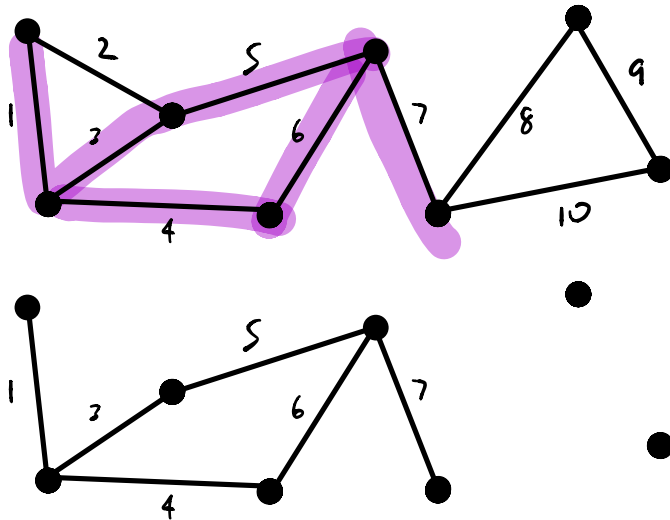
2) Subsets of independent sets are independent

3) If A and B are independent and $|B| < |A|$, then there is some $x \in A \setminus B$ such that $B \cup \{x\} \in \mathcal{I}$.

Trees (aka bases) are maximal independent sets.

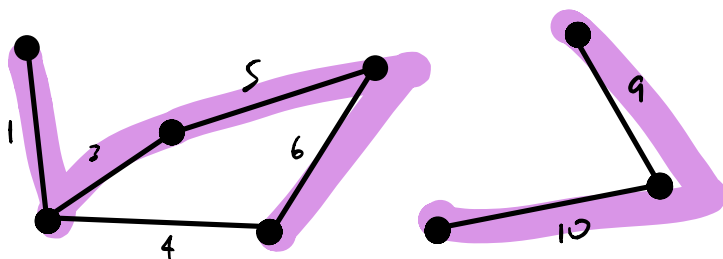
Cycles are minimal dependent sets.

If $S \subseteq E$, then we can consider the matroid restricted to S , $M|_S$.



Definition of rank of a subset $S \subseteq E$

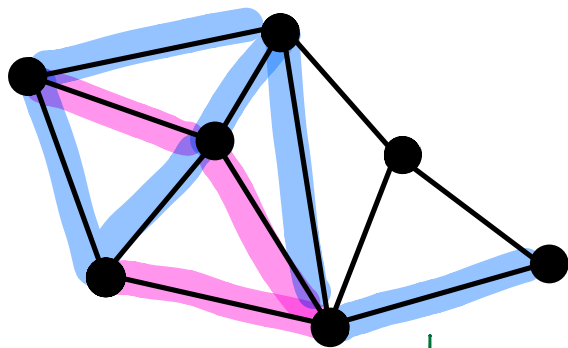
$r(S)$ is the size of the bases in $M|_S$,
 i.e. The number of edges in the maximal
 spanning forests of the restriction to S .



$$r(S) = 6$$

Definition of Flats of a Matroid.

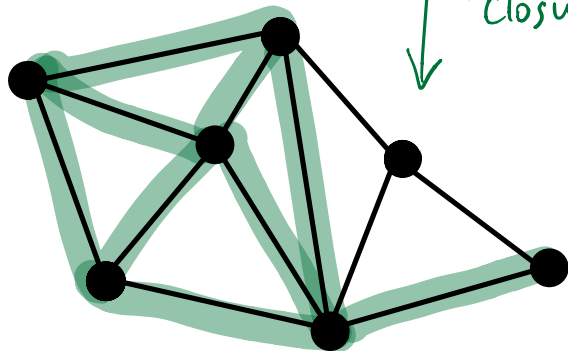
A set V is a flat of M if adding any edge to V would increase its rank.



Not a flat

(pink edges could be added without increasing the rank)

"closure"



Closure $cl(S)$ gives a flat with $r(cl(S)) = r(S)$

The Tutte Polynomial

Can be expressed on matroids.

For graphs G ,

$$T_G(x, 0) \propto \chi_G(1-x) \quad (\text{chromatic})$$

$$T_G(0, y) \propto F_G(1-y) \quad (\text{Flow})$$

For matroids M ,

$$T_M(x, y) = T_{M^*}(y, x).$$

$$T_M(x, y) = \sum_{A \subseteq M} T_{M|_A}(0, y) T_{M/A}(x, 0)$$

$$\sum \sim \chi_{M/A}(1-x) \sim \chi_w(1-y)$$

Situation

Goal A symmetric Tutte Polynomial

What we have

- 1) A good symmetric **chromatic polynomial**
(The Stanley Symmetric Chromatic Function)
- 2) A good way of dealing with **contractions**
and **homogeneous degrees** (Sesquigraphs)
- 3) A way of simplifying the Tutte Polynomial
into a **product of chromatic polynomials**

What's Next?

Jam the Stanley Chromatic function into the Convolutional formula!!

$$C_G(\bar{p}, \bar{q}) \equiv \\ \equiv \sum_{S \subseteq E} (-1)^{|A/S| + |B/(E-S)|} X_{A/S}(\bar{p}) X_{B/(E-S)}(\bar{q})$$

It turns out that

$$C_G(-\bar{p}, -\bar{q}) = D_G(\bar{p}, \bar{q})$$