Goeritz Matrix and Knot Coloring

G. Black, M Kikta, C. Li, L. Wiljanen, Y. Xuan

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Checkerboard Coloring of Link Diagram

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Definition 1

Let D be a link diagram in \mathbb{R}^2 . A checkerboard coloring of D is a coloring of $\mathbb{R}^2 \setminus D$ by $\{0, 1\}$ such that no arc has two monochromatic sides.

Theorem 2

Every classical link is checkerboard colorable.

Proof: Jordan curve theorem.



Checkerboard coloring of Virtual Links

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G. Black, M Kikta, C. Li, L. Wiljanen, Y. Xuan Note that not all Virtual links can be checkerboard colored



Checkerboard Colorable

Not Checkerboard Colorable

Goeritz Matrix

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Definition 3

The sign of a crossing $\eta(c)$:



Definition 4

Take a checkerboard colored link diagram D with the unbounded region shaded. Enumerate the shaded regions by 1, 2, ..., n. The *pre-Goeritz matrix* of D is a $n \times n$ matrix (g_{ij}) :

$$g_{ij} = \begin{cases} \sum_{\text{crossings between } i,j} \eta(c), \ i \neq j \\ -\sum_{k \neq i} g_{ik}, \ i = j \end{cases}$$

Goeritz Matrix

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Definition 5

Since the sum of rows of pre-Goeritz matrix is always 0 by definition, we can always delete one row and one column without losing any information. What we obtain is called *Goeritz matrix G*. Usually we delete the row and column corresponding to the unbounded region.

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Take the knot 8_{19} and assign each crossing $\pm 1{:}$



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Its pre-Goeritz matrix is

$$\vec{G} = \begin{bmatrix} 4 & -1 & 0 & -1 & -2 \\ -1 & -1 & 1 & 1 & 0 \\ 0 & 1 & -2 & 0 & 1 \\ -1 & 1 & 0 & -1 & 1 \\ -2 & 0 & 1 & 1 & 0 \end{bmatrix}$$

Delete the first row and column, we get the Goeritz matrix:

$$G = \begin{bmatrix} -1 & 1 & 1 & 0 \\ 1 & -2 & 0 & 1 \\ 1 & 0 & -1 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

Goeritz Matrix and Knot Coloring

G. Black, M. Kikta, C. Li, L. Wiljanen, Y. Xuan An arc in the link diagram starts and ends where it goes underneath.

Definition 6

A Fox coloring is a coloring of a link diagram by $\mathbb{Z}/p\mathbb{Z}$ (or $\mathbb{Z}/n\mathbb{Z}$) such that at each crossing the equation x + z = 2y holds:



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Goeritz Matrix and Knot Coloring

Definition 7

G. Black, M. Kikta, C. Li, L. Wiljanen, Y. Xuan Given an oriented link diagram *D*, a *Dehn coloring* is a coloring of $\mathbb{R}^2 \setminus D$ by $\mathbb{Z}/p\mathbb{Z}$. At each crossing it follows the equation A + B = C + D:



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Remark 1

The following properties hold for both Fox coloring and Dehn coloring:

- Any $k \in \mathbb{Z} / n\mathbb{Z}$ gives a monochromatic coloring.
- Adding up two colorings gives a coloring.
- Multiplying a coloring by k gives a coloring.

In conclusion:

Theorem 8

Every Fox/Dehn coloring of a link diagram is a $\mathbb{Z}/n\mathbb{Z}$ module (a vector space when n is prime).

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We can quotient out the monochromatic colorings, which can be done by setting one arc/the unbounded region to be 0. There is a one-to-one relation between the quotients of two colorings:

Each Dehn coloring gives a Fox coloring: assign the arc with the sum of the colors on both sides. The coloring is well defined and satisfies the rule x + z = 2y.



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Each Fox coloring gives a Dehn coloring: assign 0 to the unbounded region. For the rest region, draw a line connecting the center of the region with the interior of the unbounded region. Once the line enters from region R_1 to region R_2 crossing arc *c*, color R_2 by color of *c* minus the color of R_1 . We can prove the coloring is well defined and satisfying the rule A + B = C + D.

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Warning

The two maps are not inverses of each other.

In fact, the Fox coloring and Dehn coloring can be given via Goeritz matrix:

Theorem 9

The quotiented Dehn coloring, the quotiented Fox coloring, and the solution space of Goeritz matrix are isomorphic as $\mathbb{Z}/n\mathbb{Z}$ modules.

Alexander Numbering

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Remark 2

Checkerboard coloring is a $\mathbb{Z} \,/ 2 \, \mathbb{Z}$ Dehn coloring.

Checkerboard coloring can be generalized to virtual links via Alexander numbering(mod 2):

Definition 10

For an oriented virtual link diagram, an *Alexander numbering* is a coloring of its semi-arcs(arc between two classical crossings) given by the following relation:



Alexander Numbering

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Theorem 11

The following are equivalent:

- A virtual link L has a diagram D which is Alexander numberable(mod 2).
- (Σ, L) has a (oriented) spanning surface, where (Σ, L) is the embedding of L in the Carter surface corresponding to D.
- Σ\D' is checkerboard colorable, where D' is the link diagram on the Carter surface Σ corresponding to D.
- D has no "windmill" crossing.

Pseudo Goeritz Matrix

Goeritz Matrix and Knot Coloring

G. Black, M. Kikta, C. Li, L. Wiljanen, Y. Xuan N. Kamada has generalized Goeritz matrix to virtual link diagrams:

Definition 12

The first and second local index η :



fist local index



second local index

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Pseudo Goeritz Matrix

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Definition 13

For an oriented virtual diagram D with n semi-arcs(arcs between two classical crossings), the *pseudo Goeritz matrix* $G = (g_{ij})$ is an $n \times n$ matrix:

$$g_{ij} = \begin{cases} \sum_{\text{crossings between } i, j \eta, \ i \neq j \\ -\sum_{k \neq i} g_{ik}, \ i = j \end{cases}$$

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The knot 2-1 with second indices:



Then its pseudo Goeritz matrix is

$$\begin{bmatrix} -2 & 0 & 1 & 1 \\ 0 & 2 & -1 & -1 \\ 1 & -1 & 0 & 0 \\ 1 & -1 & 0 & 0 \end{bmatrix}$$

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Fox coloring matrix

Goeritz Matrix and Knot Coloring

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Definition 15

Definition 14

The Coloring matrix C of a Link diagram is obtained by deleting any row and any column of the pre-Coloring matrix \vec{C}

Remark 3

Any non-constant vector x in $\mathbb{Z}/n\mathbb{Z}$ such that $\vec{C}x = 0$ corresponds with a non-trivial fox n-coloring of the Link diagram

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The coloring equations are:

$$x_1 + x_3 = 2x_2 x_2 + x_4 = 2x_3 x_3 + x_5 = 2x_4 x_4 + x_1 = 2x_5 x_5 + x_2 = 2x_1$$

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Goeritz Matrix and Knot Coloring

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$$\overrightarrow{C} = \begin{bmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 1 & -2 & 1 & 0 \\ 0 & 0 & 1 & -2 & 1 \\ 1 & 0 & 0 & 1 & -2 \\ -2 & 1 & 0 & 0 & 1 \end{bmatrix}$$

By deleting the first row and column we get the coloring matrix

$$C = \begin{bmatrix} 1 & -2 & 1 & 0 \\ 0 & 1 & -2 & 1 \\ 0 & 0 & 1 & -2 \\ 1 & 0 & 0 & 1 \end{bmatrix}$$

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Goeritz Matrix and Knot Coloring

G. Black, M Kikta, C. Li L. Wiljanen, Y. Xuan To find an example of a coloring we can see the following equality holds in $\mathbb{Z}/5\mathbb{Z}$ and leads the the 5-coloring depicted below.

$$\begin{bmatrix} 1 & -2 & 1 & 0 & 0 \\ 0 & 1 & -2 & 1 & 0 \\ 0 & 0 & 1 & -2 & 1 \\ 1 & 0 & 0 & 1 & -2 \\ -2 & 1 & 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} 1 \\ 2 \\ 3 \\ 4 \\ 0 \end{bmatrix} = \begin{bmatrix} 0 \\ 0 \\ 0 \\ 0 \\ 0 \end{bmatrix}$$



Link determinants

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Definition 16

The Link determinant of a Link L is the determinant of the Coloring Matrix ${\it C}$

Theorem 17

The Link determinant of a Link is well defined and $|\det(C)|$ is a Knot Invariant

Theorem 18

for a prime number p a link has a nontrivial p-coloring iff p divides the link determinate

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Goeritz Matrix and Knot Coloring

G. Black, M Kikta, C. Li, L. Wiljanen, Y. Xuan By deleting the top row and left column we get the Coloring matrix:

$$C = \begin{bmatrix} 1 & 1 & 0 & -2 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 & 0 & -2 & 0 \\ 0 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & -2 & 0 & 1 & 1 & 0 \\ -2 & 0 & 0 & 0 & 0 & 1 & 1 \\ 0 & 0 & 0 & -2 & 0 & 0 & 1 \end{bmatrix}$$

Taking the determinant we get det(C) = 3. So we can conclude the only non-trivial prime coloring of this knot is a 3-coloring.

Knot determinant's and the Goeritz matrix

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Theorem 19

The absolute value of the determinant of the Goeritz Matrix of a link L is equal to the absolute value of Link determinant for a link L. $|\det(G)| = |\det(C)|$

With this theorem we can calculate the knot determinate faster because the Goeritz matrix is always smaller than the coloring matrix of a knot.

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Goeritz Matrix and Knot Coloring

G. Black, M Kikta, C. Li, L. Wiljanen, Y. Xuan We calculated the Goeritz matrix for 8_{19} on an earlier slide which was:

$$G = \begin{bmatrix} -1 & 1 & 1 & 0 \\ 1 & -2 & 0 & 1 \\ 1 & 0 & -1 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$$

So when we take the determinant of this matrix we get det(G) = -3. We can compare this to the determinant of the coloring matrix we calculated earlier and we can see that we have |det(G)| = 3 = |det(C)|.