# Recursion for the Partial-Dual Euler Genus Polynomial

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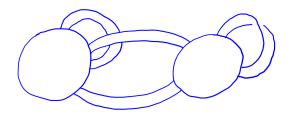
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# Ribbon Graphs

#### **Definition**

A ribbon graph G = (V, E) is a surface with boundary expressed as the union of vertex and edge discs such that

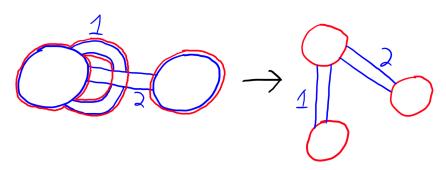
- Vertices and edges intersect at disjoint line segments
- Each such line segment lies on the boundary of exactly one vertex and edge
- Each edge contains exactly two line segments



# Partial Duality

### Definition (Chmutov, 2009)

For a ribbon graph G and  $A \subseteq E(G)$ , the partial dual  $G^A$  of G with respect to A is obtained by sewing a disc into each boundary component of the spanning ribbon subgraph induced by A and deleting the interiors of all vertices of G.



Partial duality at edge 1

### **Euler Genus**

#### Definition

The Euler genus  $\varepsilon(G)$  of a connected ribbon graph G is

$$\varepsilon(G) = \begin{cases} 2g(G) & \text{if G is orientable} \\ g(G) & \text{if G is non-orientable} \end{cases}$$

If G is not connected, then  $\varepsilon(G)$  is the sum of the Euler genus of its components.

In terms of the Euler characteristic,

$$\chi(G) = \nu(G) - e(G) + f(G) = 2k(G) - \varepsilon(G).$$

# Partial-Dual Polynomial

## Definition (Gross, Mansour, and Tucker, 2020)

The partial-dual Euler genus polynomial of a ribbon graph G is

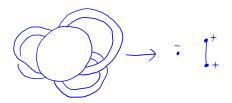
$${}^{\partial}\varepsilon_{G}(z)=\sum_{A\subset E(G)}z^{\varepsilon(G^{A})},$$

the generating function enumerating the edge subsets of G by the Euler genus of their partial duals.

# Signed Intersection Graph

#### **Definition**

Given a bouquet (one-vertex ribbon graph) B, its signed intersection graph G is the graph whose vertices are the edges of B, such that two vertices are adjacent if and only if the corresponding ribbon edges intersect. The sign of a vertex is positive (negative) if its corresponding ribbon edge is untwisted (twisted).



If v is a vertex of G then  $G^{\overline{v}}$  denotes the graph G with the opposite sign at v.

## Intersection Polynomial

## Theorem (Yan and Jin, 2022)

If two bouquets  $B_1$  and  $B_2$  have the same signed intersection graph, then  ${}^{\partial}\varepsilon_{B_1}(z)={}^{\partial}\varepsilon_{B_2}(z)$ .

Thus we can define the *intersection polynomial*  $\operatorname{IP}_G(z)$  of a signed intersection graph G as the partial-dual Euler genus polynomial of any bouquet whose signed intersection graph is G.

### Leaf Recursion

### Theorem (Yan and Jin, 2022)

Let G be a signed intersection graph and let  $v_1, v_2$  be adjacent vertices of G such that  $v_1$  is positive and has degree 1. Then

$$IP_G(z) = IP_{G-v_1}(z) + (2z^2) IP_{G-v_1-v_2}(z).$$

This allows us to recursively compute the intersection polynomial of any positive tree.

### **Cut-vertex Recursion**

### Theorem (L.)

Let G be a signed intersection graph such that  $G = G_1 \cup G_2$ ,  $G_1 \cap G_2 = \{v\}$ , and v is a cut-vertex of G. Let

$$\widetilde{Q_1} = \frac{2z^2 \, \mathsf{IP}_{G_2 - \nu} - (z+1) \, \mathsf{IP}_{G_2} + z \, \mathsf{IP}_{G_2^{\nabla}}}{2(2z^2 - z - 1)}$$

$$\widetilde{Q_2} = \frac{2z^2 \, \mathsf{IP}_{G_2 - \nu} - (z+1) \, \mathsf{IP}_{G_2^{\nabla}} + z \, \mathsf{IP}_{G_2}}{2(2z^2 - z - 1)}$$

$$\widetilde{Q_3} = \frac{z^2 (\mathsf{IP}_{G_2} + \mathsf{IP}_{G_2^{\nabla}}) - 2(z^3 + z^2) \, \mathsf{IP}_{G_2 - \nu}}{2z^2 - z - 1}$$

If v is positive, then  $\operatorname{IP}_G = \widetilde{Q_1} \operatorname{IP}_{G_1} + \widetilde{Q_2} \operatorname{IP}_{G_1^{\overline{v}}} + \widetilde{Q_3} \operatorname{IP}_{G_1 - v}$ . If v is negative, then  $\operatorname{IP}_G = \widetilde{Q_2} \operatorname{IP}_{G_1} + \widetilde{Q_1} \operatorname{IP}_{G_1^{\overline{v}}} + \widetilde{Q_3} \operatorname{IP}_{G_1 - v}$ .

# A Corollary

### Corollary (L.)

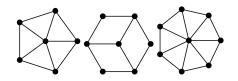
Let G be a signed intersection graph and v be a vertex of G. Then

$$\mathsf{IP}_{G-v}\left(-\frac{1}{2}\right) - \mathsf{IP}_{G}\left(-\frac{1}{2}\right) - \mathsf{IP}_{G^{\overline{v}}}\left(-\frac{1}{2}\right) = 0.$$

Moreover, if all vertices of G are positive, then

$$2z^{2} \operatorname{IP}_{G-v} + z \operatorname{IP}_{G} - (z+1) \operatorname{IP}_{G^{\overline{v}}} = 0.$$

# Non-intersection Graphs



## Theorem (Bouchet, 1994)

A simple graph is an intersection graph if and only if it does not have a vertex minor isomorphic to one of the above graphs.

Q: Can the intersection polynomial be extended to all graphs?

A: Yes, using the theory of delta-matroids.

### Delta-Matroids

#### Definition

A set system is a pair  $D=(E,\mathcal{F})$  of a set E and a collection  $\mathcal{F}$  of subsets of E. We call D a delta-matroid if  $\mathcal{F}$  is non-empty and satisfies the Symmetric Exchange Axiom:

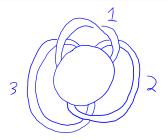
For all  $X, Y \in \mathcal{F}$  and  $u \in X \triangle Y$ , there exists  $v \in X \triangle Y$  such that  $X \triangle \{u, v\} \in \mathcal{F}$ .

E is called the ground set and elements of  $\mathcal F$  are called feasible sets.

# Ribbon-Graphic Delta-Matroids

#### Definition

From a ribbon graph G, we can obtain a delta-matroid  $D(G) = (E, \mathcal{F})$  by taking E to be the edges of G and  $\mathcal{F}$  to be the edge subsets corresponding to spanning ribbon subgraphs with exactly one boundary component. We call D(G) a ribbon-graphic delta-matroid.



If G is the ribbon graph above, then

$$D(G) = (\{1,2,3\}, \{\emptyset, \{1\}, \{1,2\}, \{1,3\}, \{2,3\}, \{1,2,3\}\}).$$

## **Twist**

### Definition

Let  $D = (E, \mathcal{F})$  be a delta-matroid and let  $A \subseteq E$ . Then the **twist** of D by A is the delta-matroid

$$D*A = (E, \{X \triangle A : X \in \mathcal{F}\}).$$

### Proposition

Let G be a ribbon graph and let  $A \subseteq E(G)$ . Then

$$D(G^A) = D(G) * A.$$

## Width

### Definition

Let  $D = (E, \mathcal{F})$  be a delta-matroid. The **width** of D is

$$w(D)=\max\{|F|:F\in\mathcal{F}\}-\min\{|F|:F\in\mathcal{F}\}.$$

### Proposition

Let G be a ribbon graph. Then

$$\varepsilon(G) = w(D(G)).$$

# Twist Polynomial

### Definition (Yan and Jin, 2021)

Let  $D = (E, \mathcal{F})$  be a delta-matroid. Then the *twist polynomial* of D is

$${}^{\partial}w_D(z) = \sum_{A \subseteq E} z^{w(D*A)}.$$

From the previous two propositions, it is immediate that if G is a ribbon graph, then  ${}^{\partial}\varepsilon_G(z)={}^{\partial}w_{D(G)}.$ 

# Binary Delta-Matroids

#### Definition

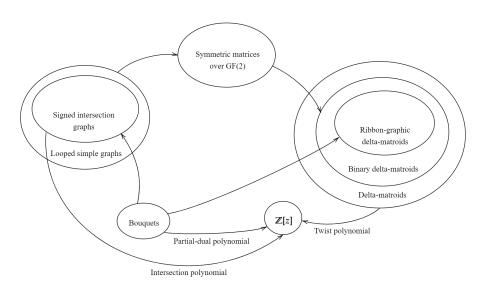
Let E be a finite set and M be a symmetric matrix over GF(2) with rows and columns indexed by E. Let  $D(M) = (E, \mathcal{F})$ , where

$$\mathcal{F} = \{A \subseteq E : M[A] \text{ is invertible}\}.$$

We take  $M[\emptyset]$  to be invertible. Then D(M) is a delta-matroid and D\*A for any  $A\subseteq E$  is called a *binary delta-matroid*.

 ${\it M}$  can be viewed as the adjacency matrix of a looped simple graph.

# A Commutative Diagram



# Adjacency Polynomial

#### **Definition**

Let G be a looped simple graph. We define the adjacency polynomial  $AP_G(z)$  to be the twist polynomial of the binary delta-matroid obtained from the adjacency matrix of G.

## Theorem (L.)

The cut-vertex recurrence relation for the intersection polynomial extends to the adjacency polynomial.