

Recurrence Relationships in Quasi-Tree Subgraphs

Logan Keck Rohan Mawalkar

June 27, 2025

Table of Contents

- 1 Matrix Quasi-Tree Theorem
- 2 Fibonacci and Lucas Numbers
- 3 Other Theorems
- 4 Recurrence in Ribbon Graphs without Matrices

Table of Contents

- 1 Matrix Quasi-Tree Theorem
- 2 Fibonacci and Lucas Numbers
- 3 Other Theorems
- 4 Recurrence in Ribbon Graphs without Matrices

Definition 1

A chord diagram consists of $2n$ distinct points on a circle and n chords connecting pairwise disjoint pairs of points.

Definition 1

A chord diagram consists of $2n$ distinct points on a circle and n chords connecting pairwise disjoint pairs of points. The intersection graph of a chord diagram D is the graph $G(D) = (V, E)$ where V is the set of chords and $uv \in E$ iff chords u and v intersect.

Definition 1

A chord diagram consists of $2n$ distinct points on a circle and n chords connecting pairwise disjoint pairs of points. The intersection graph of a chord diagram D is the graph $G(D) = (V, E)$ where V is the set of chords and $uv \in E$ iff chords u and v intersect.

Definition 2

A chord diagram is framed if each chord is assigned 0 or 1.

Definition 1

A chord diagram consists of $2n$ distinct points on a circle and n chords connecting pairwise disjoint pairs of points. The intersection graph of a chord diagram D is the graph $G(D) = (V, E)$ where V is the set of chords and $uv \in E$ iff chords u and v intersect.

Definition 2

A chord diagram is framed if each chord is assigned 0 or 1.

Definition 3

A bouquet is a ribbon graph with only one vertex.

Definition 3

A bouquet is a ribbon graph with only one vertex.

Given a bouquet B , we can create its associated framed chord diagram $F(B)$. First, consider the boundary of the vertex as the circle of the diagram. Draw a chord between the ends of each edge of B and assign it 0 iff its ribbon loop in B is orientable (untwisted).

Definitions

Given a framed chord diagram $D = \{\{a_i, b_i\} \mid 1 \leq i \leq n\}$, we can construct its intersection matrix $A(D)$. For each chord $\{a_i, b_i\}$, arbitrarily choose a pair: (a_i, b_i) or (b_i, a_i) . Each matrix entry $A_{i,j}$ is 0 if its chord $\{a_i, b_i\}$ is assigned 0, and 1 otherwise. For $i < j$, entry $A_{i,j}$ is 1 if the corresponding chords intersect with cyclic order a_i, a_j, b_i, b_j , -1 if they intersect with cyclic order a_i, b_j, b_i, a_j , and 0 if the chords do not intersect. For $i < j$, $A_{i,j} = -A_{j,i}$.

Theorem 1

Given an orientable bouquet B with n edges,

$$\kappa(B) = \det(I_n + A(F(B)))$$

Matrix Quasi-Tree Theorems

Theorem 1

Given an orientable bouquet B with n edges,

$$\kappa(B) = \det(I_n + A(F(B)))$$

Theorem 2

Given a bouquet B with n edges and exactly one non-orientable loop e_1 ,

$$\kappa(B) = \det(I_n + A(F(B)))$$

Table of Contents

- 1 Matrix Quasi-Tree Theorem
- 2 Fibonacci and Lucas Numbers**
- 3 Other Theorems
- 4 Recurrence in Ribbon Graphs without Matrices

Fibonacci Numbers

$$f_n = f_{n-1} + f_{n-2}$$

$$f_1 = 1, f_2 = 1$$

Fibonacci Numbers

$$f_n = f_{n-1} + f_{n-2}$$

$$f_1 = 1, f_2 = 1$$

Lucas Numbers

$$\ell_n = \ell_{n-1} + \ell_{n-2}$$

$$\ell_1 = 1, \ell_2 = 3$$

Fibonacci Numbers

$$f_n = f_{n-1} + f_{n-2}$$

$$f_1 = 1, f_2 = 1$$

Lucas Numbers

$$\ell_n = \ell_{n-1} + \ell_{n-2}$$

$$\ell_1 = 1, \ell_2 = 3$$

Relation

$$f_n + f_{n-2} = \ell_{n-1}$$

Special Bouquet Sequences

\mathbb{F}_n

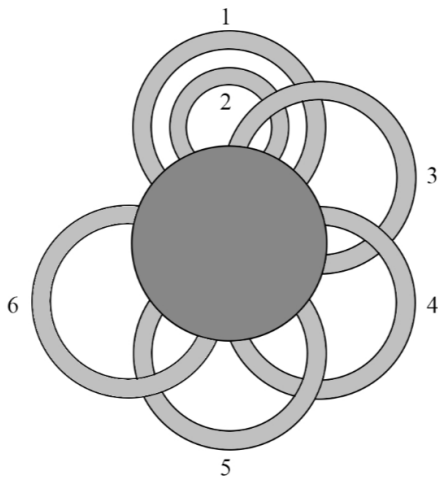
$\forall n \geq 0$, let \mathbb{F}_n denote the bouquet with the signed rotation $(1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$ whose chord diagram $D(\mathbb{F}_n)$ has pairs $\{(1, 3), (2, 5), (4, 7), \dots, (2n-4, 2n-1), (2n-2, 2n)\}$

Theorem

$$\kappa(\mathbb{F}_n) = f_{n+1}$$

\mathbb{F}'_n

$\forall n \geq 2$, let \mathbb{F}'_n denote the bouquet with the signed rotation $(1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$ whose chord diagram $D(\mathbb{F}'_n)$ has pairs $\{(1, 5), (2, 4), (3, 7), (6, 9), \dots, (2n-4, 2n-1), (2n-2, 2n)\}$



Theorem

Theorem

$$\forall n \geq 2, \kappa(\mathbb{F}'_n) = \ell_{n-1}$$

Theorem

Theorem

$$\forall n \geq 2, \kappa(\mathbb{F}'_n) = \ell_{n-1}$$

Proof

For $n = 2$, $\kappa(\mathbb{F}'_2) = 1 = \ell_1$ (the subgraph with no edges). Now, let $n \geq 3$.

$$I_n + A(F(\mathbb{F}'_n)) = \begin{pmatrix} 1 & 0 & 1 & \dots & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & \dots & 0 & 0 \\ -1 & -1 & 1 & 1 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 1 & 1 & 0 \\ 0 & 0 & 0 & \dots & -1 & 1 & 1 \\ 0 & 0 & 0 & \dots & 0 & -1 & 1 \end{pmatrix}, \text{ and by the Matrix}$$

Quasi-Tree Theorem, $\kappa(\mathbb{F}'_n) = \det(I_n + A(F(\mathbb{F}'_n)))$.

Proof Cont.

For $n = 2$, $\kappa(\mathbb{F}'_2) = 1 = \ell_1$ (the subgraph with no edges). Now, let $n \geq 3$.

$$I_n + A(F(\mathbb{F}'_n)) = \begin{pmatrix} 1 & 0 & 1 & \dots & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 & \dots & 0 & 0 \\ -1 & -1 & 1 & 1 & \dots & 0 & 0 \\ \dots & \dots & \dots & \dots & \dots & \dots & \dots \\ 0 & 0 & 0 & \dots & 1 & 1 & 0 \\ 0 & 0 & 0 & \dots & -1 & 1 & 1 \\ 0 & 0 & 0 & \dots & 0 & -1 & 1 \end{pmatrix}, \text{ and by the Matrix}$$

Quasi-Tree Theorem, $\kappa(\mathbb{F}'_n) = \det(I_n + A(F(\mathbb{F}'_n)))$.

Let $M := I_n + A(F(\mathbb{F}'_n))$.

$$\kappa(\mathbb{F}'_n) = \det(M) = \det(M[1, 1]) - \det(M[3, 1]) =$$

$$\kappa(\mathbb{F}_{n-1}) - (-1)m_{1,3}m_{2,2}\kappa(\mathbb{F}_{n-3}) = \kappa(\mathbb{F}_{n-1}) + \kappa(\mathbb{F}_{n-3}) = f_n + f_{n-2} = \ell_{n-1}$$

Theorem

Definition

Let $\mathbb{F}_n^{\prime 1}$ denote the bouquet of \mathbb{F}_n' with the first edge twisted (non-orientable). Its signed rotation is $(-1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$ whose chord diagram $D(\mathbb{F}_n')$ has pairs $\{(1, 5), (2, 4), (3, 7), (6, 9), \dots, (2n-4, 2n-1), (2n-2, 2n)\}$ and chord $(1, 5)$ is assigned 1 while all others 0.

Theorem

$$\forall n \geq 2, \kappa(\mathbb{F}_n^{\prime 1}) = f_n + \ell_{n-1}$$

Table of Contents

- 1 Matrix Quasi-Tree Theorem
- 2 Fibonacci and Lucas Numbers
- 3 Other Theorems**
- 4 Recurrence in Ribbon Graphs without Matrices

Other Theorems

Now, we'll try to understand these recurrence relations without the use of matrices.

Other Theorems

Now, we'll try to understand these recurrence relations without the use of matrices.

Lemma

Let $B(G)$ = the number of boundary components of G for any ribbon graph G . For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator, $B(P \vee Q) = B(P) + B(Q) - 1$

Other Theorems

Now, we'll try to understand these recurrence relations without the use of matrices.

Lemma

Let $B(G)$ = the number of boundary components of G for any ribbon graph G . For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator, $B(P \vee Q) = B(P) + B(Q) - 1$

Consequently,

Theorem 1

For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator,

$$\kappa(P \vee Q) = \kappa(P)\kappa(Q)$$

Subgraphs with More Boundary Components

Let $G = (V, E)$ be a ribbon graph, and let n be a positive integer.

Definition

$$\mathcal{F}_n(G) = \{F \subseteq E \mid (V, F) \text{ has } n \text{ boundary components}\}$$

Subgraphs with More Boundary Components

Let $G = (V, E)$ be a ribbon graph, and let n be a positive integer.

Definition

$$\mathcal{F}_n(G) = \{F \subseteq E \mid (V, F) \text{ has } n \text{ boundary components}\}$$

Then,

Definition

$f_n(G)$ is the cardinality of $\mathcal{F}_n(G)$

Subgraphs with More Boundary Components

Let $G = (V, E)$ be a ribbon graph, and let n be a positive integer.

Definition

$$\mathcal{F}_n(G) = \{F \subseteq E \mid (V, F) \text{ has } n \text{ boundary components}\}$$

Then,

Definition

$f_n(G)$ is the cardinality of $\mathcal{F}_n(G)$

Put simply, $f_n(G)$ counts the number of spanning subgraphs of G with exactly n boundary components (and, in particular, $f_1(G) = \kappa(G)$).

Generalization of Theorem 1

Now, using f_n rather than κ , we can generalize Theorem 1.

Generalization of Theorem 1

Now, using f_n rather than κ , we can generalize Theorem 1.

Theorem 2

For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator,

$$f_n(P \vee Q) = \sum_{i=1}^n f_i(P) f_{n+1-i}(Q)$$

Generalization of Theorem 1

Now, using f_n rather than κ , we can generalize Theorem 1.

Theorem 2

For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator,

$$f_n(P \vee Q) = \sum_{i=1}^n f_i(P) f_{n+1-i}(Q)$$

Given that graph P has k boundary components and that graph Q has j boundary components, we can conclude that $P \vee Q$ has $k + j - 1$ boundary components.

Generalization of Theorem 1

Now, using f_n rather than κ , we can generalize Theorem 1.

Theorem 2

For disjoint ribbon graphs P and Q , where \vee is the disjoint union operator,

$$f_n(P \vee Q) = \sum_{i=1}^n f_i(P) f_{n+1-i}(Q)$$

Given that graph P has k boundary components and that graph Q has j boundary components, we can conclude that $P \vee Q$ has $k + j - 1$ boundary components.

So, each subgraph of $P \vee Q$ with n boundary components must have i boundary components “contributed” by P , and $n + 1 - i$ boundary components “contributed” from Q for some i .

First Partial Duality Theorem

The next two theorems have to do with partial duality.

First Partial Duality Theorem

The next two theorems have to do with partial duality.

Theorem 3

Let G be a ribbon graph and $e \in E(G)$. Then, for any n ,

$$\begin{aligned} f_n(G) &= f_n(G \setminus e) + f_n(G/e) \\ &= f_n(G \setminus e) + f_n(G^{\delta(e)} \setminus e) \end{aligned}$$

First Partial Duality Theorem

The next two theorems have to do with partial duality.

Theorem 3

Let G be a ribbon graph and $e \in E(G)$. Then, for any n ,

$$\begin{aligned}f_n(G) &= f_n(G \setminus e) + f_n(G/e) \\ &= f_n(G \setminus e) + f_n(G^{\delta(e)} \setminus e)\end{aligned}$$

We can find an obvious 1-1 correspondence between subgraphs of G that do not contain e and subgraphs of $G \setminus e$.

There is a slightly less obvious 1-1 correspondence between subgraphs of G that *do* contain e and subgraphs of G/e .

First Partial Duality Theorem cont'd

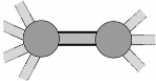
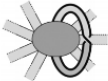
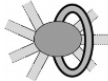

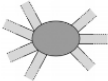
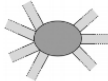
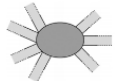
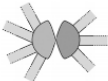

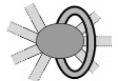
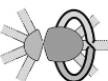
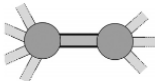
	Non-loop	Non-orientable loop	Orientable loop
G			
$G \setminus e$			
$G/e = G^{\delta(e)} \setminus e$			
$G^{\delta(e)}$			

Figure: Table of Partial Duals

Second Partial Duality Theorem

Another theorem involving partial duality falls from the first.

Second Partial Duality Theorem

Another theorem involving partial duality falls from the first.

Theorem 4

Let $G = (V, E)$ be a ribbon graph, $A \subseteq E$. Then, for any n ,

$$f_n(G) = f_n(G^{\delta(A)})$$

Second Partial Duality Theorem

Another theorem involving partial duality falls from the first.

Theorem 4

Let $G = (V, E)$ be a ribbon graph, $A \subseteq E$. Then, for any n ,

$$f_n(G) = f_n(G^{\delta(A)})$$

It is sufficient to verify this for a single edge. Let $e \in E$. Then,

$$\begin{aligned} f_n(G) &= f_n(G \setminus e) + f_n(G/e) \\ &= f_n\left((G^{\delta(e)})^{\delta(e)} \setminus e\right) + f_n(G^{\delta(e)} \setminus e) \\ &= f_n(G^{\delta(e)}) \end{aligned}$$

Table of Contents

- 1 Matrix Quasi-Tree Theorem
- 2 Fibonacci and Lucas Numbers
- 3 Other Theorems
- 4 Recurrence in Ribbon Graphs without Matrices

Sequences of Graphs

Let's remind ourselves of a few relevant sequences of ribbon graphs:

Sequences of Graphs

Let's remind ourselves of a few relevant sequences of ribbon graphs:

$$\mathbb{F}_n = (1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}'_n = (1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}^1_n = (-1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}'^1_n = (-1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

Sequences of Graphs

Let's remind ourselves of a few relevant sequences of ribbon graphs:

$$\mathbb{F}_n = (1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}'_n = (1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}_n^1 = (-1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}_n'^1 = (-1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

Where we have:

$$\kappa(\mathbb{F}_n) = f_{n+1}$$

$$\kappa(\mathbb{F}'_n) = \ell_{n-1}$$

$$\kappa(\mathbb{F}_n^1) = f_{n+2}$$

$$\kappa(\mathbb{F}_n'^1) = f_n + \ell_{n-1}$$

Sequences of Graphs

Let's remind ourselves of a few relevant sequences of ribbon graphs:

$$\mathbb{F}_n = (1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}'_n = (1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}^1_n = (-1, 2, 1, 3, 2, 4, 3, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

$$\mathbb{F}'^1_n = (-1, 2, 3, 2, 1, 4, 3, 5, 4, \dots, i, i-1, i+1, i, \dots, n-1, n-2, n, n-1, n)$$

Where we have:

$$\kappa(\mathbb{F}_n) = f_{n+1}$$

$$\kappa(\mathbb{F}'_n) = \ell_{n-1}$$

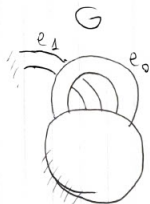
$$\kappa(\mathbb{F}^1_n) = f_{n+2}$$

$$\kappa(\mathbb{F}'^1_n) = f_n + \ell_{n-1}$$

Fact

Let $G = (E, V)$ containing $e_0, e_1 \in E$ and $v \in V$ s.t. the signed rotation of v contains the sequence (e_0, e_1, e_0) and e_0 is a loop. Then

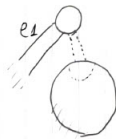
$$\kappa(G) = \kappa(G \setminus e_0) + \kappa(G \setminus e_0 \setminus e_1)$$



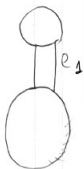
$G \setminus e_0$



G/e_0



$$K(G) = K(G \setminus e_0) + K(G/e_0)$$



G/e_0

$$= (G \setminus e_0 \setminus e_1) \vee P_2$$

$$K(G/e_0) = K((G \setminus e_0 \setminus e_1) \vee P_2) = K(G \setminus e_0 \setminus e_1) \cdot 1$$

$$(K(P_2) = 1)$$

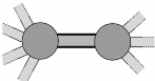
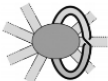
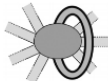

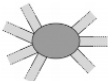
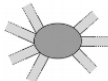
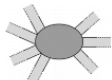
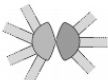

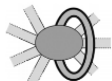
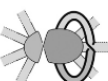
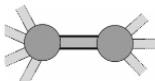
	Non-loop	Non-orientable loop	Orientable loop
G			
$G \setminus e$			
$G/e = G^{\delta(e)} \setminus e$			
$G^{\delta(e)}$			

Figure: Table of Partial Duals

Another Interesting Sequence

Consider the sequence,

$$\mathbb{F}_n^r = (1, 2, -1, 3, -2, 4, -3, \dots, i, 1-i, i+1, -i, \dots, n-1, 2-n, n, 1-n, -n)$$

(which is just \mathbb{F}_n with all ribbons twisted)

Another Interesting Sequence

Consider the sequence,

$$\mathbb{F}_n^r = (1, 2, -1, 3, -2, 4, -3, \dots, i, 1-i, i+1, -i, \dots, n-1, 2-n, n, 1-n, -n)$$

(which is just \mathbb{F}_n with all ribbons twisted)

An interesting result is that

Fact

$$\kappa(\mathbb{F}_n^r) = \kappa(\mathbb{F}_{n-1}^r) + \kappa(\mathbb{F}_{n-2}^r) + \kappa(\mathbb{F}_{n-3}^r)$$

Another Interesting Sequence

Consider the sequence,

$$\mathbb{F}_n^r = (1, 2, -1, 3, -2, 4, -3, \dots, i, 1-i, i+1, -i, \dots, n-1, 2-n, n, 1-n, -n)$$

(which is just \mathbb{F}_n with all ribbons twisted)

An interesting result is that

Fact

$$\kappa(\mathbb{F}_n^r) = \kappa(\mathbb{F}_{n-1}^r) + \kappa(\mathbb{F}_{n-2}^r) + \kappa(\mathbb{F}_{n-3}^r)$$

More generally, any sequence with such an ending will follow this recurrence relation.

$$\kappa(\mathbb{F}_n^r) = \kappa(\mathbb{F}_n^r \setminus n) + \kappa(\mathbb{F}_n^r / n)$$

Of course, $\mathbb{F}_n^r \setminus n = \mathbb{F}_{n-1}^r$. But also, by our previous result,

$$\begin{aligned} \kappa(\mathbb{F}_n^r / n) &= \kappa(\mathbb{F}_n^r / n \setminus (n-1)) + \kappa(\mathbb{F}_n^r / n \setminus (n-1) \setminus (n-2)) \\ &= \kappa(\mathbb{F}_n^r \setminus n \setminus (n-1)) + \kappa(\mathbb{F}_n^r \setminus n \setminus (n-1) \setminus (n-2)) \\ &= \kappa(\mathbb{F}_{n-2}^r) + \kappa(\mathbb{F}_{n-3}^r) \end{aligned}$$

Another interesting result arises from considering f_n where $n > 1$.

Another interesting result arises from considering f_n where $n > 1$. Consider \mathbb{F}_n (or any sequence with such an ending). We can use our theorems to deduce that $f_2(\mathbb{F}_n) = f_2(\mathbb{F}_{n-1}) + f_2(\mathbb{F}_{n-2}) + f_1(\mathbb{F}_{n-2})$.

Another interesting result arises from considering f_n where $n > 1$. Consider \mathbb{F}_n (or any sequence with such an ending). We can use our theorems to deduce that $f_2(\mathbb{F}_n) = f_2(\mathbb{F}_{n-1}) + f_2(\mathbb{F}_{n-2}) + f_1(\mathbb{F}_{n-2})$. Using the recurrence relation on $f_1(\mathbb{F}_n)$, we can find that

$$f_2(\mathbb{F}_n) = 2f_2(\mathbb{F}_{n-1}) + f_2(\mathbb{F}_{n-2}) - 2f_2(\mathbb{F}_{n-3}) - f_2(\mathbb{F}_{n-4})$$

Another interesting result arises from considering f_n where $n > 1$. Consider \mathbb{F}_n (or any sequence with such an ending). We can use our theorems to deduce that $f_2(\mathbb{F}_n) = f_2(\mathbb{F}_{n-1}) + f_2(\mathbb{F}_{n-2}) + f_1(\mathbb{F}_{n-2})$. Using the recurrence relation on $f_1(\mathbb{F}_n)$, we can find that

$$f_2(\mathbb{F}_n) = 2f_2(\mathbb{F}_{n-1}) + f_2(\mathbb{F}_{n-2}) - 2f_2(\mathbb{F}_{n-3}) - f_2(\mathbb{F}_{n-4})$$

Interestingly, the characteristic polynomial of this recurrence relation is $x^4 - 2x^3 - x^2 + 2x + 1 = (x^2 - x - 1)^2$, where $x^2 - x - 1$ is the characteristic polynomial of $f_1(\mathbb{F}_n)$.

Q. Deng, X. Jin, Q. Yan, The number of quasi-trees of bouquets with exactly one non-orientable loop, Preprint arXiv:2406.11648v1 [math.CO]