

## Examples Class 1

4)  $\mathcal{L} = \{+, 0\}$ . Let  $T$  be the theory of nontrivial torsion free divisible abelian groups.

$$TF: \forall x (nx=0 \rightarrow x=0) \quad nx \text{ is } x+\dots+x \text{ (n times)}$$

$$D: \forall x \exists y (x=ny)$$

Claim:  $T$  is  $\kappa$ -categorical  $\forall \kappa > \aleph_0$ .

Proof: Models of  $T$  "are" vector spaces over  $\mathbb{Q}$ .

NTTFDAG  $G$ ,  $a \in G$ ,  $\lambda = \frac{m}{n} \in \mathbb{Q}$ . Define  $\lambda a$  to be the unique solution of  $ma = nx$  in  $G$ .

$M, N \models T$  are isomorphic iff  $\dim(M) = \dim(N)$   $\star$

If  $M \models T$  is uncountable then  $\dim(M) = |M|$ .  $\square$

7)  $\mathcal{N}$ ,  $A \subseteq \mathcal{N}$ . Build substructure generated by  $A$  "by hand"

Remark: If  $M \subseteq \mathcal{N}$  and  $A \subseteq M$ , then for any  $\mathcal{L}$ -term  $t(x_1, \dots, x_n)$  and  $\bar{a} \in A^n$ ,  $t^M(\bar{a}) = t^{\mathcal{N}}(\bar{a})$  (see Claim in proof of Thm 2.2)

Define  $M = \{t^M(\bar{a}) : \text{any } \mathcal{L}\text{-term } t, \text{ any } \bar{a} \text{ from } A\}$ . Define  $M$

• constant symbol  $c$  :  $\underline{c^M} = c^{\mathcal{N}} \in M$

•  $n$ -ary function symbol  $f$ : Given  $b_1, \dots, b_n \in M$ , define  $f^M(b_1, \dots, b_n)$ .

Fix terms  $t_1, \dots, t_n$  and tuples  $\bar{a}_1, \dots, \bar{a}_n$  from  $A$  st  $b_i = t_i^{\mathcal{N}}(\bar{a}_i)$

Let  $\underline{f^M(b_1, \dots, b_n)} = \underline{f^{\mathcal{N}}(t_1^{\mathcal{N}}(\bar{a}_1), \dots, t_n^{\mathcal{N}}(\bar{a}_n))} = t^{\mathcal{N}}(\bar{a}_1, \dots, \bar{a}_n) \in M$

where  $t$  is  $f(t_1, \dots, t_n)$ .

•  $n$ -ary relation symbol  $R$  :  $\underline{R^M} := R^{\mathcal{N}} \cap M^n$

By construction  $M \subseteq \mathcal{N}$ .  $A \subseteq M$ . It's the smallest by the Remark.

8) Chain  $(M_i)_{i < \alpha}$ . Define  $N = \bigcup_{i < \alpha} M_i$ . Define  $N = \bigcup_{i < \alpha} M_i$ .

• constant  $c$ :  $\forall i < j < \alpha$ ,  $c^{M_i} = c^{M_j}$  since  $M_i \subseteq M_j$ .

Let  $c^N = c^{M_i}$ .

• n-ary  $f$ : Given  $a_1, \dots, a_n \in N$ . Choose  $i < \alpha$  st  $a_1, \dots, a_n \in M_i$

$\forall j > i$ ,  $f^{M_i}(\bar{a}) = f^{M_j}(\bar{a})$  since  $M_i \subseteq M_j$ .

Define  $f^N(\bar{a}) = f^{M_i}(\bar{a})$  for some/any  $i$  st  $a_1, \dots, a_n \in M_i$ .

• n-ary relation  $R$ :  $\forall i < j < \alpha$   $R^{M_i} = R^{M_j} \cap M_i^n$  since  $M_i \subseteq M_j$ .

Define  $R^N = \bigcup_{i < \alpha} R^{M_i}$ . Then  $\forall i$ ,  $R^N \cap M_i^n = R^{M_i}$ .

a) Clear from construction that  $M_i \subseteq N \forall i < \alpha$ .

b) Assume  $(M_i)_{i < \alpha}$  is elementary. Fix  $\phi(x_1, \dots, x_n)$   $\mathcal{L}$ -formula

★ WTS:  $\forall i < \alpha \forall \bar{a} \in M_i^n$ ,  $M_i \models \phi(\bar{a}) \iff N \models \phi(\bar{a})$ .

Induction on  $\phi$ :  $\phi$  is atomic, use  $M_i \subseteq N$ .  $\wedge, \neg$  clear.

Assume the result for  $\phi(x_1, \dots, x_n, y)$ . Consider  $\exists y \phi(\bar{x}, y)$ .

Fix  $i < \alpha$ . Fix  $\bar{a} \in M_i^n$ . Suppose  $M_i \models \exists y \phi(\bar{a}, y)$ . Then  $\exists b \in M_i$

st  $M_i \models \phi(\bar{a}, b)$ . Then  $\exists b \in N$  st  $N \models \phi(\bar{a}, b)$  by induction.

So  $N \models \exists y \phi(\bar{a}, y)$ .

Suppose  $N \models \exists y \phi(\bar{a}, y)$ . Fix  $b \in N$  st  $N \models \phi(\bar{a}, b)$ . Fix  $j \geq i$

st  $b \in M_j$ . Then  $M_j \models \phi(\bar{a}, b)$  by induction. So  $M_j \models \exists y \phi(\bar{a}, y)$ .

Since  $M_i \subseteq M_j$ , we have  $M_i \models \exists y \phi(\bar{a}, y)$ . □.

⑨  $M, A \subseteq M$ . Want  $N \subseteq M$  st  $A \subseteq N$  +  $|N| \leq |A| + |Z| + \aleph_0$

Attempt #1 Let  $N$  be the substructure generated by  $A$ .  $|N| \leq |A| + |Z| + \aleph_0$  (by #7)

Only know  $N \subseteq M$ .

Ex: Let  $M$  be infinite set in  $Z = \emptyset$ . Let  $A \subseteq M$  be finite.

Then  $N = A$ .  $M \neq N$ .

Attempt #2 Apply DLST to  $\text{Th}_A(M)$  to get  $N \models \text{Th}_A(M)$  with

$|N| \leq |A| + |Z| + \aleph_0$ . There may be no elementary embedding of  $N$  into  $M$  fixing  $A$ .

Ex:  $M = (\mathbb{Q}, +, 0)$   $A = \{0\}$   $N = (\mathbb{Q}^2, +, 0)$

$N$  does not embed into  $M$ .

Proof of 9:  $Z_0 = Z_A$ . Expand  $M$  to an  $Z_0$ -structure.

Given  $Z_k$  and expansion of  $M$  to  $Z_k$ -structure.

Let  $Z_{k+1} = Z_k \cup \left\{ C_{\phi, \bar{a}} : \begin{array}{l} \phi(x_1, \dots, x_n, y) \text{ is an } Z_k\text{-formula} \\ \bar{a} \in A^n \text{ st } M \models \exists y \phi(\bar{a}, y) \end{array} \right\}$

Expand  $M$  to  $Z_{k+1}$ -structure st  $M \models \phi(\bar{a}, C_{\phi, \bar{a}}^M)$ .

Let  $Z^* = \bigcup Z_k$ . Let  $N$  be the  $Z^*$ -substructure of  $M$  generated by  $A$ .

$A \subseteq N$ .  $|N| \underset{\#7}{\leq} |A| + |Z^*| + \aleph_0 \leq |A| + |Z| + \aleph_0$ .

Fix  $Z$ -formula  $\phi(x_1, \dots, x_n)$  and  $b_1, \dots, b_n \in N$

WTS:  $M \models \phi(\bar{b})$  iff  $N \models \phi(\bar{b})$ .

Induction on  $\phi$ : atomic is clear since  $N \subseteq M$ ,  $\wedge, \neg$  easy.

Assume the result for  $\phi(x_1, \dots, x_n, y)$ . Consider  $\exists y \phi(\bar{x}, y)$ .

Fix  $b_1, \dots, b_n \in N$ .  $N \models \exists y \phi(\bar{b}, y) \Rightarrow M \models \exists y \phi(\bar{b}, y)$  by induction.

Suppose  $M \models \exists y \phi(\bar{b}, y)$ . Choose  $Z^*$ -terms  $t_1, \dots, t_n$  st  $b_i = t_i^M(\bar{a}_i)$

For some tuple  $\bar{a}_i$  from  $A$ . let  $\Psi(\bar{v}_1, \dots, \bar{v}_n, y)$  be

$$\Phi(t_1(\bar{v}_1), \dots, t_n(\bar{v}_n), y) \leftarrow \mathcal{L}^* \text{-formula.}$$

$$\underline{M \models \exists y \Psi(\bar{a}_1, \dots, \bar{a}_n, y)}. \text{ So if } c = c_{\Psi, \bar{a}}^M, \text{ then } \underline{M \models \Psi(\bar{a}, c)}$$

$$\text{So } \underline{M \models \Phi(\bar{b}, c)}. \text{ So } N \models \Phi(\bar{b}, c) \text{ (since } c \in N).$$

$$\text{So } N \models \exists y \Phi(\bar{b}, y)$$

⑥  $M, N, M \subseteq N, h: M \rightarrow N$  inclusion.

WTS:  $h$  is an  $\mathcal{L}$ -embedding (i.e.  $M \subseteq N$ ) iff  $h$  preserves g.f. formulas.

( $\Rightarrow$ ): By proof of Thm 2.2 (only needed surjectivity for quantifier step)

( $\Leftarrow$ ): constant symbol  $c: \Phi(x)$  be  $x=c$

$$M \models \Phi(c^M) \text{ so } N \models \Phi(c^M) \text{ i.e. } c^M = c^N$$

$n$ -ary  $f$ : let  $\Phi(x_1, \dots, x_n, y)$  be  $y = f(x_1, \dots, x_n)$

$$\text{Given } \bar{a} \in M^n, M \models \Phi(\bar{a}, f^M(\bar{a})). \text{ So } N \models \Phi(\bar{a}, f^M(\bar{a}))$$

$$\text{i.e. } f^M(\bar{a}) = f^N(\bar{a}).$$

$n$ -ary  $R$ : let  $\Phi(x_1, \dots, x_n)$  be  $R(x_1, \dots, x_n)$

$$\text{Given } \bar{a} \in M^n, \bar{a} \in R^M \text{ iff } M \models \Phi(\bar{a}) \text{ iff } N \models \Phi(\bar{a})$$

$$\text{iff } \bar{a} \in R^N. \text{ So } R^M = R^N \cap M^n$$

⑤  $M = (V, E)$  is a graph which is finitely  $k$ -colorable.

$\mathcal{L}^0 = \{E\}, \mathcal{L} = \{E, C_1, \dots, C_k\}$  each  $C_i$  is a unary relation symbol.

let  $\sigma$  say " $C_1, \dots, C_k$  partition the universe + yield a  $k$ -coloring" i.e.

$$\forall x \bigvee_{i=1}^k C_i(x) \wedge \forall x \bigwedge_{i \neq j} \neg (C_i(x) \wedge C_j(x)) \wedge \forall x \forall y (E(x, y) \rightarrow \bigwedge_{i=1}^k \neg (C_i(x) \wedge C_i(y)))$$

let  $T = \mathcal{D}(M) \cup \{\text{graph axioms}\} \cup \{\sigma\}$

A model of  $T$  is a  $k$ -colored graph which has  $M$  as a subgraph.

Fix  $\Sigma \subseteq T$  finite. Let  $A$  be the set of elements in  $V$  whose constant symbols occur in  $\Sigma$ .  $A$  is a finite subgraph of  $M$ .

Expand  $A$  to an  $\mathcal{L}_V$ -structure satisfying  $\Sigma$ . By Compactness  $T$  has a model.

## ADDED AFTER CLASS

---

① Let  $\Delta = \{\sigma_n : n \geq 1\}$  where  $\sigma_n$  says "there are at least  $n$  distinct elements."

By assumption,  $T \cup \Delta$  is finitely satisfiable.

By Compactness,  $T \cup \Delta$  has a model.

---

② (i)  $\Leftrightarrow$  (ii) by definition.

$\neg(i) \Rightarrow \neg(iii)$  If  $T$  is not complete then there is an  $\mathcal{L}$ -sentence  $\phi$  st  $T \not\models \phi$  and  $T \not\models \neg\phi$ . So there is  $M \models T$  st  $M \models \neg\phi$  and  $N \models T$  st  $N \models \phi$ . So  $M \not\equiv N$ .

$\neg(iii) \Rightarrow \neg(i)$ . Suppose  $M, N \models T$  and  $M \not\equiv N$ . Then there is an  $\mathcal{L}$ -sentence  $\phi$  st  $M \models \phi$  and  $N \models \neg\phi$ . So  $T \not\models \neg\phi$  and  $T \not\models \phi$ .

---

③ If  $\mathcal{L} = \emptyset$  then " $\mathcal{L}$ -isomorphism" = "bijection."

So if  $M, N$  are infinite and  $|M| = |N|$ , then  $M \cong N$ .

By Vaught's Test, the theory of infinite sets is complete.