

Examples Class 3

1) T complete, cttle language.

a) $M \models T$ is atomic. WTS M is \aleph_0 -hom.

Fix $\bar{a}, \bar{b} \in M^n$ with $\bar{a} \equiv_{\emptyset} \bar{b}$. Fix $c \in M$. $tp(\bar{a}, c)$ is isolated by some $\phi(\bar{y}, x)$. $M \models \exists x \phi(\bar{a}, x)$. So $M \models \exists x \phi(\bar{b}, x)$. Choose $d \in M$ st $M \models \phi(\bar{b}, d)$. Claim $(\bar{a}, c) \equiv_{\emptyset} (\bar{b}, d)$

$\phi(\bar{y}, x) \in tp(\bar{b}, d)$ ie $tp(\bar{b}, d) \in [\phi(\bar{y}, x)] = \{tp(\bar{a}, c)\}$.

b) $M, N \models T$ cttle, \aleph_0 -hom., realize the same types over \emptyset .

$M = \{a_n : n \geq 1\}$, $N = \{b_n : n \geq 1\}$. $\mathcal{F}_0 = \emptyset$ is partial elementary since $M \equiv N$.

Suppose we have \mathcal{F}_n . Add a_{n+1} and b_{n+1} . Enumerate $\text{dom}(\mathcal{F}_n) = \bar{c}$ and $\text{Im}(\mathcal{F}_n) = \bar{d}$ (so $\bar{c} \equiv_{\emptyset} \bar{d}$).

$[(\bar{c}, a_{n+1}) \equiv_{\emptyset} (\bar{d}, b) \text{ want for some } b \in N]$

$tp(\bar{c}, a_{n+1})$ is realized in M . So $\exists \bar{d}', b'$ in N st

$$(\bar{c}, a_{n+1}) \equiv_{\emptyset} (\bar{d}', b')$$

We know $\bar{d}' \equiv_{\emptyset} \bar{c} \equiv_{\emptyset} \bar{d}$. So $\bar{d}' \equiv_{\emptyset} \bar{d}$

Since N is \aleph_0 -hom. $\exists b \in N$ st $(\bar{d}', b) \equiv_{\emptyset} (\bar{d}, b)$

So $(\bar{c}, a_{n+1}) \equiv_{\emptyset} (\bar{d}, b)$.

Similarly find $a \in M$, st $(\bar{c}, a_{n+1}, a) \equiv_{\emptyset} (\bar{d}, b, b_{n+1})$

Let $\mathcal{F}_{n+1} = \mathcal{F}_n \cup \{(a_{n+1}, b), (a, b_{n+1})\}$.

Ex: Suppose $1 < I(T, \aleph_0) < 2^{\aleph_0}$. Then T has a prime model M and a cttle saturated model N . $M \not\equiv N$. Both are \aleph_0 -hom.

c) $M, N \models T$ are prime M, N are stable + atomic

M, N are \aleph_0 -hom. by (a). M, N realize the same types. So $M \cong N$ by (b).

4) M, N saturated, $M \equiv N$, $|M| = |N| = \kappa$. Then $M \cong N$.

PF: Enumerate $M = \{a_i\}_{i < \kappa}$ $N = \{b_i\}_{i < \kappa}$.
↑ ordinal

Build $f_0 \subseteq f_1 \subseteq \dots \subseteq f_i \subseteq \dots$ for $i < \kappa$. $|\text{dom}(f_i)| < \kappa$.

$f_0 = \emptyset$. $\alpha < \kappa$ - limit: $f_\alpha = \bigcup_{i < \alpha} f_i$. $|\text{dom}(f_\alpha)| < \kappa$.

$$|\text{dom}(f_i)| = \begin{cases} < \aleph_0 & i < \aleph_0 \\ \aleph_0 & i \geq \aleph_0 \end{cases}$$

Given f_i , let $p = \text{tp}(a_i / \text{dom}(f_i)) \in S_1^M(\text{dom}(f_i))$.

let $q = f_i(p) \in S_1^N(\text{Im}(f_i))$.

let $b \in N$ realize q . Then $f_i \cup \{(a_i, b)\}$ is partial elementary.

Go back to find $a \in M$ st $f_i \cup \{(a_i, b), (a, b_i)\}$ is partial elementary.
↑ f_{i+1}

let $f = \bigcup_{i < \kappa} f_i$. Then $f: M \rightarrow N$ is an isomorphism.

3) M is κ -saturated. $A \in M$, $|A| < \kappa$. Then M is κ -saturated as an \mathcal{L}_A -structure.

$B \in M$, $|B| < \kappa$. $p \in S_1^{\text{Th}_A(M)}(B)$

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$\hat{p} \in S_1^{\text{Th}(M)}(A \cup B)$ $|A \cup B| < \kappa$.

2) $k \geq 1$. $\mathcal{L} = \{<, U_1, \dots, U_k, c_0, c_1, c_2, \dots\}$

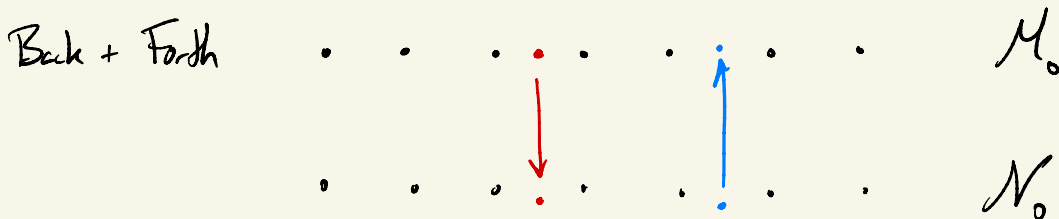
T says DLO + " U_1, \dots, U_k partition into dense sets" + " $c_0 < c_1 < c_2 < \dots$ "
+ " $c_n \in U_i, \forall n$ "

T is complete: Fix $M, N \models T$ cld. WTS $M \equiv N$.

Fix $n \geq 1$. Let M_0, N_0 be reducts to $\{<, U_1, \dots, U_k, c_0, c_1, \dots, c_n\}$

We show $M_0 \equiv N_0$. In fact, $M_0 \cong N_0$.

Start with $\bar{c}^{M_0} \rightarrow \bar{c}^{N_0}$.



$I(T, \mathcal{N}_0) = k+2$.

PF: For $1 \leq t \leq k$, we have M_t st $\sup c_n$ exists + satisfies U_t .

M_{k+1} : (c_n) is bounded above, no sup in model

M_{k+2} : (c_n) is unbounded.

For any $M \models T$ cld $\exists 1 \leq t \leq k+2$ st $M \cong M_t$.

⑤ $T, \kappa \geq |\mathcal{L}| + \aleph_0$.

Assume T is κ -stable. ($|S_1(M)| = \kappa$ for any $M \models T, |M| = \kappa$).

Fix $n \geq 1, M \models T, A \in M, |A| \leq \kappa$. WTS $|S_n(A)| \leq \kappa$.

Note: $\exists N \leq M$ st $A \in N + |N| \leq \kappa, |S_n(A)| \leq |S_n(N)|$

wlog $A = M, |M| = \kappa$.

Claim: $\forall M \models T, |M| = \kappa$, we have $|S_n(M)| = \kappa$.

PF: Induction on n . $n=1 \checkmark |S_1(M)| = \kappa$

Fix $N \cong M$ st $|N| = \kappa$ and N realizes all types in $S_1(M)$.

Given $p \in S_{n+1}(M)$, let $p_0 = \{ \varphi(x_{n+1}) : \varphi(x_{n+1}) \wedge \bar{x} = \bar{x} \in p \}$

So $p_0 \in S_1(M)$. Choose $a_p \in N$ st $a_p \models p$.

Let $q_p \in S_n(N)$ be any extension of $p(x_1, \dots, x_n, a_p)$

Show $p \mapsto (a_p, q_p)$ is injective. So $|S_{n+1}(M)| \leq |N \times S_n(N)| \leq \kappa$.

Suppose $(a_{p_1}, q_{p_1}) = (a_{p_2}, q_{p_2})$

$$\varphi(x_1, \dots, x_{n+1}) \in p_1 \iff \varphi(x_1, \dots, x_n, a_{p_1}) \in q_{p_1}$$

$$\iff \varphi(x_1, \dots, x_n, a_{p_2}) \in q_{p_2}$$

$$\iff \varphi(x_1, \dots, x_{n+1}) \in p_2.$$

⑥ Compactness.

$$\Gamma(\bar{x}_n, \bar{y}_n)_{n \geq 1} = \{ \varphi(\bar{x}_i, \bar{y}_j) : i \leq j \} \cup \{ \neg \varphi(\bar{x}_i, \bar{y}_j) : i > j \}.$$

$\Gamma \cup T$ is fin. satisfiable.

So $\exists M \models T$ and $(\bar{a}_n)_{n \geq 1}, (\bar{b}_n)_{n \geq 1}$ from M

$$\text{st } M \models \varphi(\bar{a}_i, \bar{b}_j) \iff i \leq j.$$

⑦ $p \in S_n(M) \quad \phi(\bar{x}, \bar{y}), \psi(\bar{x}, \bar{z})$

If $\chi(\bar{y})$ is a ϕ -def. for p and $\theta(\bar{z})$ is a ψ -def. for p .

Then $\neg \chi(\bar{y})$ is a $(\neg \phi)$ -def. for p and

$\chi(\bar{y}) \wedge \theta(\bar{z})$ is a $(\phi \wedge \psi)$ -def. for p .

⑧ $T = \text{Th}(\mathbb{Z}, <)$ $p \in S_n(\mathbb{Z})$ Then p is definable.

$\boxed{M \cong \mathbb{Z} \quad \longleftrightarrow \longleftrightarrow \longleftrightarrow \longleftrightarrow}$

Fix x_1, \dots, x_n . Atomic formulas $\phi(\bar{x}, \bar{y})$ are

1. $x_i = s^\alpha(y)$ or 2. $x_i < s^\alpha(y)$ $\alpha \in \mathbb{Z}$

[e.g. $x_i > s^\alpha(y) \Leftrightarrow \neg (x_i < s^\alpha(y) \vee x_i = s^\alpha(y))$]

Case 1: $x_i = a \in p$ for some $a \in \mathbb{Z}$.

$a = s^\alpha(y)$ def. for 1.

$a < s^\alpha(y)$ def. for 2.

Case 2: $x_i \neq a \in p$ for all $a \in \mathbb{Z}$.

For 1, use $y \neq y$ then $x_i = s^\alpha(b) \in p$ iff $b \neq y$

For 2. case 2a. $x_i > a \in p \quad \forall a \in \mathbb{Z}$. Use $y \neq y$.

case 2b. $x_i < a \in p \quad \forall a \in \mathbb{Z}$. Use $y = y$.
