

Part III Model Theory, Lecture 1, 12 Oct

\mathcal{L} is a fixed language.

An \mathcal{L} -theory T is finitely satisfiable if every finite subset of T is satisfiable.

Compactness Theorem satisfiable \Leftrightarrow finitely satisfiable

Downward Löwenheim-Skolem Theorem Any satisfiable \mathcal{L} -theory has a model of cardinality at most $|\mathcal{L}| + \aleph_0$

see notes

non-examinable

Löwenheim-Skolem Theorem

Suppose T is an \mathcal{L} -theory with infinite models. Then T has a model of cardinality κ for any $\kappa \geq |\mathcal{L}| + \aleph_0$.

Proof

Let $\mathcal{L}^* = \mathcal{L} \cup \{c_i : i \in \kappa\}$ where each c_i is a new constant symbol.

Let $T^* = T \cup \{c_i \neq c_j : i \neq j\}$. Suppose $\Sigma \in T^*$ is finite.

So $\Sigma \subseteq T \cup \{c_i \neq c_j : i, j \in I\}$ for some finite set I .

Let $\mathcal{M} \models T$ be infinite (\mathcal{M} is an \mathcal{L} -structure). Expand \mathcal{M} to an \mathcal{L}^* -structure \mathcal{M}^* by interpreting $c_i^{\mathcal{M}^*}$ as distinct elements for $i \in I$, and interpreting $c_i^{\mathcal{M}^*}$ for $i \notin I$ arbitrarily. Then $\mathcal{M}^* \models \Sigma$.

By Compactness, T^* is satisfiable. By DLST, T^* has a model

\mathcal{N}^* of cardinality at most $|\mathcal{L}^*| + \aleph_0 = \kappa$. So \mathcal{N}^* has

cardinality κ . Let \mathcal{N} be the reduct of \mathcal{N}^* to \mathcal{L} .

Then $\mathcal{N} \models T$ and $|\mathcal{N}| = \kappa$. \square

Complete Theories

Def 1.1 Let T be an \mathcal{L} -theory and ϕ an \mathcal{L} -sentence.

Then $T \models \phi$ ("T models ϕ ", "T implies ϕ ") if any model of T is a model of ϕ .

Example 1.2

1) $\{\phi, \psi\} \models \phi \wedge \psi$

2) If T is consistent then $T \models \exists x (x=x)$. So $\emptyset \models \exists x (x=x)$
(ie, satisfiable)

3) Let T be the theory of groups in $\mathcal{L} = \{*, e\}$

$$T \models \forall x \forall y \forall z ((x * y = e \wedge x * z = e) \rightarrow y = z)$$

Def 1.3 An \mathcal{L} -theory T is complete if, for any \mathcal{L} -sentence ϕ ,
 $T \models \phi$ or $T \models \neg \phi$.

Example 1.4

1) The theory of groups is not complete. Consider $\forall x \forall y (x * y = y * x)$

2) ZFC is not complete. Consider the Continuum Hypothesis.

Def 1.5 Let \mathcal{M} be an \mathcal{L} -structure. The theory of \mathcal{M} is

$$\text{Th}(\mathcal{M}) = \text{Th}_{\mathcal{L}}(\mathcal{M}) := \{\phi : \phi \text{ is an } \mathcal{L}\text{-sentence} \wedge \mathcal{M} \models \phi\}$$

Note that $\text{Th}(\mathcal{M})$ is complete.

Def 1.6 Two \mathcal{L} -structures \mathcal{M} and \mathcal{N} are elementarily equivalent, written $\mathcal{M} \equiv \mathcal{N}$,
if $\text{Th}(\mathcal{M}) = \text{Th}(\mathcal{N})$.

Note that \equiv is an equivalence relation on \mathcal{L} -structures.

May use $\equiv_{\mathcal{L}}$ for emphasis.

Exercise (ESI #2) Let T be an \mathcal{L} -theory. TFAE

i) T is complete.

ii) For any \mathcal{L} -sentence ϕ , if $T \not\models \phi$ then $T \models \neg \phi$.

iii) Any two models of T are elementarily equivalent.

Example 1.7 Let $\mathcal{L} = \emptyset$ and $T = \{\phi_n : n \geq 2\}$ where

$$\phi_n \text{ is } \exists x_1, \dots, \exists x_n \bigwedge_{i \neq j} x_i \neq x_j.$$

T is "the theory of infinite sets"

Then T is complete (ESI #3). So, as \mathcal{L} -structures,

$$\mathbb{N} \equiv \mathbb{Z} \equiv \mathbb{Q} \equiv \mathbb{R} \equiv \mathbb{C} \equiv \mathcal{P}(\mathbb{C}) \equiv \text{any infinite set.}$$

Theorem 1.8 (Vaught's Test)

Let T be an L -theory st

a) T has no finite models,

b) $\exists \kappa \geq |L| + \aleph_0$ st any two models of T of cardinality κ are elementarily equivalent.

Then T is complete.

Proof Suppose T is not complete. Then there is a sentence ϕ st $T \cup \{\neg\phi\}$ is satisfiable and $T \cup \{\phi\}$ is satisfiable. By (a) these theories have infinite models. By LST, these theories have models of size κ . This contradicts (b).