

Diagrams + Extensions

Let \mathcal{M} be an \mathcal{L} -structure.

Remark 5.8 If $h: \mathcal{M} \rightarrow \mathcal{N}$ is an \mathcal{L} -embedding then after identifying $a \in \mathcal{M}$ with $h(a) \in \mathcal{N}$, we can view \mathcal{M} as a substructure of \mathcal{N} . Similarly, if h is an elementary embedding, then \mathcal{M} "is" an elementary substructure of \mathcal{N} .

Given $A \in \mathcal{M}$, let $\mathcal{L}_A = \mathcal{L} \cup \{ \underline{a} : a \in A \}$ where \underline{a} is a new constant symbol.

Then \mathcal{M} is canonically an \mathcal{L}_A -structure where $\underline{a}^{\mathcal{M}} = a$.

Def 5.9 The diagram of \mathcal{M} , written $D(\mathcal{M})$, is the $\mathcal{L}_{\mathcal{M}}$ -theory consisting of all quantifier-free $\mathcal{L}_{\mathcal{M}}$ -sentences ϕ st $\mathcal{M} \models \phi$.

The elementary diagram of \mathcal{M} is $Th_{\mathcal{M}}(\mathcal{M}) := Th_{\mathcal{L}_{\mathcal{M}}}(\mathcal{M})$.

Proposition 5.10 Suppose \mathcal{M} is an \mathcal{L} -structure and \mathcal{N}^* is an $\mathcal{L}_{\mathcal{M}}$ -structure st $\mathcal{N}^* \models D(\mathcal{M})$. Let \mathcal{N} be the reduct of \mathcal{N}^* to \mathcal{L} . Define $h: \mathcal{M} \rightarrow \mathcal{N}$ st $h(a) = \underline{a}^{\mathcal{N}^*}$. Then h is an \mathcal{L} -embedding. Moreover, if $\mathcal{N}^* \models Th_{\mathcal{M}}(\mathcal{M})$ then h is an elementary embedding.

Proof Use Corollary 2.4. Let $\phi(x_1, \dots, x_n)$ be a quantifier-free \mathcal{L} -formula, and fix $a_1, \dots, a_n \in \mathcal{M}$. Then $\mathcal{M} \models \phi(a_1, \dots, a_n)$ iff $\mathcal{M} \models \phi(\underline{a}_1, \dots, \underline{a}_n)$ iff

$$\phi(\underline{a}_1, \dots, \underline{a}_n) \in D(\mathcal{M}) \text{ iff } \mathcal{N}^* \models \phi(\underline{a}_1, \dots, \underline{a}_n) \text{ iff } \mathcal{N} \models \phi(h(a_1), \dots, h(a_n))$$

The "moreover" is similarly. □.

Application to Groups: Recall that an abelian group G is orderable if there is a linear order $<$ on G st $\forall x, y, z \in G$, if $x < y$ then $x+z < y+z$.

Note that any orderable abelian group is torsion-free. E.g. $x > 0 \Rightarrow x < 2x < 3x < \dots$

Theorem 5.11 (Levi 1942)

Any torsion-free abelian group is orderable.

Proof Let $\mathcal{L}^0 = \{+, 0\}$ be the language of (abelian) groups. Set $\mathcal{L} = \mathcal{L}^0 \cup \{<\}$, where $<$ is a binary relation symbol. Let σ be the \mathcal{L} -sentence

$$\forall x \forall y \forall z (x < y \rightarrow x + z < y + z)$$

Now let G be a torsion-free abelian group, viewed as an \mathcal{L}^0 -structure.

Define the \mathcal{L}_G -theory $T = \underbrace{D(G)}_{\mathcal{L}^0\text{-theory}} \cup \{\text{axioms for abelian groups}\} \cup \{\text{axioms for linear order}\} \cup \{\sigma\}$

Suppose T has a model \mathcal{M} . Then $(\mathcal{M}, +^{\mathcal{M}}, 0^{\mathcal{M}}, <^{\mathcal{M}})$ is an ordered abelian group and $G \subseteq (\mathcal{M}, +^{\mathcal{M}}, 0^{\mathcal{M}})$ by Proposition 5.10. So G is a subgroup of an ordered abelian group, hence is orderable. So it suffices to show that T has a model.

Fix $\Sigma \in T$ finite. Let $A = \{a \in G : a \text{ is used in some } \mathcal{L}_G\text{-sentence in } \Sigma\}$. \uparrow in $D(G)$

Let $H = \langle A \rangle \leq G$. Then $H \cong \mathbb{Z}^n$ for some $n \geq 0$ by the structure theorem for finitely generated abelian groups. View H as an \mathcal{L}_A -structure st $a^H = a$ and $<^H$ is the lexicographic ordering. Then $H \models_{\mathcal{L}_A} \Sigma$ and so

$H \models \Phi$ for any $\Phi \in D(G)$ using only extra constants from A (by Corollary 2.4)

So $H \models \Sigma$. □

Quantifier Elimination

Idea: let T be an \mathcal{L} -theory and $\mathcal{M} \models T$. Then $X \subseteq M^n$ is definable if there is an \mathcal{L} -formula $\phi(x_1, \dots, x_n)$ st $X = \{\bar{a} \in M^n : \mathcal{M} \models \phi(\bar{a})\}$.

Goal: Study definable subsets of models of T .

Quantifiers make this difficult. X might be nice but the projection

$$Y = \{(a_1, \dots, a_{n-1}) \in M^{n-1} : (\bar{a}, b) \in X \text{ for some } b \in M\} \text{ (defined by } \exists x_n \phi(\bar{x}))$$

might be complicated.

Def 5.12 An \mathcal{L} -theory T has quantifier elimination if for any \mathcal{L} -formula $\varphi(x_1, \dots, x_n)$ there is a quantifier-free \mathcal{L} -formula $\psi(x_1, \dots, x_n)$ st

$$T \models \forall \bar{x} (\varphi(\bar{x}) \leftrightarrow \psi(\bar{x})).$$

(So φ and ψ define the same set in any $\mathcal{M} \models T$).

Example 5.13

① $T = \text{Th}(F)$ where F is a field. $\varphi(w, x, y, z)$ is " $\begin{pmatrix} w & x \\ y & z \end{pmatrix}$ has an inverse"

i.e. $\exists s \exists t \exists u \exists v \left(\begin{pmatrix} w & x \\ y & z \end{pmatrix} \begin{pmatrix} s & t \\ u & v \end{pmatrix} = \begin{pmatrix} 1 & 0 \\ 0 & 1 \end{pmatrix} \right)$

Then $T \models \forall w \forall x \forall y \forall z (\varphi(w, x, y, z) \leftrightarrow wz - xy \neq 0)$.