Implicitization of surfaces via Geometric Tropicalization

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Three references:

Sturmfels, Tevelev, Yu: The Newton polytope of the implicit equation (2007)
Sturmfels, Tevelev: Elimination theory for tropical varieties (2008)
MAC: arXiv:1105.0509 (2011)
(and many, many more!)

Implicitization problem: Classical vs. tropical approach

Input: Laurent polynomials $f_1, f_2, \ldots, f_n \in \mathbb{C}[t_1^{\pm 1}, \ldots, t_d^{\pm 1}]$.

Algebraic Output: The *prime* ideal I defining the Zariski closure Y of the image of the map:

$$\mathbf{f} = (f_1, \ldots, f_n) \colon \mathbb{T}^d \dashrightarrow \mathbb{T}^n$$

The ideal I consists of all polynomial relations among $f_1, f_2, \ldots, f_n$.

Existing methods: Gröbner bases and resultants.

- GB: always applicable, but often too slow.
- Resultants: useful when $n = d + 1$ and I is *principal*, with limited use.

Geometric Output: Invariants of Y, such as dimension, degree, etc.

Punchline: We can *effectively* compute them using tropical geometry.

TODAY: Study the case when $\mathbf{d = 2}$ and $\mathbf{Y}$ is a **surface**.

Example: parametric surface in $\mathbb{T}^3$

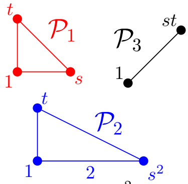
`Input`: Three Laurent polynomials in two unknowns:

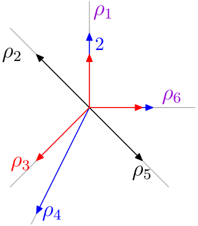
$$\begin{cases} x = f_1(s,t) = 3 + 5\,s + 7\,t, \\ y = f_2(s,t) = 17 + 13\,t + 11\,s^2, \\ z = f_3(s,t) = 19 + 47\,st. \end{cases}$$

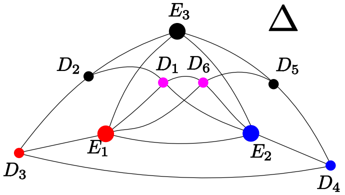
`Output`: The Newton polytope of the implicit equation $g(x, y, z)$.

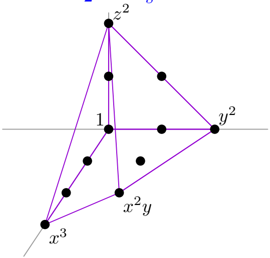
STRATEGY: Recover the Newton polytope of $g(x, y, z)$ from the **Newton polytopes** of the `input polynomials` f_1, f_2, f_3.

$$Y = \begin{cases} x = f_1(s,t) = 3 + 5\,s + 7\,t, \\ y = f_2(s,t) = 17 + 13\,t + 11\,s^2, \\ z = f_3(s,t) = 19 + 47\,st. \end{cases} \quad \rightsquigarrow \text{ Newton polytope of } g(x,y,z).$$

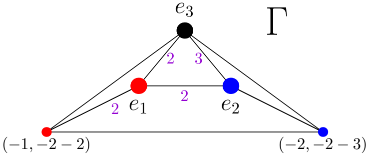








f-vector= (5, 8, 5)



- Γ is a balanced weighted *planar* graph in $\mathbb{R}^3$. It is the tropical variety $\mathcal{T}(g(x,y,z))$, dual to the Newton polytope of g.
- We can recover $g(x,y,z)$ from Γ using *numerical linear algebra*.

What is Tropical Geometry?

Given a variety $X \subset \mathbb{T}^n$ with defining ideal $I \subset \mathbb{C}[x_1^{\pm 1}, \ldots, x_n^{\pm 1}]$, the tropicalization of X equals:

$$\mathcal{T}X = \mathcal{T}I := \{w \in \mathbb{R}^n \,|\, \mathrm{in}_w I \text{ contains no monomial}\}.$$

1. It is a rational polyhedral fan in $\mathbb{R}^n \rightsquigarrow \mathcal{T}X \cap \mathbb{S}^{n-1}$ is a spherical polyhedral complex.
2. If I is prime, then $\mathcal{T}X$ is pure of the same dimension as X.
3. Maximal cones have canonical multiplicities attached to them.

Example (hypersurfaces):

- Maximal cones in $\mathcal{T}(g)$ are dual to edges in the Newton polytope $\mathrm{NP}(g)$, and m_σ is the lattice length of the associated edge.
- Multiplicities are essential to recover $\mathrm{NP}(g)$ from $\mathcal{T}(g)$.

What is Geometric Tropicalization?

AIM: Given $Z \subset \mathbb{T}^N$ a **surface**, compute $\mathcal{T}Z$ from the *geometry* of Z.
KEY FACT: $\mathcal{T}Z$ can be characterized in terms of divisorial valuations.

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Theorem (Geometric Tropicalization [Hacking - Keel - Tevelev, *C*.])

Consider $\mathbb{T}^N$ with coordinate functions $\chi_1, \ldots, \chi_N$, and let $Z \subset \mathbb{T}^N$ be a closed smooth **surface**. *Suppose $\overline{Z} \supset Z$ is any normal and $\mathbb{Q}$-factorial compactification, whose boundary divisor has m irreducible components $D_1, \ldots, D_m$ with no triple intersections (***C.N.C.***). Let Δ be the graph:*

$$V(\Delta) = \{1, \ldots, m\} \quad ; \quad (i,j) \in E(\Delta) \iff D_i \cap D_j \neq \emptyset.$$

Realize Δ *as a graph* $\Gamma \subset \mathbb{R}^N$ *by* $[D_k]{:=}(val_{D_k}(\chi_1), \ldots, val_{D_k}(\chi_N)) \in \mathbb{Z}^N$.

Then, $\mathcal{T}Z$ is the **cone over the graph** Γ.

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Theorem (Combinatorial formula for multiplicities [C.])

$$m_{([D_i],[D_j])} = (D_i \cdot D_j)\left[\,(\mathbb{Z}\langle[D_i],[D_j]\rangle)^{sat} : \mathbb{Z}\langle[D_i],[D_j]\rangle\,\right]$$

QUESTION: How to compute $\mathcal{T}Y$ from a parameterization

$$\mathbf{f} = (f_1, \ldots, f_n) \colon \mathbb{T}^2 \dashrightarrow Y \subset \mathbb{T}^n \quad ?$$

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ANSWER: Compactify the domain $X = \mathbb{T}^2 \smallsetminus \bigcup_{i=1}^{n} (f_i = 0)$ and use the map $\mathbf{f}$ to translate back to Y.

Proposition

Given $\mathbf{f} \colon X \subset \mathbb{T}^2 \to Y \subset \mathbb{T}^n$ generically finite map of degree δ, let $\overline{X}$ be a normal, $\mathbb{Q}$-factorial, CNC compactification with intersection complex Δ. Map each vertex D_k of Δ in $\mathbb{Z}^n$ to a vertex $[\widetilde{D_k}]$ of $\Gamma \subset \mathbb{R}^n$, where

$$[\widetilde{D_k}] = val_{D_k}(\chi \circ \mathbf{f}) = \mathbf{f}^{\#}([D_k]).$$

Then, $\mathcal{T}Y$ is the cone over the graph $\Gamma \subset \mathbb{R}^n$, with multiplicities

$$m_{([\widetilde{D_i}],[\widetilde{D_j}])} = \frac{1}{\delta} (D_i \cdot D_j) \left[(\mathbb{Z}\langle[\widetilde{D_i}], [\widetilde{D_j}]\rangle)^{sat} : \mathbb{Z}\langle[\widetilde{D_i}], [\widetilde{D_j}]\rangle \right].$$

Implicitization of *generic* surfaces

SETTING: Let $f = (f_1, \dots, f_n)\colon \mathbb{T}^2 \dashrightarrow Y \subset \mathbb{T}^n$ of $\deg(f) = \delta$, where we **fix the Newton polytope** of each f_i and allow **generic coefficients**.
GOAL: Compute the graph Γ of $\mathcal{T}Y$ from the Newton polytopes $\{\mathcal{P}_i\}_{i=1}^n$.

IDEA: Compactify X inside the proj. toric variety $X_{\mathscr{N}}$, where $\mathscr{N}$ is the common refinement of all $\mathscr{N}(P_i)$. *Generically*, $\overline{X}$ is smooth with CNC.

The vertices and edges of the boundary intersection complex Δ are

$$V(\Delta) = \{E_i : \dim \mathcal{P}_i \neq 0, 1 \leq i \leq n\} \bigcup \{D_\rho : \rho \in \mathscr{N}^{[1]}\},$$

- $(D_\rho, D_{\rho'}) \in E(\Delta)$ iff ρ, ρ' are *consecutive* rays in $\mathscr{N}$.
- $(E_i, D_\rho) \in E(\Delta)$ iff $\rho \in \mathscr{N}(\mathcal{P}_i)$.
- $(E_i, E_j) \in E(\Delta)$ iff $(f_i = f_j = 0)$ has a solution in $\mathbb{T}^2$.

Then, Γ is the realization of Δ via

$$[E_i] := e_i \quad (1 \leq i \leq n), \quad [D_\rho] := \operatorname{trop}(\mathbf{f})(\eta_\rho) \quad \forall \text{ ray } \rho \ (\eta_\rho \text{ prim. vector.})$$

Theorem [Sturmfels-Tevelev-Yu, C.]: $\mathcal{T}Y$ is the weighted cone over Γ.

Implicitization of *non-generic* surfaces

Non-genericity $\leftrightarrow$ CNC condition is violated.

Solution 1:

1. Embed X in $X_{\mathscr{N}}$.
2. Resolve triple intersections and singularities by classical blow-ups, and carry divisorial valuations along the way.

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Solution 2:

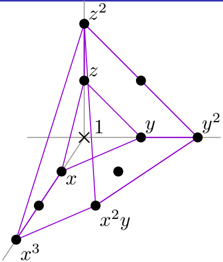
1. Embed X in $\mathbb{P}^2_{(s,t,u)} \rightsquigarrow n+1$ boundary divisors
$$E_i = (f_i = 0) \quad (1 \leq i \leq n), \quad E_\infty = (u = 0).$$
2. Resolve triple intersections and singularities by blow-ups $\pi\colon \tilde{X} \to X$, and read divisorial valuations by *columns*

$$(f \circ \pi)^*(\chi_i) = \pi^*(E_i - \deg(f_i)E_\infty) = E_i' - \deg(f_i)E_\infty' - \sum_{j=1}^{r} b_{ij}H_j \quad \forall i.$$

The graph Δ is obtained by gluing resolution diagrams and adding pairwise intersections.

Example (non-generic surface)

$$Y = \begin{cases} x = f_1(s,t) = s - t, \\ y = f_2(s,t) = t - s^2, \\ z = f_3(s,t) = -1 + s\,t, \end{cases}$$



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Affine Charts: