

Conclude $(1+2+\dots+(n-1)+n) + (n+(n-1)+\dots+1) = (n+1)n$

So $2(1+2+\dots+n) = 2 \sum_{k=1}^n k = (n+1)n$

$$\sum_{k=1}^n k = \frac{(n+1)n}{2}$$

• This is a very elegant argument, but it won't help us with finding $\sum_{k=1}^n k^2$, $\sum_{k=1}^n k^3$, etc. For this, we give a second way of finding the formula for $\sum_{k=1}^n k$.

• 2nd approach Use telescopic sums

Recall: $(k+1)^2 = k^2 + 2k + 1$, so $(k+1)^2 - k^2 = 2k + 1$ (*)

Q: What happens if we sum (*) as we vary k between 1 and n?

A: A lot of cancellations occur!

$$\sum_{k=1}^n ((k+1)^2 - k^2) = (\cancel{2^2} - 1^2) + (\cancel{3^2} - \cancel{2^2}) + \dots + (\cancel{n^2} - \cancel{(n-1)^2}) + ((n+1)^2 - \cancel{n^2})$$

← only survivors! →

$$= (n+1)^2 - 1^2$$

In general, this will lead to a telescopic sum:

$$(a_2 - a_1) + (a_3 - a_2) + \dots + (a_{n+1} - a_n) = \sum_{k=1}^n (a_{k+1} - a_k) = \overset{\text{from last term}}{a_{n+1}} - \underset{\text{from first term}}{a_1}$$

Now, we can sum this via the right-hand side of (*):

$$\sum_{k=1}^n (2k+1) = \begin{array}{|l} 2 \cdot 1 \\ 2 \cdot 2 \\ 2 \cdot 3 \\ \vdots \\ 2 \cdot n \end{array} + \begin{array}{|l} 1 \\ 1 \\ 1 \\ \vdots \\ 1 \end{array}$$

Again, we "add column-by-column"

① = $2 \cdot 1 + 2 \cdot 2 + \dots + 2 \cdot n = 2(1+2+\dots+n)$
= $2 \sum_{k=1}^n k$

② = $\underbrace{1+1+\dots+1}_{n \text{ times}} = n$

So we get $\sum_{k=1}^n (2k+1) = 2 \sum_{k=1}^n k + n$

by (*) ||

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$$\sum_{k=1}^n ((k+1)^2 - k^2) = (n+1)^2 - 1 = n^2 + 2n + 1 - 1 = n(n+2)$$

Conclude: $2 \left(\sum_{k=1}^n k \right) + n = n(n+2)$

$$2 \left(\sum_{k=1}^n k \right) = n(n+2) - n = n(n+1)$$

$$\sum_{k=1}^n k = \frac{n(n+1)}{2}$$

This approach can be generalized to compute $\sum_{k=1}^n k^2$, $\sum_{k=1}^n k^3$ (HWS)

Proposition 2: $\sum_{k=1}^n k^2 = \frac{n(n+1)(2n+1)}{6}$

Why? We use a different telescopic sum.

$$(k+1)^3 - k^3 = k^3 + 3k^2 + 3k + 1 - k^3 = 3k^2 + 3k + 1$$

(**)

So $\sum_{k=1}^n ((k+1)^3 - k^3) = (\cancel{2^3} - 1^3) + (\cancel{3^3} - \cancel{2^3}) + \dots + ((n+1)^3 - \cancel{n^3}) = (n+1)^3 - 1 = n^3 + 3n^2 + 3n$

→ only survivors ←
↓ telescopic!

On the other hand $\sum_{k=1}^n (3k^2 + 3k + 1) =$

$3 \cdot 1^2$	+	$3 \cdot 1$	+	1
$3 \cdot 2^2$	+	$3 \cdot 2$	+	1
\vdots				
$3 \cdot n^2$	+	$3 \cdot n$	+	1

Add columns one at a time

① + ② + ③

- where
- ① = $3 \cdot 1^2 + 3 \cdot 2^2 + \dots + 3 \cdot n^2 = 3(1^2 + 2^2 + \dots + n^2) = 3 \sum_{k=1}^n k^2$
 - ② = $3 \cdot 1 + 3 \cdot 2 + \dots + 3 \cdot n = 3(1 + 2 + \dots + n) = 3 \sum_{k=1}^n k$ WHAT WE WANT!
 - ③ = $\underbrace{1 + 1 + \dots + 1}_{n \text{ times}} = n$ $= 3 \frac{(n+1)n}{2}$ (by Prop)

So we get $n^3 + 3n^2 + 3n = \sum_{k=1}^n ((k+1)^3 - k^3) = \sum_{k=1}^n (3k^2 + 3k + 1) =$
 $= \textcircled{1} + \textcircled{2} + \textcircled{3} = 3 \sum_{k=1}^n k^2 + 3 \frac{(n+1)n}{2} + n$
by ()*
WHAT WE WANT!

So $3 \sum_{k=1}^n k^2 = n^3 + 3n^2 + 3n - 3 \frac{(n+1)n}{2} - n$
 $= n^3 + 3n^2 + 3n - \frac{3n^2}{2} - \frac{3n}{2} - n$
 $= \frac{1}{2} (2n^3 + 3n^2 + n) = \frac{1}{2} (n (2n^2 + 3n + 1))$
 $= \frac{1}{2} n (2n+1)(n+1)$

Conclude $\sum_{k=1}^n k^2 = \frac{n(n+1)(2n+1)}{6}$

Proposition 3: $\sum_{k=1}^n k^3 = \frac{n^2(n+1)^2}{4}$ (see HW5)

• General procedure to compute $\sum_{k=1}^n k^s$ for $s=1, 2, 3, \dots$

STEP 1 Use the Binomial Theorem to write $(k+1)^{s+1} - k^{s+1}$.

$(k+1)^{s+1} - k^{s+1} = \binom{s+1}{1} k^s + \binom{s+1}{2} k^{s-1} + \dots + \binom{s+1}{s} k + 1$

STEP 2: Add up both sides:

• Left-hand side: telescopic sum:

$\sum_{k=1}^n (k+1)^{s+1} - k^{s+1} = (\cancel{2^{s+1}} - \cancel{1^{s+1}}) + (\cancel{3^{s+1}} - \cancel{2^{s+1}}) + \dots + ((n+1)^{s+1} - \cancel{n^{s+1}})$
only 2 survivors
 $= (n+1)^{s+1} - 1^{s+1}$

• Right-hand side:

$\sum_{k=1}^n \left(\binom{s+1}{1} k^s + \binom{s+1}{2} k^{s-1} + \dots + \binom{s+1}{s} k + 1 \right) =$ arrange in columns & add them up

$= \binom{s+1}{1} \sum_{k=1}^n k^s + \binom{s+1}{2} \sum_{k=1}^n k^{s-1} + \dots + \binom{s+1}{s} \sum_{k=1}^n k + n \cdot 1$

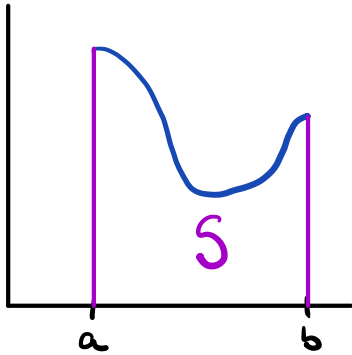
WHAT WE WANT!

formulas for lower exponents are assumed to be known!

So we can compute $\sum_{k=1}^n k^s$ using this approach!

§ 2. The area under a curve

GOAL: Given a continuous function $f: [a, b] \rightarrow \mathbb{R}_{\geq 0}^{\{x \geq 0\}}$, we want to find the area of the region S that lies under the curve $y = f(x)$ & above the x -axis



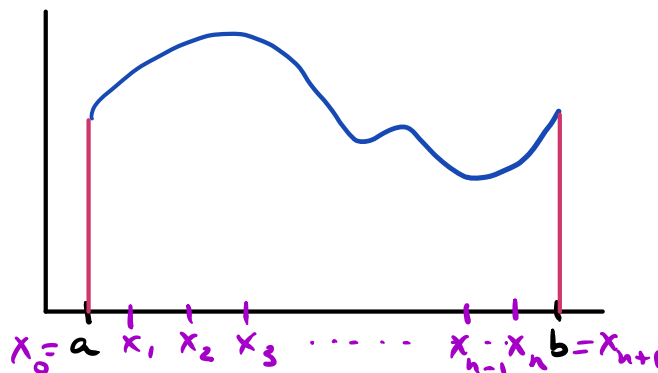
We use 2 approximations to compute $\text{Area}(S)$

- ① Overestimate with larger rectangles (intrap S)
 \rightsquigarrow Upper Riemann Sums
- ② Underestimate with small rectangles (exhaust S)
 \rightsquigarrow Lower Riemann Sums

Both types of rectangles will have:

- base on the x -axis of small length
- height: either a max (for ①) or a min (for ②) of f restricted to each base

STEP 1 Pick points x_1, \dots, x_n that determine the bases of these \square



We write the differences as

$$\begin{aligned} \Delta x_1 &= x_1 - a \\ \Delta x_2 &= x_2 - x_1 \\ &\vdots \\ \Delta x_n &= x_n - x_{n-1} \\ \Delta x_{n+1} &= b - x_n \end{aligned}$$

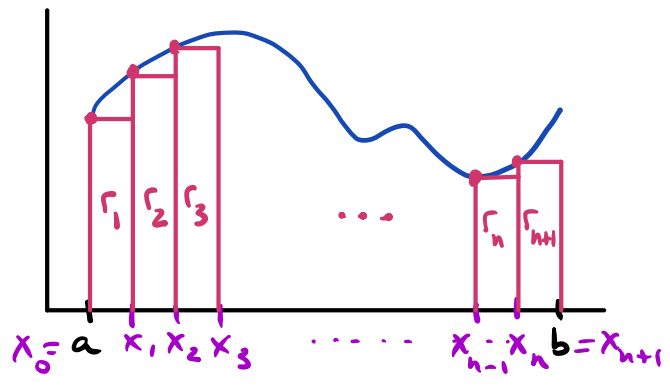
we assume they are all small (which means n is going to be larger!)

To fix ideas, we can make $\Delta x_k = \frac{b-a}{n+1}$ for all $k = 1, 2, \dots, n+1$.

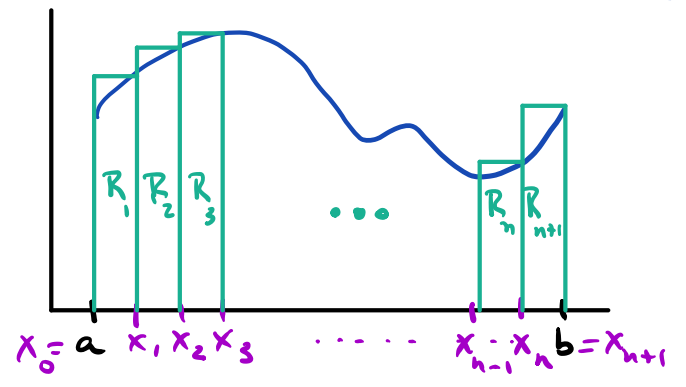
STEP 2: Use these intervals as the bases of our rectangles.

$n+1$ small rectangles = $\Gamma_1, \dots, \Gamma_{n+1}$ & height of $\Gamma_k = \min$ of f on $[x_{k-1}, x_k]$

$n+1$ large rectangles = R_1, \dots, R_{n+1} & height of $R_k = \max$ of f in $[x_{k-1}, x_k]$



LOWER SUMS



UPPER SUMS

If f is continuous $m_k = \text{height of } r_k = f(\underline{x}_k^*)$ for some \underline{x}_k^* in $[x_{k-1}, x_k]$
 $M_k = \text{height of } R_k = f(\bar{x}_k^*)$ for some \bar{x}_k^* in $[x_{k-1}, x_k]$

The existence of \underline{x}_k^* & \bar{x}_k^* is guaranteed by the Extreme Values Thm.

STEP 3: We compare S with the areas covered by small & larger rectangles:

$$\boxed{A(f)} := \sum_{k=1}^{n+1} \text{Area}(r_k) \leq \text{Area}(S) \leq \sum_{k=1}^{n+1} \text{Area}(R_k) := \boxed{A(R)}$$

As we increase n & we decrease all Δx_k (bases of the rectangles r_k & R_k), the heights $f(\underline{x}_k^*)$ & $f(\bar{x}_k^*)$ get closer to each other. This gives:

Theorem: If f is continuous, $\sum_{k=1}^n \text{Area}(r_k)$ & $\sum_{k=1}^{n+1} \text{Area}(R_k)$ have the same limit as $n \rightarrow \infty$. This is the value of $\text{Area}(S)$

Why? Since f is continuous, the EVT gave us $f(\bar{x}_k^*) = M_k$
 $f(\underline{x}_k^*) = m_k$.

If we pick any other point x_k^* in $[x_{k-1}, x_k]$, we'll have

$$\text{Area}(r_k) = f(\underline{x}_k^*) \Delta x_k \leq f(x_k^*) \Delta x_k \leq f(\bar{x}_k^*) \Delta x_k = \text{Area}(R_k)$$

Call $\sum_{k=1}^n \underbrace{f(x_k^*) \Delta x_k}_{\text{area of rectangle with base } = [x_{k-1}, x_k] \text{ height } = f(x_k^*)} = \text{Riemann Sum}$

area of rectangle with base = $[x_{k-1}, x_k]$
height = $f(x_k^*)$

We get $A(r) \leq \text{Riemann Sum} \leq A(R)$

"
Lower Riemann Sum

"
Upper Riemann Sum

If $\delta = \max_{k=1, \dots, n} \Delta x_k \rightarrow 0$, then all $\Delta x_k \rightarrow 0$ & this

forces $n \rightarrow \infty$.

• If we show $\lim_{\delta \rightarrow 0} A(r) = \lim_{\delta \rightarrow 0} A(R)$ (*), then by the Squeeze

Theorem we'll have $\text{Area}(S) = \text{this limit} = \lim_{\delta \rightarrow 0} \sum_{k=1}^n f(x_k^*) \Delta x_k$
 $\delta = \max_{k=1, \dots, n} \Delta x_k$ & x_k^* any point in $[x_{k-1}, x_k]$

(Special choices: $x_k^* = \text{leftmost point} = x_{k-1}$)

$x_k^* = \text{rightmost point} = x_k$

$x_k^* = \text{midpoint} = \frac{x_{k-1} + x_k}{2}$

Definition: We set $\int_a^b f(x) dx$ as the area of S , viewed as a limit of Riemann Sums. We call it the definite integral.

Names: $a = \text{lower limit of integration}$

, $f(x) = \text{integrand}$

$b = \text{upper}$ _____

$x = \text{variable of integration}$.

Q: How can we guarantee (*)? Continuity of f will show that

$m_k = f(\underline{x}_k^*)$ & $M_k = f(\bar{x}_k^*)$ will be close enough if Δx_k is sufficiently small (See Appendix A5 for details).

Next Time: Examples of integrable & non-integrable functions.