

Lecture 3: §1.2 Echelon forms & Gauss Jordan Elimination

Last Time:

System of m Linear Eqs in n variables

$\rightsquigarrow$

$m \times (n+1)$ Matrix

$$\begin{cases} \text{Eqn 1: } a_{11}x_1 + a_{12}x_2 + \dots + a_{1n}x_n = b_1 \\ \text{Eqn 2: } a_{21}x_1 + a_{22}x_2 + \dots + a_{2n}x_n = b_2 \\ \vdots \\ \text{Eqn } m: a_{m1}x_1 + a_{m2}x_2 + \dots + a_{mn}x_n = b_m \end{cases} \rightsquigarrow B = \left(\begin{array}{ccc|c} a_{11} & a_{12} & \dots & a_{1n} & b_1 \\ a_{21} & a_{22} & \dots & a_{2n} & b_2 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ a_{m1} & a_{m2} & \dots & a_{mn} & b_m \end{array} \right)$$

$[x_1] \ [x_2] \ \dots [x_n] \ \text{const}$

3 Elementary Operations to simplify systems / put matrices in Echelon form

EQUATIONS

- ① Interchange 2 Equations $[E_i \leftrightarrow E_j] \ i \neq j$
- ② Replace an Equation by a non-zero constant multiple α $[E_i \rightarrow \alpha E_i]$
- ③ Replace an equation by adding to it a constant multiple of a different equation $[E_i \rightarrow E_i + \alpha E_j] \ i \neq j$

MATRICES

- ① Interchange 2 Rows $[R_i \leftrightarrow R_j] \ i \neq j$
- ② Replace a Row by a non-zero constant multiple α $[R_i \rightarrow \alpha R_i]$
- ③ Replace a Row by adding to it a constant multiple of a different Row $[R_i \rightarrow R_i + \alpha R_j] \ i \neq j$

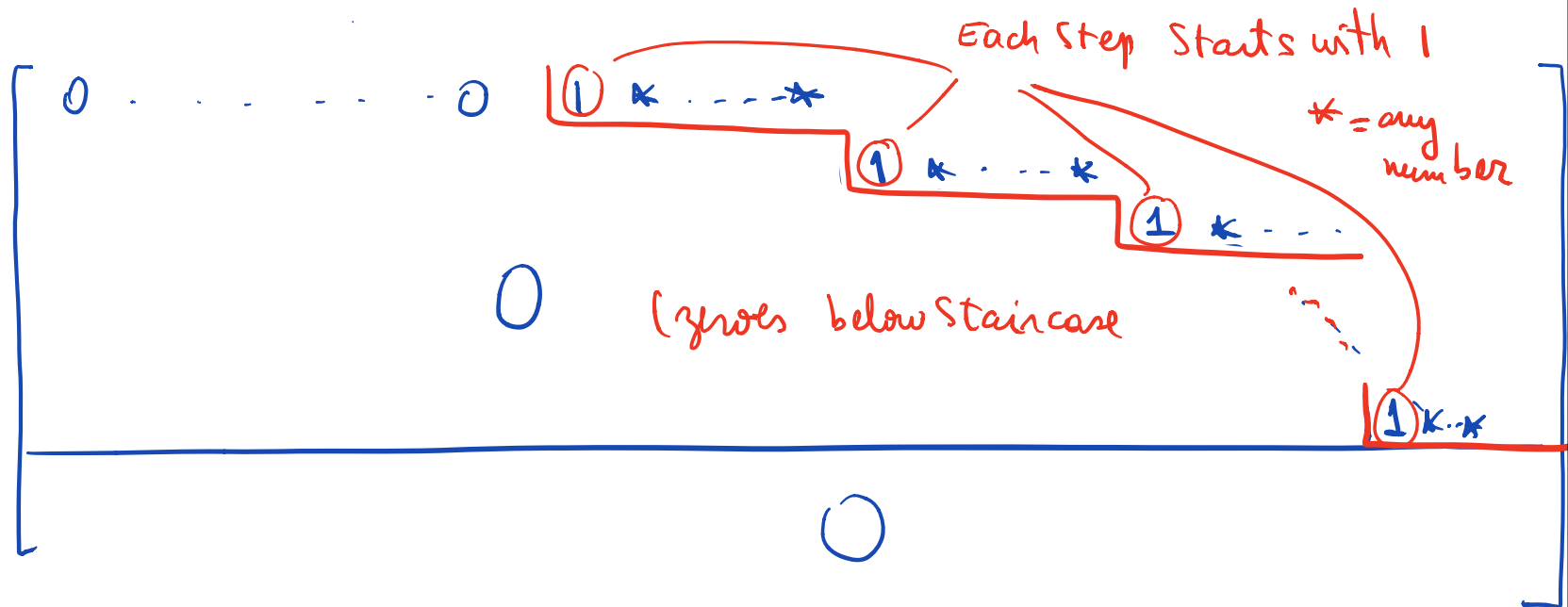
Main result: Elementary operations preserve solution sets (Equiv Systems, Row equiv Matrices)

Favorite Matrices = (Reduced) Echelon form

Def. An $m \times n$ matrix A is in echelon form (E.F.) if

- (1) all rows containing only 0's are grouped together at the bottom of A
- (2) in every nonzero row, the first nonzero entry (from the left) is a 1.
- (3) if a row is nonzero, the first nonzero entry is to the right of the first nonzero entry of the previous row ("staircase" shape)

In general :



Def An $m \times n$ matrix A is in reduced echelon form (REF) if

- (1) A is in EF (staircase shape + start with 1 each nonzero row)
- (2) The first nonzero entry in a row is the ONLY nonzero entry on the corresponding column

Q: Why do we care about REF matrices?

A: They give the simplest systems to solve.

Algorithm:

System $\rightsquigarrow B = [A | b]$

$\updownarrow$ Same Solutions!

$\downarrow$ Row operations

New System
easy to solve

$\rightsquigarrow B' = [A' | b']$ with A' in REF

- dependent vars
- indep vars

Example 1:

$$\begin{cases} x_2 + x_3 - x_4 = 3 \\ x_1 + 2x_2 - x_3 + x_4 = 1 \\ -x_1 + x_2 + 7x_3 - x_4 = 0 \end{cases}$$

$$\leadsto B = \left[\begin{array}{cccc|c} 0 & 1 & 1 & -1 & 3 \\ 1 & 2 & -1 & 1 & 1 \\ -1 & 1 & 7 & -1 & 0 \end{array} \right]$$

Obs: We can check these are solutions of the original system

How? Substitute the values for x_1, x_2, x_3 in the equations!

$$(x_1, x_2, x_3, x_4) = (-13 - 6x_4, \frac{17}{3} + 2x_4, -\frac{8}{3} - x_4, x_4)$$

means
$$\begin{cases} x_1 = -13 - 6x_4 \\ x_2 = \frac{17}{3} + 2x_4 \\ x_3 = -\frac{8}{3} - x_4 \end{cases}$$

$\leadsto$ substitute this in 3 equations & check that we get the expected value

$$\begin{cases} x_2 + x_3 - x_4 = 3 \\ x_1 + 2x_2 - x_3 + x_4 = 1 \\ -x_1 + x_2 + 7x_3 - x_4 = 0 \end{cases}$$

This check does not help us verify we found all the solutions, but we can catch some arithmetic errors this way.

Example 2: $B = \left[\begin{array}{ccc|c} 1 & 0 & 8 & 5 \\ 0 & 1 & 4 & 7 \\ 0 & 0 & 0 & 1 \end{array} \right] = B'$

Example 3 $B = \left[\begin{array}{ccc|c} 1 & 0 & 0 & 5 \\ 0 & 1 & 4 & 7 \\ 0 & 0 & 0 & 0 \end{array} \right] = B'$

Theorem: Given a matrix $B = [A|b]$ of size $m \times (n+1)$, there is a unique matrix $B' = [A'|b']$ of size $m \times (n+1)$ that is row equivalent to B with A' in REF. We find B' using Gauss-Jordan Elimination Algorithm

Obs: We can get A' in many ways, but end result is the same.

Example 4: Solve

$$\begin{cases} x_2 - x_3 + x_4 - x_5 = 1 \\ x_1 - 3x_2 + x_3 - x_4 + x_5 = 3 \\ -2x_2 + 2x_3 + x_4 - x_5 = 2 \\ x_2 - x_3 + 7x_4 - 7x_5 = 9 \end{cases}$$

$$\leadsto B = \left[\begin{array}{ccccc|c} 0 & 1 & -1 & 1 & -1 & 1 \\ 1 & -3 & 1 & -1 & 1 & 3 \\ 0 & -2 & 2 & 1 & -1 & 2 \\ 0 & 1 & -1 & 7 & -7 & 9 \end{array} \right]$$



Example 5: Solve
$$\begin{cases} 2x_1 + 3x_2 - 4x_3 = 3 \\ x_1 - 2x_2 - 2x_3 = -2 \\ -x_1 + 16x_2 + 2x_3 = \frac{33}{2} \end{cases} \rightsquigarrow B = \left[\begin{array}{ccc|c} 2 & 3 & -4 & 3 \\ 1 & -2 & -2 & -2 \\ -1 & 16 & 2 & \frac{33}{2} \end{array} \right]$$

Obs: Can solve many systems with the same coefficient matrix.

$\left(A \mid \underset{\text{system 1}}{b_1} \mid \underset{\text{system 2}}{b_2} \mid \dots \right)$ since operations are always dictated by A

GAUSS-JORDAN ELIMINATION

Input: $B = [A|b]$
 $m \text{ rows}, n+1 \text{ cols}$

Output: $B' = [A'|b']$ A REF
 $B \sim_{\text{row}} B'$

STEP 0 If B has all zero entries, then $B' = B$

STEP 1 Pick the first (left-most) column with a nonzero entry (call it j)
 [Find Col]

STEP 2 Exchange Rows so that j^{th} column has $a_{1j} \neq 0$ (nonzero entry on row 1)
 [Swap Step]

STEP 3 If $\alpha = a_{1j}$ is first nonzero entry on row 1, do $R_1 \rightarrow \frac{1}{\alpha} R_1$
 [Rescale Step] (After this, new value of $a_{1j} = 1$)

STEP 4 Replace each row R_i with $i = 2, 3, \dots, m$ with $R_i \rightarrow R_i - a_{ij} R_1$
 [Replacement Step]

$\leadsto$ get

$$\begin{bmatrix} 0 & \dots & 0 & 1 & * & \dots & * \\ 0 & \dots & 0 & a_{2j} & * & \dots & * \\ \vdots & & & \vdots & & & \\ 0 & \dots & 0 & a_{mj} & * & \dots & * \end{bmatrix} \xrightarrow{\quad} \underbrace{\begin{bmatrix} 0 & \dots & 0 & 1 & * & \dots & * \\ 0 & \dots & 0 & 0 & * & \dots & * \\ \vdots & & & \vdots & & & \\ 0 & \dots & 0 & 0 & * & \dots & * \end{bmatrix}}_{B_1}$$

Fix

$$B \xrightarrow[\text{STEPS 0-4}]{\sim} \left[\begin{array}{cccc|c|ccc} & & & & i & & & \\ 0 & \dots & 0 & 1 & \times & \dots & \times \\ 0 & \dots & 0 & 0 & \times & & \times \\ & & & \vdots & & & \\ 0 & \dots & 0 & \times & & & \times \end{array} \right] \quad B_1 = m \text{ rows \& } n+1 \text{ cols}$$

B_1

STEP 5: Repeat Steps 0 through 4 for the smaller matrix B_1 until we get an EF matrix.

STEP 6: [From EF to REF]

Work our way backwards to put 0's on top of each 1 starting a row.
(fix columns by moving from right to left)

$$S \left[\begin{array}{cccc|c|ccc} & & & & t & & & \\ & & & & \times & & & \\ & & & & \vdots & & & \\ & & & & \times & & & \\ & & & & 1 & & & \\ & & & & 0 & & & \\ & & & & \vdots & & & \\ & & & & 0 & & & \end{array} \right] \quad \left. \begin{array}{c} \times \\ \vdots \\ \times \end{array} \right\} \text{fix these!}$$

$R_1 \rightarrow R_1 - a_{1,t} R_5$
 $R_2 \rightarrow R_2 - a_{2,t} R_5$
 $R_{s-1} \rightarrow R_{s-1} - a_{s-1,t} R_5$

↑ fix this column first

$$\left[\begin{array}{cccc|c|ccc} & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & & \\ & & & & & & 1 & \\ & & & & & & 0 & \\ & & & & & & \vdots & \\ & & & & & & 0 & \end{array} \right] \quad \left. \begin{array}{c} 0 \\ \vdots \\ 0 \end{array} \right\} \text{fix these!}$$

← MOVE TO THE LEFT.