

Lecture XIX: Recurrence relations & ordinary generating functions (o.g.f.)

Recall: 1. $F_0 = 1, F_1 = 1, F_n = F_{n-1} + F_{n-2} \quad \forall n$ $F_n = \# \{ a_1 + \dots + a_k = n \quad a_i = 1 \text{ or } 2 \}$
(Recursion) Fibonacci Numbers.

2. Showed
$$\sum_{n \geq 0} F_n x^n = \frac{1}{1-x-x^2}$$

We can use partial fractions to give a closed formula for F_n (Binet's Formula from Lecture 1)

3. Factor $1-x-x^2 = (1-\phi x)(1-\bar{\phi} x) \quad \phi = \frac{1+\sqrt{5}}{2}, \bar{\phi} = 1-\phi = \frac{1-\sqrt{5}}{2}$

4.
$$\sum_{n \geq 0} F_n x^n = \frac{1}{1-x-x^2} = \frac{\alpha}{1-\phi x} + \frac{\beta}{1-\bar{\phi} x} = \sum_{n \geq 0} (\alpha \phi^n + \beta \bar{\phi}^n) x^n$$

series

Conclude $F_n = \alpha \phi^n + \beta \bar{\phi}^n \quad \forall n$

Use $F_0 = 1$ & $F_1 = 1$ to solve for α & $\beta \Rightarrow \alpha = \frac{\phi}{\sqrt{5}}, \beta = -\frac{\bar{\phi}}{\sqrt{5}}$

Binet's Formula:
$$F_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^{n+1} - \left(\frac{1-\sqrt{5}}{2}\right)^{n+1}}{\sqrt{5}} \quad \text{for all } n.$$

Note: denominator encoded reflected recurrence $[F_n - F_{n-1} - F_{n-2} \Leftrightarrow 1-x-x^2]$

Claim: We can use the same method to find closed formulas for any sequence admitting a linear recurrence relation over \mathbb{C} (or any alg. closed field of char 0, eg. $\mathbb{C}[[t]]$)

Def: Let $c(z) = 1 + c_1 z + c_2 z^2 + \dots + c_d z^d \in \mathbb{C}[[z]]$ of degree $= d \geq 1$ ($c_d \neq 0$)

We define the reflected polynomial $c^R(z) = z^d + c_1 z^{d-1} + c_2 z^{d-2} + \dots + c_d, c_d \neq 0$

Obs: $c(z) = z^d c^R\left(\frac{1}{z}\right)$

In particular, if $c^R(z) = (z-\alpha_1) \dots (z-\alpha_d)$ (we allow $\alpha_i = \alpha_j$)

$$\Rightarrow c(z) = z^d \left(\frac{1}{z} - \alpha_1\right) \dots \left(\frac{1}{z} - \alpha_d\right) = (1-\alpha_1 z) \dots (1-\alpha_d z)$$

In the example above $c(z) = 1-z-z^2$ & $c^R(z) = z^2 - z - 1 = (z-\phi)(z-\bar{\phi}) = (1-\phi z)(1-\bar{\phi} z)$

The roots of c^R become the powers in the relevant summands for F_n . This will be true in general.

The following is the main theorem in this subject:

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Theorem: Fix a sequence c_1, \dots, c_d of d complex numbers with $c_d \neq 0$ & set $c(z) = 1 + c_1 z + \dots + c_d z^d = \prod_{i=1}^k (1 - \alpha_i z)^{d_i}$ where $\alpha_1, \dots, \alpha_k$ are the roots of $c^R(z)$ (with $\text{mult}(\alpha_i, c^R) = d_i$).

Consider a counting function $F: \mathbb{N}_0 \rightarrow \mathbb{C}$ & write $F_n := F(n) \forall n$.

The following conditions are equivalent:

(A1) [Recurrence of order d] $F_{n+d} + c_1 F_{n+d-1} + \dots + c_d F_n = 0 \quad \forall n \geq 0$

(A2) [Generating function] $F(z) = \sum_{n \geq 0} F_n z^n = \frac{P(z)}{c(z)}$

where $P \in \mathbb{C}[z]$ has $\deg(P) < d$.

(A3) [Partial fractions] $F(z) = \sum_{n \geq 0} F_n z^n = \sum_{i=1}^k \frac{g_i(z)}{(1 - \alpha_i z)^{d_i}}$

where $g_1, \dots, g_k \in \mathbb{C}[z]$ has $\deg(g_i) < d_i \quad \forall i = 1, \dots, k$.

(A4) [Explicit form] $F_n = \sum_{i=1}^k p_i(n) \alpha_i^n$

where $p_1, \dots, p_k \in \mathbb{C}[z]$ $\deg p_i(n) < d_i \quad \forall i = 1, \dots, k$.

Proof: We define 4 sets, one for each (Ai) condition:

$$V_i := \{ F: \mathbb{N}_0 \rightarrow \mathbb{C} : (F_n) \text{ satisfies (Ai)} \} \quad i = 1, \dots, 4.$$

- Each V_i is a \mathbb{C} -vector space (functions are a.v.s & conditions (Ai) are linear in $(F_n)_n$)
- $\dim_{\mathbb{C}} V_i = d \quad \forall i = 1, \dots, 4$

Prf/ $F \in V_1$: $f(0), \dots, f(d-1)$ can be chosen arbitrarily & the rest are determined
 Basis $\{ \underbrace{(1, 0, \dots, 0)}_d, \underbrace{(0, 1, 0, \dots, 0)}_d, \dots, \underbrace{(0, \dots, 0, 1)}_d \} \Rightarrow V_1 \cong \mathbb{C}^d$

$F \in V_2$: f is determined by $P(z) = c_0 + c_1 z + \dots + c_{d-1} z^{d-1}, c_0, \dots, c_{d-1} \in \mathbb{C}$
 Again: $V_2 \cong \mathbb{C}^d \quad P \leftrightarrow (c_0, \dots, c_{d-1})$.

$F \in V_3$: Use (coeff $g_1, \dots, \text{coeff } g_k$) & $\sum_{i=1}^k d_i = d$ by construction.
 So $V_3 \cong \mathbb{C}^{d_1} \times \dots \times \mathbb{C}^{d_k} \cong \mathbb{C}^d$.

$F \in V_4$: Use (coeff $p_1, \dots, \text{coeff } p_k$) $\Rightarrow V_4 \cong \mathbb{C}^{d_1} \times \dots \times \mathbb{C}^{d_k} \cong \mathbb{C}^d$

Strategy: To show (Ai)'s are all equivalent, we must show $V_1 = V_2 = V_3 = V_4$.

By dimension, enough to show $V_2 \subseteq V_3$, $V_3 \subseteq V_2$, $V_3 \subseteq V_4$.

$V_2 \subseteq V_1$: Pick $f \in V_2$ Then $c(z) \overline{F}(z) = P(z)$

$$\sum_{n \geq 0} (1 + c_1 z + \dots + c_d z^d) F_n z^n = \sum_{n=0}^{\infty} \left(\sum_{i=0}^d c_i F_{n+d-i} \right) z^{n+d} = P(z)$$

Write $P(z) = a_0 + a_1 z + \dots + a_{d-1} z^{d-1} + 0 \cdot z^d + 0 z^{d+1} + \dots$ ($c_0 = 1$)

This gives $F_{n+d} + F_{n+d-1} c_1 + \dots + F_n c_d = 0 \quad \forall n$ (coeff z^{n+d})

So $f \in V_1$.

$V_3 \subseteq V_2$: Pick $f \in V_3$ Then:

$$\sum_{n=0}^{\infty} F_n z^n = \sum_{i=1}^k \frac{g_i(z)}{(1-d_i z)^{d_i}} = \frac{\sum_{i=1}^k g_i(z) \prod_{j \neq i} (1-d_j z)^{d_j}}{\prod_{i=1}^k (1-d_i z)^{d_i}} = \frac{P(z)}{C(z)}$$

deg < d
deg = $\sum_{j \neq i} d_j$
deg < d_i
= C(z).

and $\deg P \leq \max_i \left(\deg g_i + \sum_{j \neq i} d_j \right) < \max_i \left(\sum_{j=1}^k d_j \right) = d$

So $f \in V_2$.

$V_3 \subseteq V_4$: Pick $f \in V_3$ & look at $\frac{g_i(z)}{(1-d_i z)^{d_i}}$ for each $i=1, \dots, k$.

The Binomial Series (Lecture 17) gives

$$\frac{1}{(1-d_i z)^{d_i-1}} = \sum_{n \geq 0} \binom{d_i+n-1}{n} d_i^n z^n = \sum_{n \geq 0} \binom{d_i+n-1}{d_i-1} d_i^n z^n$$

Multiply by $g_i(z) = g_0^{(i)} + g_1^{(i)} z + \dots + g_{d_i-1}^{(i)} z^{d_i-1}$ gives:

$$\begin{aligned} \frac{g_i(z)}{(1-d_i z)^{d_i}} &= \sum_{n \geq 0} \left(\sum_{j=0}^{d_i-1} g_j^{(i)} \binom{d_i+n-j-1}{d_i-1} d_i^{n-j} \right) z^n \\ &= \sum_{n \geq 0} \underbrace{\left(\sum_{j=0}^{d_i-1} \frac{g_j^{(i)}}{d_i^j} \binom{d_i+n-j-1}{d_i-1} \right)}_{=: P_i(n)} d_i^n z^n \end{aligned}$$

Note: $\binom{d_i+n-j-1}{d_i-1} \in \mathbb{Q}[n]$ has degree = $d_i-1 \Rightarrow \deg P_i(n) \leq d_i-1$.

Conclude $\sum_{n=0}^{\infty} F_n z^n = \sum_{n=0}^{\infty} \left(\sum_{i=1}^k P_i(n) d_i^n \right) z^n$ & so $F_n = \sum_{i=1}^k P_i(n) d_i^n \quad \forall n$
So $f \in V_4$ \square