

Lecture 51

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(51.0) Recall that we discussed (in last lecture) irreducibility of a polynomial (one variable; coefficients from a unique factorization domain). We also made the following observation:

Observation: Let K be a field (e.g. $K = \mathbb{Q}, \mathbb{R}, \mathbb{C}, \mathbb{Z}/p\mathbb{Z}$)

Let $f(x) \in K[x], \alpha \in K$. Then ($p \in \mathbb{Z}_{\geq 2}$
prime)

$$(x - \alpha) \text{ divides } f(x) \iff f(\alpha) = 0$$

Proof. - (\Rightarrow) If $f(x) = (x - \alpha) \cdot g(x)$ then $f(\alpha) = 0 \cdot g(\alpha) = 0$.
(for some $g(x) \in K[x]$)

(\Leftarrow) Write $f(x) = (x - \alpha) \cdot q(x) + r(x)$ using division, so that $\text{degree}(r(x)) < \text{degree}(x - \alpha) = 1$. (divide $f(x)$ by $x - \alpha$)
i.e. $r(x) = r \in K$.

$$\text{Now } \left\{ \begin{array}{l} f(\alpha) = 0 \cdot q(\alpha) + r = r \\ f(\alpha) = 0 \quad (\text{given}) \end{array} \right\} \Rightarrow f(x) = (x - \alpha) \cdot q(x).$$

[Ex.] Carry out the same argument to prove:

$$(x - \alpha)^N \text{ divides } f(x) \iff f(\alpha) = \underbrace{f'(\alpha) = f''(\alpha) \dots = f^{(N-1)}(\alpha)}_{\text{derivatives of } f, f', \dots} = 0$$

$$\left[\text{Recall } f(x) = a_n x^n + a_{n-1} x^{n-1} + \dots + a_0 \right. \\ \left. \Rightarrow f'(x) = n a_n x^{n-1} + (n-1) a_{n-1} x^{n-2} + \dots + 2 a_2 x + a_1. \right]$$

Corollary of Observation: (Again $f(x) \in K[x]$).

(2)

$f(x)$ has m distinct roots in $K \Rightarrow m \leq n$.

[side note: we actually proved & used this result in Lecture 21 - see page 4 §21.3 - to conclude that

$$\text{Aut}_{\text{gp}}(\mathbb{Z}/p\mathbb{Z}) = \mathbb{F}_p^\times \text{ is cyclic.}]$$

Example. Let $x^2 + x + 1 \in \mathbb{F}_2[x]$. Then $\alpha \in \mathbb{F}_2$ can
 $f(x) =$

only be 0 or 1. $f(0) = 1$ and $f(1) = 3 \equiv 1 \pmod{2}$

i.e. $\forall \alpha \in \mathbb{F}_2$, $f(\alpha) \neq 0$. Thus $f(x)$ cannot be written as a product of two linear factors in $\mathbb{F}_2[x]$; hence it is irreducible.

(51.1) Polynomial rings in many variables - coefficients from a field K ($\text{many} \geq 1$).

Let $n \geq 1$ and let us consider

$$R = K[x_1, x_2, \dots, x_n]$$

ring of polynomials in n variables with coefficients from a field K .

Let us quickly review our proof of Hilbert Basis Theorem (3)

which implies that R is Noetherian (i.e., every ideal in R is finitely generated).

$$R = K[x_1, \dots, x_n] = \underbrace{(K[x_1, \dots, x_{n-1}])}_{\text{call it } A} [x_n]$$

call this variable u for now.

$$R = A[u]$$

$\overset{+U}{I}$ ideal \rightsquigarrow ideal in A generated by leading coefficients of $p(u) \in I \subseteq A[u]$
say $L(I) \subseteq A$

Step 1. - Use the fact that A is Noetherian to get

$$(a_1, \dots, a_k) = L(I) \text{ in } A$$

Pick $p_1(u), \dots, p_k(u) \in A[u]$ so that

Leading Coeff. of $p_j(u) = a_j$

Division algorithm \Rightarrow modulo $p_1(u), \dots, p_k(u)$

degree of any $q(u) \in A[u]$

can be brought down to be

(strictly) less than $\max_{1 \leq j \leq k} \{\deg(p_j(u))\}$

!!
D

Step 2 - Pick "generators" of $I < D$ (finitely many) " $\{ p(u) \in I \mid \deg(p(u)) < D \}$ (4)

Note: (1) by writing $u = x_n$ we made a choice: x_n is "better" than x_1, \dots, x_{n-1} (NOT democratic).

(2) Step 2 is not as "algorithmic" as Step 1.

We will fix these issues now.

(51.2) Monomials, orderings and related notation.

Monomial = Polynomial with only one term

(e.g. $x_1^2 x_2 \in K[x_1, x_2]$).

We can write a polynomial in $K[x_1, \dots, x_n]$ as

finit sum of monomials

(e.g. typical poly. in $K[x_1, x_2]$ is of the form

$$\sum_{k, l \geq 0} c(k, l) x_1^k x_2^l .)$$

finit (i.e. $k, l \leq N$ for some $N \in \mathbb{Z}_{\geq 0}$)

Convenient short-hand: \underline{x} for x_1, \dots, x_n ⑤

$\underline{\alpha} = \alpha_1, \dots, \alpha_n$ "exponents"; $\alpha_1, \dots, \alpha_n \in \mathbb{Z}_{\geq 0}$

$\underline{x}^{\underline{\alpha}} = x_1^{\alpha_1} x_2^{\alpha_2} \dots x_n^{\alpha_n}$: typical monomial. } in $K[x_1, \dots, x_n]$

$\sum_{\substack{\underline{\alpha} \in (\mathbb{Z}_{\geq 0})^n \\ \text{finite sum}}} c(\underline{\alpha}) \underline{x}^{\underline{\alpha}}$: typical polynomial.

Definition: A monomial order is a total order on the set of monomials $\left\{ \underline{x}^{\underline{\alpha}} \mid \underline{\alpha} \in (\mathbb{Z}_{\geq 0})^n \right\}$

s.t. $\underline{x}^{\underline{\alpha}} \leq \underline{x}^{\underline{\beta}}$

$\Rightarrow \underline{x}^{\underline{\alpha} + \underline{\gamma}} \leq \underline{x}^{\underline{\beta} + \underline{\gamma}} \quad \forall \underline{\gamma} \in (\mathbb{Z}_{\geq 0})^n$

(51.3) For example:

Dictionary order or lexicographic order - we fix an

arbitrary ordering on variables, and then decide

which monomial is "bigger" according to the following rule
(how to look up words in a dictionary!)

$$\boxed{x_1^{k_1} x_2^{k_2} x_3^{k_3} \geq x_1^{l_1} x_2^{l_2} x_3^{l_3}}$$

iff $k_1 > l_1$
 or $k_1 = l_1$ and $k_2 > l_2$
 or $k_1 = l_1$; $k_2 = l_2$ and $k_3 > l_3$

(to fix ideas - our language has 3 alphabets x_1, x_2, x_3)
 $\parallel x_1$ comes before x_2 .
 $\parallel x_2$ comes before x_3 .
 (our choice - can be changed).
 $x_1 > x_2 > x_3$

e.g. $x_1^2 x_2^1 x_3^1 \geq x_1^2 x_2^1 x_3^0$

(5.4) Leading term :- Let us assume we fixed a monomial order for monomials in $K[x_1, \dots, x_n]$.

For $f(x_1, \dots, x_n) \in K[x_1, \dots, x_n]$
 \parallel
 $\sum c(\alpha) x^\alpha$

LT(f) = $c(\alpha_0) x^{\alpha_0}$

leading term of f

where α_0 is such that

- $c(\alpha_0) \neq 0$
- $c(\alpha) \neq 0 \Rightarrow \alpha_0 \geq \alpha$

Convention: $LT(0) = 0$

let $I \subsetneq K[x_1, \dots, x_n]$ be ^(an) ideal. Define, similarly (7)

$$LT(I) = \text{ideal in } K[x_1, \dots, x_n] \text{ generated by} \\ \{LT(f) \mid f \in I\}$$

(51.5) Example. Say $n=2$ and let us just write x & y for our variables. So, $R = K[x, y]$. Let us also choose lexicographic order, with $x > y$.

$$f(x, y) = x^3y - xy^2 + 1 \quad \Rightarrow \quad LT(f) = x^3y$$

$$g(x, y) = x^2y^2 - y^3 - 1 \quad \Rightarrow \quad LT(g) = x^2y^2$$

Catch: $LT(f), LT(g)$ do not generate $LT(I)$

if $I = (f, g) \subset K[x, y]$.

e.g. $x + y = y \cdot f - x \cdot g \in I$

$$\Rightarrow LT(I) \ni x$$