

Lecture 10

0. Summary so far. Recall (from Lecture 4) that we introduced the following rational connection and obtained a criterion for its flat-ness (or consistency)

$$\nabla = d - \sum_{x \in X} \frac{dx}{x} t_x$$

In the set up of lecture 4; we fixed $V =$ a finite-dimensional \mathbb{C} -vector space; $X \subset V^* \setminus \{0\}$ a finite set of linear forms. We proved

Kohno's Lemma: ∇ is flat (i.e. defines a consistent system of PDE's) \Leftrightarrow for every subset $Y \subset X$ maximal s.t. $\text{Span } Y$ is 2-dim'l

we have
$$\left[\sum_{y \in Y} t_y, t_z \right] = 0.$$

Now we are going to assume that the underlying hyperplane arrangement (i.e. $X \subset V^*$ or $\{H_x = \text{Ker}(x)\}_{x \in X}$ - hyperplanes in V) comes from a root system (see Lecture 5, §3).

1. Complexification. Let $\underline{E}, R \subset E^* \setminus \{0\}$ be a finite root system.

Set $\mathfrak{h} = E \otimes_{\mathbb{R}} \mathbb{C}$ (i.e. extend scalars from \mathbb{R} to \mathbb{C}).

Each $\alpha: E \rightarrow \mathbb{R}$ is then extended to a \mathbb{C} -linear form

The role of " $X \in V^* \setminus \{0\}$ " will be played by R_+ = set of positive roots. Thus our rational connection will take the

following form:
$$\nabla = d - \sum_{\alpha \in R_+} \frac{d\alpha}{\alpha} \cdot t_\alpha$$

For concrete-ness, let us fix another f.d. \mathbb{C} -vector space, say F , and assume that $t_\alpha \in \text{End}(F) \quad \forall \alpha \in R_+$.

As $\frac{d(-\alpha)}{(-\alpha)} = \frac{d\alpha}{\alpha}$, we are going to consider $t_\alpha \in \text{End} F$

for all roots $\alpha \in R$ and impose $t_\alpha = t_{-\alpha}$.

2. W -action. Let us assume that we are given, in addition, an action of W on F . That is, a group homomorphism $W \xrightarrow{a} GL(F)$. This allows us to conjugate elements of $\text{End} F$

by $w \in W$ as: for $w \in W$ $w \cdot t := a(w) t a(w)^{-1} \in \text{End} F$
 $t \in \text{End}(F)$

We say the connection ∇ is W -equivariant if $w \cdot t_\alpha = t_{w(\alpha)} \quad \forall w \in W$ and $\alpha \in R$.

3. A W -equivariant connection $\nabla = d - \sum_{\alpha \in R_+} \frac{d\alpha}{\alpha} t_\alpha$

allows us to define the monodromy representation, which is

a group homomorphism $\mu : \pi_1(\mathfrak{h}^{reg}/W) \rightarrow GL(F)$.

[Notation: $\mathfrak{h}^{reg} = \mathfrak{h} \setminus \bigcup_{\alpha \in R_+} H_\alpha$ "regular elements of \mathfrak{h} ".]

μ in fact depends on

- choice of a base point in \mathfrak{h}^{reg}/W ; say p_0 .
(defn. of π_1 involves choosing a base point)

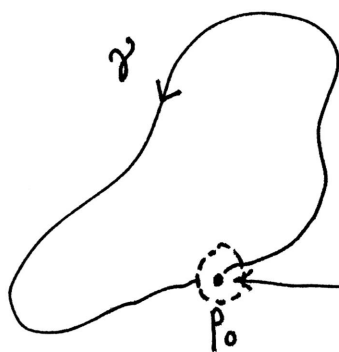
- a $GL(F)$ -valued solution ψ of $\nabla\psi = 0$; near p_0 .

4. Definition of $\mu_{(\psi, p_0)}$: Given a loop $\gamma : [0, 1] \rightarrow \mathfrak{h}^{reg}/W$

consider the analytic continuation of ψ along γ

$$\mu(\gamma) := \psi^{-1} \cdot \tilde{\psi} \in GL(F)$$

(constant since both ψ & $\tilde{\psi}$ solve the same equation).



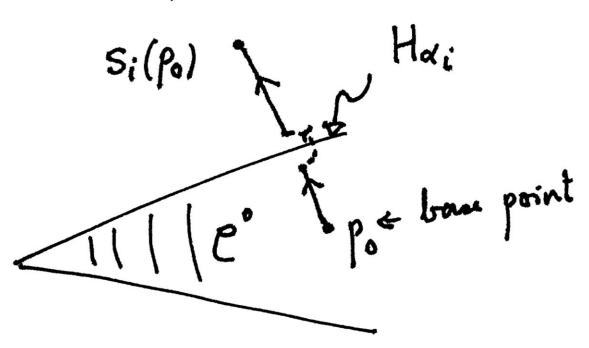
two solutions ψ and $\tilde{\psi}$
//
continuation of ψ along γ .

5. The fundamental group $\pi_1(\mathfrak{h}^{\text{reg}}/W)$ was computed by Brieskorn (DIE FUNDAMENTALGRUPPE DES RAUMES DER REGULÄREN ORBITS EINER ENDLICHEN KOMPLEXEN SPIEGELUNGSGRUPPE - Invent Math 12 (57-61). 1971).

Theorem. $\pi_1(\mathfrak{h}^{\text{reg}}/W, p_0) = B_W$ (braid group of W) (see Lecture 9, §4, page 7).

The base point is assumed to be in the fundamental chamber (i.e. $p_0 \in \mathfrak{h}^{\text{reg}}$ is such that $\alpha_i(p_0) \in \mathbb{R}_{>0}$.) For each $i \in I$, the generator T_i of B_W is then mapped to the following loop:

γ_i joins p_0 to $s_i(p_0)$ by approaching the wall H_{α_i} in the real vector space E ; and circling around it in a counterclockwise direction - in the imaginary coordinates.



Picture of the loop γ_i ($i \in I$).

6. Brieskorn's Theorem was reproved and extended ($\pi_k(\mathfrak{h}^{\text{reg}}/W) = \{1\}$ $\forall k \geq 2$)

by Deligne - Les immeubles des groupes de tresses généralisés (Invent. Math. 1972).

7. Good / fundamental solutions.

[De Concini, Procesi - Hyperplane arrangements and holonomy equations
Selecta Math. (1995).]

Now we will try to find solutions ψ of $\nabla\psi = 0$ for which the monodromy is "easy to compute". This should be viewed as several variable generalization of (Lecture 1 §4):

$$F'(z) = \left(\frac{\Lambda}{z} + A_{\text{reg}}(z) \right) F \quad \leadsto \text{fundamental solution } H(z). z^{\Lambda} = \psi$$

$$\leadsto \mu_{\psi} \left(\begin{array}{c} \circlearrowleft \\ \circ \end{array} \right) = e^{2\pi i \Lambda}$$

The problem is that of "normal crossing" (i.e., we usually have too many hyperplanes passing through $0 \in \mathfrak{h}$. To make this (more than $\dim_{\mathbb{C}} \mathfrak{h}$)

point clear, let us ~~prove~~ that consider the "normal crossing" case first

8. Let $n \geq 2$ and assume we have the following system of PDE's:

$$\frac{\partial f}{\partial x_i} = \left(\frac{t_i}{x_i} + R_i(\underline{x}) \right) f \quad (1 \leq i \leq n). \quad (*) \quad (6)$$

Here $f(x_1, \dots, x_n)$ is (unknown) function of n variables; taking values in $GL(F)$; $t_1, \dots, t_n \in \text{End}(F)$; and $R_i(\underline{x})$ are $\text{End}(F)$ -valued, holomorphic functions near $\underline{x} = \underline{0}$.

Exercise: $(*)$ is consistent \Leftrightarrow $[t_i, t_j] = 0 \quad \forall i, j$
 $[t_i, R_j] = 0 \quad \forall i \neq j$

$$\frac{\partial R_i}{\partial x_j} - \frac{\partial R_j}{\partial x_i} + [R_i, R_j] = 0 \quad \forall i \neq j$$

Exercise: $\exists!$ solution of $(*)$ of the form (under a non-resonance hypothesis akin to the one from Thm. 4 of Lecture 1, page 6.)

$$\psi(\underline{x}) = H(x_1, \dots, x_n) \cdot \prod_{i=1}^n x_i^{t_i}$$

where

unambiguous since t_i 's commute

$H(\underline{0}) = 1$, and H is holomorphic near $\underline{x} = \underline{0}$.

Then:

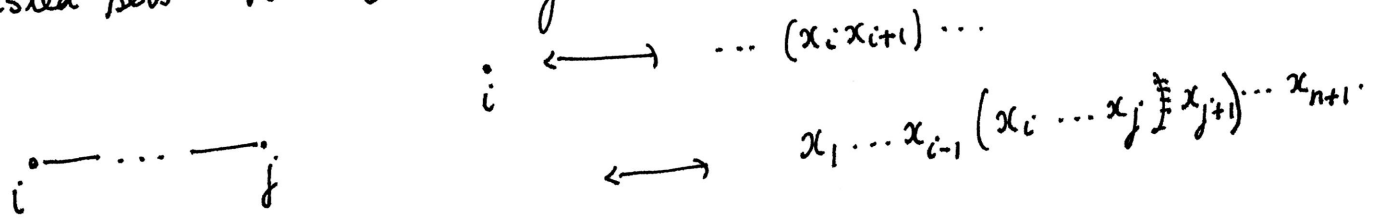
$$\mu_\psi \left(\overset{x_i=0}{\cancel{p_0}} = (1, 1, \dots, 1) \right) = \exp(2\pi\sqrt{-1} t_i)$$

A maximal nested set for $n=3$

$$F_1 = \left\{ \overset{\cdot}{1} \text{---} \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{1}, \overset{\cdot}{3} \right\}; \quad F_2 = \left\{ \overset{\cdot}{1} \text{---} \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{1} \text{---} \overset{\cdot}{2}, \overset{\cdot}{3} \right\}$$

$$F_3 = \left\{ \overset{\cdot}{1} \text{---} \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{1} \text{---} \overset{\cdot}{2}, \overset{\cdot}{2} \right\}$$

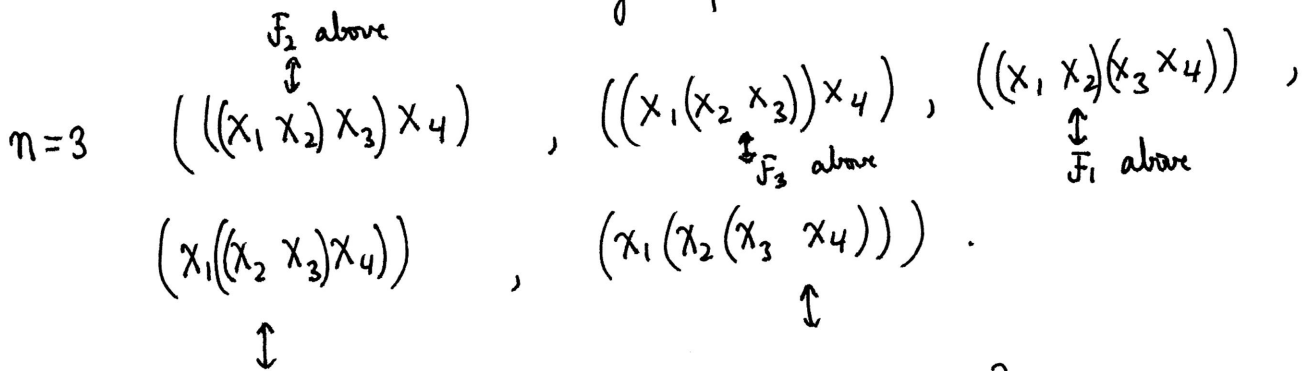
Nested sets vs. bracketings on $x_1 x_2 \dots x_{n+1}$



Connected subdiagram of A_n

• Compatibility - same for brackets.

• Maximal nested set - complete bracketings
 eg. for $n=2$ $((x_1 x_2) x_3)$ or $(x_1 (x_2 x_3))$



$$\left\{ \overset{\cdot}{1} \text{---} \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{2}, \overset{\cdot}{2} \text{---} \overset{\cdot}{3} \right\} \quad \left\{ \overset{\cdot}{1} \text{---} \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{2} \text{---} \overset{\cdot}{3}, \overset{\cdot}{3} \right\}$$