

## LIE GROUPS: HOMEWORK 4 (HINT FOR PROBLEM 4)

**Problem 4.** Let  $W$  be a group (not necessarily finite) generated by a finite subset of distinct elements  $S = \{s_1, \dots, s_N\} \subset W$  such that  $s_i \neq 1$  and  $s_i^2 = 1$  for each  $1 \leq i \leq N$ . As we did in class, we can define the length function on  $W$  as: for each  $w \in W$ ,  $l(w)$  is the smallest non-negative integer such that  $w$  can be expressed as a product of  $l(w)$  elements of  $S$ . An expression of  $w = s_{i_1} \cdots s_{i_l}$  where  $l = l(w)$  is called a reduced expression of  $w$ .

Assume the following property (exchange property): for  $i \in \{1, \dots, N\}$  and  $w \in W$ , if  $l(ws_i) < l(w)$ , then for any reduced expression  $w = s_{i_1} \cdots s_{i_l}$  there exists  $j$ ,  $1 \leq j \leq l$  such that

$$s_{i_j} \cdots s_{i_l} = s_{i_{j+1}} \cdots s_{i_l} s_i$$

For  $s, s' \in S$  define  $m(s, s')$  to be the order of  $ss' \in W$  (possibly infinity).

Prove the following universal property: Let  $G$  be a group and let  $f : S \rightarrow G$  be a set map such that  $f(s)^2 = 1$  for every  $s \in S$  and for every  $s, s' \in S$  with  $m(s, s') < \infty$ , we have

$$\underbrace{f(s)f(s')f(s) \cdots}_{m(s, s') \text{ terms}} = \underbrace{f(s')f(s)f(s') \cdots}_{m(s, s') \text{ terms}}$$

Then there exists a unique group homomorphism  $g : W \rightarrow G$  such that  $g(s) = f(s)$  for every  $s \in S$ .

**HINT FOR A PROOF.** The only non-trivial part is to prove that the following extension  $g : W \rightarrow G$  is well defined. For any  $w \in W$ , pick a reduced expression  $w = s_{i_1} \cdots s_{i_l}$  and define  $g(w) = g(w; s_{i_1} \cdots s_{i_l}) := f(s_{i_1}) \cdots f(s_{i_l})$  (to emphasize the dependence on the reduced expression chosen). One needs to show that the right-hand side of this is independent of the reduced expression. Once this is done, the rest of the proof is not so hard!

To prove independence of  $g(w)$  from the choice of a reduced expression, one proceeds by induction on the length of the reduced expression. The base case is trivial. For the induction step, start from two reduced expressions  $w = s_{i_1} \cdots s_{i_l} = s_{j_1} \cdots s_{j_l}$ . Observe that if  $i_1 = j_1$  or  $i_l = j_l$  the induction step will imply the result. Otherwise, apply the exchange property to the reduced expression  $s_{i_1} \cdots s_{i_l}$  and  $i = j_l$ , to get another reduced expression

$$s_{i_1} \cdots s_{i_{t-1}} \widehat{s_{i_t}} s_{i_{t+1}} \cdots s_{i_l} s_{j_l}$$

Now this one shares the last term with  $s_{j_1} \cdots s_{j_l}$  and first term with  $s_{i_1} \cdots s_{i_l}$ , unless  $t = 1$ . So, if  $t \neq 1$ , we are done by previous observation. The case when  $t = 1$  gives  $s_{i_2} \cdots s_{i_l} s_{j_l}$  as the new reduced expression. Repeated application of this argument (and also for the role of  $i$ 's and  $j$ 's flipped), will result in

$$w = \cdots s_{i_l} s_{j_l} s_{i_l} s_{j_l} = \cdots s_{j_l} s_{i_l} s_{j_l} s_{i_l}$$

at which point we are done by the hypothesis on  $f$ .