

Lecture 9

(9.0) Recall: we introduced Killing form $K(\cdot, \cdot): \mathfrak{g} \times \mathfrak{g} \rightarrow \mathbb{C}$ (or \mathbb{R})

$$K(X, Y) := \text{Trace}_{\mathfrak{g}}(\text{ad}(X) \circ \text{ad}(Y))$$

Invariance: $K([X, Y], Z) + K(Y, [X, Z]) = 0$

Thm. (Cartan's criterion): (A) \mathfrak{g} is solvable $\Leftrightarrow K(X, Y) = 0 \quad \forall X \in \mathfrak{g}, Y \in [\mathfrak{g}, \mathfrak{g}]$

(B) \mathfrak{g} is semisimple $\Leftrightarrow K$ is non-degenerate

Last time we proved (B) assuming (A).

(9.1) Proof of A (easy part) [work over \mathbb{C} , see Remark in Lecture 8. section (8.9) on page 8]

\Rightarrow). Assume \mathfrak{g} is solvable. Then $[\mathfrak{g}, \mathfrak{g}]$ acts on \mathfrak{g} by nilpotent matrices (use Lie's Thm for the adjoint repn.)

So $\text{Tr}(\text{ad}(X) \cdot \text{ad}(Y)) = 0 \quad \forall X \in \mathfrak{g}, Y \in [\mathfrak{g}, \mathfrak{g}]$ since

$$\text{ad}(X) = \begin{bmatrix} * & * \\ 0 & * \end{bmatrix} \quad \text{and} \quad \text{ad}(Y) = \begin{bmatrix} 0 & * \\ 0 & 0 \end{bmatrix}. \quad \square$$

(9.2) Proof of A. (hard part). For this we need Jordan's decomposition thm and its relatively easier version

Jordan decomposition. Let $X \in M_{n \times n}(\mathbb{C})$. Then $\exists!$ X_S, X_N in $M_{n \times n}(\mathbb{C})$ s.t (i) X_S is diagonalisable, X_N is nilpotent

(ii) $X = X_S + X_N$ (iii) X_S and X_N are polynomials in X ②

(i.e. $\exists \beta, N \in \mathbb{C}[T]$ s.t. $X_S = \beta(X)$ and $X_N = N(X)$). In particular

$$[X_S, X_N] = 0.$$

Moreover $\text{ad}(X) \hookrightarrow M_{n \times n}(\mathbb{C})$ has s.s. and nilpotent parts $\text{ad}(X_S)$, $\text{ad}(X_N)$ resp. That is $\text{ad}(X)_S = \text{ad}(X_S)$

$$\text{ad}(X)_N = \text{ad}(X_N)$$

Note: if eigenvalues of X_S are $\{\alpha_1, \alpha_2, \dots, \alpha_n\}$ then eigenvalues of $\text{ad}(X_S)$ are $\{\alpha_i - \alpha_j\}_{1 \leq i, j \leq n}$.

(9.3) Proof of A. (\Leftarrow). We need to prove that \mathfrak{g} is solvable.

We will show that $[\mathfrak{g}, \mathfrak{g}]$ is nilpotent (which implies \mathfrak{g} is solvable by defn.). Let $X \in [\mathfrak{g}, \mathfrak{g}]$. By Engel's Thm and its

Corollary (8.7) of Lecture 8 on page 7, it is enough to prove

that $Y = \text{ad}(X)$ is nilpotent. ($Y \in M_{n \times n}(\mathbb{C})$, $n = \dim \mathfrak{g}$)

Assume on the contrary, $Y = Y_S + Y_N$, $Y_S \neq 0$.

$$\text{Then } \text{Tr}(\bar{Y}_S Y_S) = \sum_{\alpha: \text{eigenvalue of } Y_S} \bar{\alpha} \alpha > 0 \quad (*)$$

Complex Conjugate

$$= \text{Tr}(\bar{Y}_S (Y_S + Y_N)) = \text{Tr}(\bar{Y}_S Y)$$

(since $\bar{Y}_S Y_N$ is nilpotent)

Since $X \in [\mathfrak{g}, \mathfrak{g}]$ we can write $X = \sum_i [X_i', X_i'']$ (3)

and consequently $Y = \sum_i [Y_i', Y_i'']$ where $Y_i' = \text{ad}(X_i')$.

$$\text{So } \text{Tr}(\bar{Y}_s Y) = \sum_i \text{Tr}(\bar{Y}_s \cdot [Y_i' Y_i''])$$

$$= \sum_i \text{Tr}([\bar{Y}_s, Y_i'] \cdot Y_i'')$$

$$= \sum_i \text{Tr}(\overline{(\text{ad } Y)_s} (Y_i') \cdot Y_i'')$$

[see the proof of invariance of K in (8.8) of Lecture 8 (p.7)]

$(\text{ad } Y)_s$ is a polynomial in $\text{ad } Y$. Thus, so is $\overline{(\text{ad } Y)_s}$ (R coeff.)
 (* - see page 10) by (9.2)

$$\text{Now } (\text{ad } Y)(Y_i') = \text{ad}([X, X_i'])$$

$$(\text{ad } Y)^2(Y_i') = \text{ad}([X, [X, X_i']]) \text{ and so on...}$$

Thus $\text{Tr}(\bar{Y}_s Y)$ is a linear combination of terms of the form

$$\text{Tr}(\text{ad}([X, [X, \dots [X, X_i'] \dots]]) \text{ad}(X_i'')) = 0 \text{ by assumption}$$

↓
 $\in [\mathfrak{g}, \mathfrak{g}]$

Hence we contradicted (*) $\Rightarrow Y_s = 0$. So $Y = \text{ad}(X)$ is

nilpotent $\forall X \in [\mathfrak{g}, \mathfrak{g}]$ and thus $[\mathfrak{g}, \mathfrak{g}]$ is nilpotent \square

(9.4) From now on $\mathfrak{g} =$ a semisimple Lie alg. over \mathbb{C} . We record a few consequences of non-degeneracy of K .

(4)

Cor (i) Let $\mathfrak{a} \subset \mathfrak{g}$ be an ideal. Then $\mathfrak{b} = \mathfrak{a}^\perp = \{ X \in \mathfrak{g} : K(X, Y) = 0 \forall Y \in \mathfrak{a} \}$ is another ideal, and $\mathfrak{g} = \mathfrak{a} \oplus \mathfrak{b}$.

[Proof: \mathfrak{b} is an ideal: let $X \in \mathfrak{g}, Y \in \mathfrak{b}$. For every $Z \in \mathfrak{a}$, $K([Y, X], Z) = K(Y, [X, Z]) = 0 \Rightarrow [Y, X] \in \mathfrak{b}$.
 \mathfrak{a} as \mathfrak{a} is an ideal

If $\mathfrak{a} \cap \mathfrak{b} \neq (0)$ Then $K \equiv 0$ on $\mathfrak{a} \cap \mathfrak{b} \Rightarrow \mathfrak{a} \cap \mathfrak{b}$ is a solvable ideal contradicting semisimplicity of \mathfrak{g} . Finally as K is non-deg. (by A. of Thm (9.0))
 $\mathfrak{g} = \mathfrak{a} \oplus \mathfrak{b}$]

(ii) $\mathfrak{g} = [\mathfrak{g}, \mathfrak{g}]$.

[Proof. Since $[\mathfrak{g}, \mathfrak{g}]$ is an ideal, its complement \mathfrak{a} must be an ideal as well. But then $\mathfrak{a} \cong \mathfrak{g}/[\mathfrak{g}, \mathfrak{g}] \leftarrow$ abelian ideal in \mathfrak{g} contradicts s.s. (unless $\mathfrak{a} = 0$). $\mathfrak{a} = 0$ means $\mathfrak{g} = [\mathfrak{g}, \mathfrak{g}]$]

(iii) $\mathfrak{g} = \mathfrak{g}_1 \oplus \dots \oplus \mathfrak{g}_n$ \mathfrak{g}_j 's are simple Lie alg.
(min'l ideals of \mathfrak{g}).

[Easy consequence of (i) and induction on $\dim \mathfrak{g}$].

(iv) $\mathfrak{g} \xrightarrow{\text{ad}} \text{End}_{\mathbb{C}}(\mathfrak{g})$ is injective. (5)

[$\text{Ker}(\text{ad}) = \mathcal{Z}(\mathfrak{g}) = \text{center of } \mathfrak{g} := \{X \in \mathfrak{g} : [X, Y] = 0 \ \forall Y \in \mathfrak{g}\}$
is an abelian ideal of \mathfrak{g} , hence 0].

(9.5) More generally, let $\mathfrak{g} \subset V$ faithfully (i.e. $X \cdot v = 0 \ \forall v \in V \Rightarrow X = 0$). Let $B_V(\cdot, \cdot) : \mathfrak{g} \times \mathfrak{g} \rightarrow \mathbb{C}$ be defined by

$$B_V(X, Y) = \text{Trace}_V(\underbrace{X_V \circ Y_V}_{X_V, Y_V \in V \text{ matrices.}})$$

Assume V is simple*

Lemma. (i) B_V is non-degenerate and invariant.

(ii) Let $\{X_i\}$ and $\{X^i\}$ be two bases of \mathfrak{g} st.

$$B_V(X_i, X^j) = \delta_{ij}$$

$$\text{Then } \sum_i X_{i,V} X^i_V = \frac{\dim \mathfrak{g}}{\dim V} \cdot \text{Id}_V.$$

(* a representation V is simple / irreducible if 0 and V are the only subrepresentations).

Proof of (i) $B_V([X, Y], Z) + B_V(Y, [X, Z]) = 0$ (same as the

one for $K(\cdot, \cdot)$ [Lecture 8, (8.8)].

Now identify \mathfrak{g} with a subalg. of $\text{End}(V)$ and let

$$\mathfrak{a} = \{X \in \mathfrak{g} : B_V(X, Y) = 0 \ \forall Y \in \mathfrak{g}\} \subseteq \text{ideal } \mathfrak{g}$$

||
Tr(X, Y) via this identification

⑥

Using the same argument as in (9.3) we can prove that

$[\alpha, \alpha]$ is nilpotent, hence $\alpha \in \mathfrak{g}$ is solvable $\Rightarrow \alpha = 0$, i.e. B_V is non-degenerate.

(ii) We claim that $[X, \sum_{i=1}^n X_i X^i] = 0 \quad \forall X \in \mathfrak{g}$. ($n = \dim \mathfrak{g}$)

If true, then $C_V = \sum X_i X^i : V \rightarrow V$ is a hom. of \mathfrak{g} representations

Since V is simple $C_V = \lambda \cdot \text{Id}_V$ (let λ be an eigenvalue of C_V ;

then $\text{Ker}(C_V - \lambda \text{Id}_V)$ is a non-zero \mathfrak{g} -subrepr. of V , hence $= V$)

$\text{Trace}_V(C_V) = \lambda \cdot (\dim V) = \sum_i B_V(X_i, X^i) = \dim \mathfrak{g}$ [Schur's Lemma]

$\Rightarrow \lambda = \frac{\dim \mathfrak{g}}{\dim V}$ as required.

Proof of $[X, \sum X_i X^i] = 0$: let $[X, X_i] = \sum_j a_{ij}^i X_j$
 $[X, X^i] = \sum_j b_j^i X^j$

Then $[X, \sum_i X_i X^i] = \sum_i [X X_i] X^i + \sum_i X_i [X X^i]$
 $= \sum_{i,j} a_{ij}^i X_j X^i + \sum_{i,j} b_j^i X_i X^j = \sum_{i,j} (a_{ij}^i + b_i^j) X_i X^j$

Now $a_i^i = B([X, X_i], X^i) = -B(X_i, [X X^i]) = -b_i^i$

$\Rightarrow a_i^i + b_i^i = 0$

□

(9.6) Thm. Let V be a (f.d.) repr. of \mathfrak{g} and $V' \subset V$ be a subrepr. Then $\exists V'' \subset V$ a subrepr. s.t. $V = V' \oplus V''$ as repr. of \mathfrak{g} . (7)

Proof. We have $0 \rightarrow V' \rightarrow V \xrightarrow{\text{pr}} V/V' \rightarrow 0$ a short exact

sequence of \mathfrak{g} -reprs. We need to find a hom. of \mathfrak{g} -reprs.

$V/V' \xrightarrow{s} V$ s.t. $\text{pr} \circ s = \text{Id}_{V/V'}$. Then $V'' = \text{Image of } s$

is the required subrepr. of V .

[Aside - review of some definitions. Let A, B be two reprs of \mathfrak{g} .

$$\text{Hom}_{\mathbb{C}}(A, B) \supset \text{Hom}_{\mathfrak{g}}(A, B) := \left\{ f: A \rightarrow B \text{ such that } \begin{array}{l} f(X \cdot a) = X \cdot f(a) \quad \forall a \in A \\ X \in \mathfrak{g} \end{array} \right\}$$

\mathfrak{g} acts naturally on $\text{Hom}_{\mathbb{C}}(A, B)$:

$$\forall \xi \in \text{Hom}_{\mathbb{C}}(A, B), X \in \mathfrak{g}, a \in A$$

$$(X \cdot \xi)(a) = X \cdot (\xi(a)) - \xi(X \cdot a)$$

So that

$$\text{Hom}_{\mathfrak{g}}(A, B) = \left\{ \xi \in \text{Hom}_{\mathbb{C}}(A, B) \text{ such that } \begin{array}{l} X \cdot \xi = 0 \\ \forall X \in \mathfrak{g} \end{array} \right\}$$

Continuing with the proof, pick any $l \in \text{Hom}_{\mathbb{C}}(V/V', V)$

s.t. $\text{pr} \circ l = \text{Id}_{V/V'}$. We will modify l so that

it lands in $\text{Hom}_{\mathfrak{g}}(V/V', V)$ still satisfying $\text{pr} \circ l = \text{Id}_{V/V'}$.

Consider $f: \mathfrak{g} \longrightarrow \text{Hom}_{\mathbb{C}}(V/V', V')$. We claim that $\textcircled{8}$

$$\begin{array}{ccc} \mathfrak{g} & \longrightarrow & \text{Hom}_{\mathbb{C}}(V/V', V') \\ \cup & & \cup \\ X & \longmapsto & X \cdot \ell \end{array}$$

$f(X) \in \text{Hom}_{\mathbb{C}}(V/V', V')$. i.e. $\text{pr} \circ f(X) = 0 \quad \forall X \in \mathfrak{g}$. This is

clear since $(\forall a \in V/V')$

$$\begin{aligned} \text{pr} \circ f(X)(a) &= (\text{pr} \circ (X \cdot \ell))(a) = \text{pr}((X \cdot \ell)(a)) \\ &= \text{pr}(X \cdot (\ell(a)) - \ell(X \cdot a)) = X \cdot (\text{pr}(\ell(a)) - (\text{pr} \circ \ell)(X \cdot a)) \\ &= X \cdot a - X \cdot a = 0 \end{aligned}$$

(since $\text{pr} \circ \ell = \text{id}_{V/V'}$). (since $\text{pr}: V \rightarrow V/V'$ is a hom of \mathfrak{g} -reps.)

Moreover $f([X, Y]) = X \cdot f(Y) - Y \cdot f(X)$ (after all $f(X) = X \cdot \ell$)

Crucial Lemma: Let R be a repn. of \mathfrak{g} and assume we

have a \mathbb{C} -linear map $f: \mathfrak{g} \rightarrow R$ s.t.

$$f([X, Y]) = X \cdot f(Y) - Y \cdot f(X).$$

Then $\exists r \in R$ s.t. $f(X) = X \cdot r$

Apply it to $R = \text{Hom}_{\mathbb{C}}(V/V', V')$ and $f: \mathfrak{g} \rightarrow \text{Hom}(V/V', V')$

$$X \longmapsto X \cdot \ell$$

to find $m \in \text{Hom}_{\mathbb{C}}(V/V', V')$ s.t. $X \cdot \ell = X \cdot m$

i.e. $X \cdot (\ell - m) = 0 \quad (\text{pr} \circ (\ell - m) = \text{pr} \circ \ell = \text{id})$

i.e. $\ell - m \in \text{Hom}_{\mathfrak{g}}(V/V', V)$. \square

(9.7) Proof of Crucial Lemma.

(9)

Case 1. R is irreducible and $\sigma \rightarrow \text{End}(R)$ is injective.

Use Lemma (9.5). We claim $r = \frac{1}{\lambda} \sum_i x_i \cdot f(x^i) \in R$.

We need to prove that $f(x) = x \cdot r \quad (\forall x \in \sigma)$.

$$\begin{aligned} \lambda (f(x) - x \cdot r) &= \sum_i x_i \cdot (x^i \cdot f(x)) - \sum_i x \cdot (x_i \cdot f(x^i)) \\ &= - \sum_i [x, x^i] \cdot f(x^i) + x_i \cdot f([x, x^i]) \quad (\text{by assumption on } f) \\ &\quad \uparrow \qquad \qquad \qquad \uparrow \\ &\quad \sum_j a_i^j x_j \qquad \qquad \qquad \sum_j b_i^j x_j \\ &= - \sum_{i,j} (a_i^j + b_i^j) x_j \cdot f(x^i) = 0 \end{aligned}$$

Case 2. R is irreducible. Let $k = \{x \in \sigma : x \cdot a = 0 \forall a \in R\}$.

$k \subset \sigma$ is an ideal, hence $\sigma = k \oplus b$ (Cor 9.4 (i) p. 4)

(both k and b s.s.). By Case 1, $\exists r \in R$ s.t. $f(x) = x \cdot r \quad \forall x \in b$.

Now $f([y, z]) = y \cdot f(z) - z \cdot f(y) = 0 \quad \forall y, z \in k$.

Since $k = [k, k]$ (as k is semisimple Cor 9.4 (ii))

$0 = f(x) = x \cdot r \quad \forall x \in k$.

Case 3. R arbitrary. By induction on $\dim R$. (10)

If R is not irreducible, we have a subrepr $R' \subsetneq R$ ($0 \neq R'$)

Apply induction hypothesis to $\begin{array}{ccc} \mathfrak{g} & \longrightarrow & R/R' \\ \downarrow \psi & & \downarrow \text{pr} \\ X & \longmapsto & \text{pr}(f(X)) \end{array}$ where

$\text{pr}: R \rightarrow R/R'$. We find $r_1 \in R/R'$ st. $\text{pr}(f(X)) = X \cdot r_1$

Choose any $\tilde{r}_1 \in R$ st. $\text{pr}(\tilde{r}_1) = r_1$. Then we have

$$\text{pr}(f(X)) = \text{pr}(X \cdot \tilde{r}_1) \quad \text{since pr is a hom of } \mathfrak{g}\text{-reprs.}$$

Apply induction again to $\begin{array}{ccc} \mathfrak{g} & \longrightarrow & R' \\ X & \longmapsto & f(X) - X \cdot \tilde{r}_1 \end{array}$ and find

$$r_2 \in R' \text{ st. } f(X) - X \cdot \tilde{r}_1 = X \cdot r_2. \quad \text{Then } r = r_2 + \tilde{r}_1 \quad \square$$

(9.8) Explanation of (*) on page 3. We know $\exists \delta(T) \in \mathbb{C}[T]$ st.

$$\delta(X) = X_S = \begin{bmatrix} \lambda_1 & & 0 \\ & \ddots & \\ 0 & & \lambda_n \end{bmatrix}. \quad \text{Take } f(t) \in \mathbb{C}[t] \text{ st.}$$

$$f(0) = 0 \text{ and } f(\lambda_i) = \bar{\lambda}_i \quad (1 \leq i \leq n)$$

(in a suitable basis)
Then $(f \circ \delta)(X) = \bar{X}_S$.

(eg. if $\alpha_1, \dots, \alpha_k$ are distinct complex #'s

$$\sum_i \bar{\alpha}_i \prod_{j \neq i} \frac{(t - \alpha_j)}{\alpha_i - \alpha_j} \text{ maps } \alpha_i \mapsto \bar{\alpha}_i \quad \forall (1 \leq i \leq k)$$