

Lecture 16

①

Recall that last time we introduced the notion of homotopy of two morphisms $\alpha, \beta: C \rightarrow D$. The terminology comes from algebraic topology: $X: \text{a topological space} \rightsquigarrow C_*(X)$ (singular chain complex)

$$X \xrightarrow{f} Y \text{ cts} \rightsquigarrow f_*: C_*(X) \rightarrow C_*(Y)$$

$f, g: X \rightarrow Y$ homotopic $\rightsquigarrow f_*$ and g_* are homotopic

Today we will prove analogous assertions in homological algebra.

(16.1) Again, let R be a commutative ring with $1 \neq 0$. Let $\mathcal{A} = R\text{-mod}$. Recall $Q \in \mathcal{A}$ (resp. $P \in \mathcal{A}$) is called injective (resp. projective) if any of the following equivalent conditions hold:

(1) $\text{Hom}_{\mathcal{A}}(-, Q)$ is exact (resp. $\text{Hom}_{\mathcal{A}}(P, -)$ is exact).

(see Defn. 10.3 page 5)

(2) \forall injective morphism $Q \hookrightarrow \tilde{Q}$ (resp. surjective morphism $\tilde{P} \twoheadrightarrow P$), we have $\tilde{Q} = Q \oplus \bar{Q}$ (resp. $\tilde{P} = P \oplus \bar{P}$).

(3) \forall injective morphism $0 \rightarrow A \xrightarrow{i} B$ and arbitrary $A \xrightarrow{f} Q$, $\exists \tilde{f}: B \rightarrow Q$ s.t. $\tilde{f} \circ i = f$

(resp. \forall surjective morphism $B \xrightarrow{\pi} C \rightarrow 0$ and any $g: P \rightarrow C$, $\exists \tilde{g}: P \rightarrow B$ such that $\pi \tilde{g} = g$)

(2)

[see § 11.4, 11.5 (page 8) and Theorem 12.1 page 1].

(16.2) Definition. — An injective resolution of $M \in \mathcal{A}$ is a cochain complex $0 \rightarrow I^0 \xrightarrow{d^0} I^1 \xrightarrow{d^1} I^2 \rightarrow \dots$ ($\in \mathcal{K}(\mathcal{A})$)

s.t. $\left\{ \begin{array}{l} \cdot H^n(I^\bullet) = 0 \quad \forall n \geq 1 \quad (\text{i.e. it is exact at } I^n) \\ \cdot I^k \text{ is injective} \quad \forall k \geq 0 \\ \cdot \text{Ker}(d^0) = M = H^0(I^\bullet) \end{array} \right. \quad \forall n \geq 1$

(in other words, $0 \rightarrow M \xrightarrow{i} I^0 \rightarrow I^1 \rightarrow \dots$ is exact)

A projective (resp. other adjectives like flat, free etc.) resolution of $N \in \mathcal{A}$ is a chain complex ($\in \mathcal{K}(\mathcal{A})$)

$$\dots \rightarrow P_n \xrightarrow{d_n} P_{n-1} \rightarrow \dots \rightarrow P_1 \xrightarrow{d_1} P_0 \rightarrow 0$$

s.t. $\left\{ \begin{array}{l} \cdot H_n(P_\bullet) = 0 \quad \forall n \geq 1 \quad (\text{i.e. exact at } P_n \quad \forall n \geq 1) \\ \cdot P_k \text{ is projective} \quad \forall k \geq 0 \\ \cdot H_0(P_\bullet) = P_0 / \text{Im}(d_1) = N \end{array} \right.$

(i.e. $\dots \rightarrow P_1 \rightarrow P_0 \xrightarrow{\pi} N \rightarrow 0$ is exact)

Remarks. - (i) Injective / projective resolutions are ③
 not unique, but only up to homotopy (Today's main thm.)

(ii) In practice, projective resolutions (or free resolutions - i.e. each P_k is a free R -module, i.e. $P_k = R^{\oplus(I)} = \bigoplus_{i \in I} R$) are computationally friendlier than injective resolutions.

(iii) For example, (proof next week)

$0 \rightarrow \mathbb{Z} \xrightarrow{i} \mathbb{Q} \xrightarrow{d^0} \mathbb{Q}/\mathbb{Z} \rightarrow 0$ is an injective resolution of \mathbb{Z} in $\mathcal{A} = \mathbb{Z}\text{-mod} = \text{Ab}$.

(iv) Let $R = K[x, y]$ and let $N = (x, y) \subset R$ ^{ideal}.

$$\begin{array}{ccccccc}
 & & 1 & \xrightarrow{\quad} & (y, -x) & & \\
 0 & \rightarrow & R & \xrightarrow{d_0} & R \oplus R & \xrightarrow{\pi} & N \rightarrow 0 \\
 & & \uparrow & & \begin{array}{l} (1, 0) \xrightarrow{\quad} x \\ (0, 1) \xrightarrow{\quad} y \end{array} & & \\
 \text{is a free resolution} & & & & \uparrow & & \\
 \text{"relations"} & & & & \text{"generators"} & &
 \end{array}$$

(16.3) Theorem. - Let $M, M' \in \mathcal{A}$. Assume we are given two ~~exact sequences~~ injective resolutions.

$$0 \rightarrow M \xrightarrow{i} I^0 \xrightarrow{d^0} I^1 \xrightarrow{d^1} I^2 \xrightarrow{d^2} \dots$$

$$0 \rightarrow M' \xrightarrow{j} J^0 \xrightarrow{d^0} J^1 \xrightarrow{d^1} J^2 \rightarrow \dots$$

Then for every R -linear map $f: M \rightarrow M'$, there exists (4)

$f^*: I^* \rightarrow J^*$, unique up to homotopy, s.t. $f^* \circ i = j \circ f$.

$$\begin{array}{ccccccc}
 0 & \longrightarrow & M & \xrightarrow{i} & I^0 & \longrightarrow & I^1 & \longrightarrow & \dots \\
 & & \downarrow f & \curvearrowright & \downarrow f^0 & & \downarrow f^1 & & \dots \\
 0 & \longrightarrow & M' & \xrightarrow{j} & J^0 & \longrightarrow & J^1 & \longrightarrow & \dots
 \end{array}$$

Similarly if we are given two projective resolutions (of $N \in \mathcal{A}$)

P_* and P'_* , then $\forall f: N \rightarrow N', \exists$ unique

up to homotopy $g: P_* \rightarrow P'_*$ s.t.

$$\begin{array}{ccccccc}
 \dots & \longrightarrow & P_2 & \longrightarrow & P_1 & \longrightarrow & P_0 & \longrightarrow & N & \longrightarrow & 0 \\
 & & \downarrow g_2 & & \downarrow g_1 & & \downarrow g_0 & \curvearrowright & \downarrow g & & \\
 \dots & \longrightarrow & P'_2 & \longrightarrow & P'_1 & \longrightarrow & P'_0 & \longrightarrow & N' & \longrightarrow & 0
 \end{array}$$

(16.4) As a consequence of Theorem 16.3 above, we get that:

Corollary. Injective resolutions are unique up to homotopy.
(Projective)

Pf. Let I_1^* and I_2^* be two injective resolutions of $M \in \mathcal{A}$.
Take $f = \text{Id}_M$ in the theorem to get (unique up to homotopy)

$\alpha: I_1^* \rightarrow I_2^*$ and $\beta: I_2^* \rightarrow I_1^*$. Then $\beta\alpha: I_1^* \rightarrow I_1^*$

Lifts the identity map $M \xrightarrow{\text{id}_M} M$. That is,

$$\begin{array}{ccccccc}
 0 & \longrightarrow & M & \longrightarrow & I_1^0 & \longrightarrow & I_1^1 & \longrightarrow & I_1^2 & \longrightarrow & \dots \\
 & & \parallel & & (\beta\alpha)^0 \downarrow & & (\beta\alpha)^1 \downarrow & & (\beta\alpha)^2 \downarrow & & \dots \\
 0 & \longrightarrow & M & \longrightarrow & I_1^0 & \longrightarrow & I_1^1 & \longrightarrow & I_1^2 & \longrightarrow & \dots
 \end{array}$$

But so does Id_{I_1} . By uniqueness up to homotopy part of the theorem $\beta\alpha \sim \text{Id}_{I_1}$. Similarly $\alpha\beta \sim \text{Id}_{I_2}$. Thus I_1^i and I_2^i are homotopic. The same argument works for projective resolutions. \square

(16.5) Proof of Theorem 16.3 - I. - Existence of f^\bullet .

Base Case. - f^0 .

$$\begin{array}{ccc}
 0 & \longrightarrow & M & \xrightarrow{\text{injective}} & I^0 \\
 & & \downarrow f^0 & \swarrow \text{exists, since } J^0 \text{ is injective R-mod.} & \\
 & & J^0 & & =: f^0
 \end{array}$$

Now assume f^0, f^1, \dots, f^n have been constructed so that
 (set $f^{-1} = f : M \rightarrow M'$)

$$\begin{array}{ccccccc}
 0 & \longrightarrow & M & \xrightarrow{i} & I^0 & \longrightarrow & \dots & \longrightarrow & I^n \\
 & & \downarrow f & & \downarrow f^0 & \downarrow f^1 & \dots & \downarrow f^{n-1} & \downarrow f^n \\
 0 & \longrightarrow & M' & \xrightarrow{j} & J^0 & \longrightarrow & \dots & \longrightarrow & J^n
 \end{array}$$

on all squares commute.

Define $I^n / \text{Ker}(d_I^n) \longrightarrow J^{n+1}$ via :

$$\begin{array}{ccccccc}
 I^{n-1} & \xrightarrow{d^{n-1}} & I^n & \longrightarrow & I^n / \text{Im}(d^{n-1}) & \xrightarrow{\text{defn. of } d_J^n \circ \bar{f}^n} & \\
 \downarrow f^{n-1} & & \downarrow f^n & & \downarrow \bar{f}^n & & \\
 J^{n-1} & \xrightarrow{d^{n-1}} & J^n & \longrightarrow & J^n / \text{Im}(d^{n-1}) & \xrightarrow{d_J^n} & J^{n+1}
 \end{array}$$

(factors through $J^n / \text{Im}(d^{n-1})$ because $\text{Im}(d^{n-1}) = \text{Ker}(d^n)$ in fact equal)

Now $I^n / \text{Im}(d^{n-1}) = I^n / \text{Ker}(d^n) \hookrightarrow I^{n+1}$

$$\begin{array}{ccc}
 d_J^n \circ \bar{f}^n & \downarrow & \\
 & & J^{n+1}
 \end{array}$$

← defn of f^{n+1} exists because J^{n+1} is injective.

Remark. — We didn't use injectivity of objects I^k ($k \geq 0$).
 We didn't use exactness of J , but only that it is a complex.

(16.6) Proof of Thm 16.3 part II. — Uniqueness.

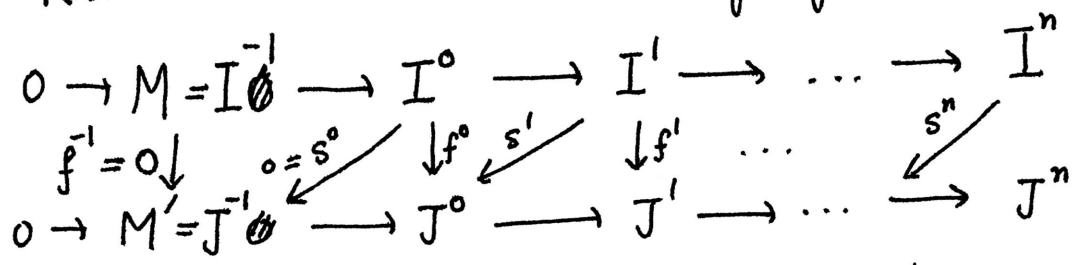
If f' and g' are two morphisms that lift the same $f: M \rightarrow M'$, their difference lifts $0: M \rightarrow M'$. Thus, it suffices to prove that if f' lifts zero morphism, then it is

homotopic to 0 (i.e. null homotopic). Thus we need

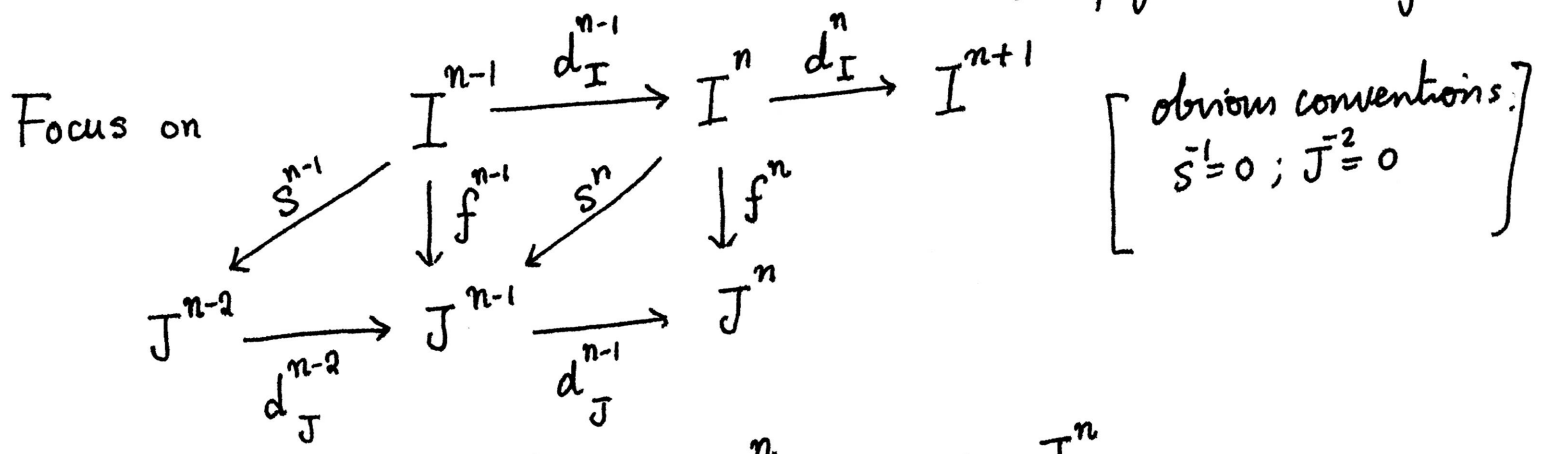
to construct $s^n: I^n \rightarrow J^{n-1}$ st. $f^n = ds + sd$.

Base case. $s^0: I^0 \rightarrow J^{-1} = 0$, ~~so~~ $s^0 := 0$.

Now assume that we successfully constructed s^0, \dots, s^n ($n \geq 0$)

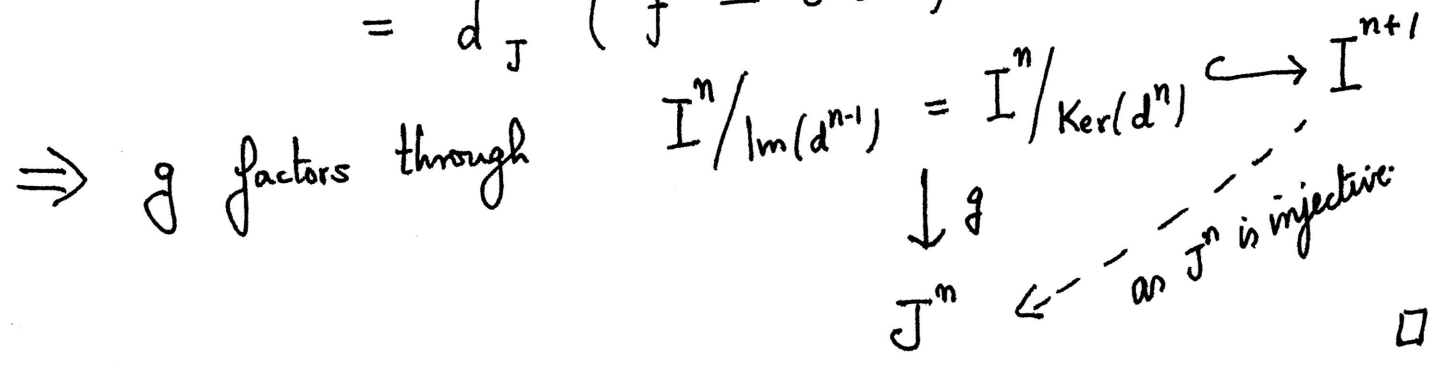


So that $f^k = d_J^{k-1} s^k + s^{k+1} d_I^k \quad \forall 0 \leq k \leq n-1$
 (empty condition if $n=0$)



Set $g = f^n - d_J^{n-1} s^n: I^n \rightarrow J^n$

Now $g \circ d_I^{n-1} = f^n d_I^{n-1} - d_J^{n-1} s^n d_I^{n-1}$
 $= d_J^{n-1} (f^{n-1} - s^n d_I^{n-1}) = d_J^{n-1} d_I^{n-2} s^{n-1} = 0$



□