

Recap: $\left[\begin{array}{l} A: \text{Cartan matrix}; D: \text{symmetrizing integers.} \\ A \in M_{I \times I}(\mathbb{Z}) \text{ s.t. } (D = \text{diag.}(d_i: i \in I) \in M_{I \times I}(\mathbb{Z})) \quad (d_i \in \mathbb{Z}_{>0}) \\ \text{(i) } a_{ii} = 2 \quad \forall i \in I \quad \text{(iii) } d_i a_{ij} = d_j a_{ji} \quad \forall i, j \in I \\ \text{(ii) } a_{ij} \leq 0 \quad \forall i \neq j \quad \text{(iv) } A \text{ is positive-definite.} \\ \text{(normalization: } \gcd(d_i: i \in I) = 1. \end{array} \right]$

To the data of (D, A) we associate: $E = \mathbb{R}\text{-span of } \{\rho_i: i \in I\}$
 $\alpha_j \in E^*$ defined by $\alpha_j(\rho_i) = a_{ij}$
 $\forall i, j \in I.$

$$\phi: E^* \xrightarrow{\sim} E$$

$$\alpha_i \mapsto d_i \rho_i$$

(\cdot, \cdot) on E^* :

$$(\alpha_i, \alpha_j) = d_i a_{ij}.$$

(\cdot, \cdot) on E - so ϕ is an isometry -

$$(\rho_i, \rho_j) = a_{ij} d_j^{-1}.$$

Simple reflections: $s_i(\gamma) = \gamma - \gamma(\rho_i) \alpha_i \quad (\forall i \in I.) \quad (\gamma \in E^*)$
 orthogonal reflection $(s_i(\alpha_j) = \alpha_j - a_{ij} \alpha_i \quad \forall i, j \in I)$

Transported via ϕ : $s_i(x) = x - \alpha_i(x) \rho_i \quad (\text{reflection in } \text{Ker}(\alpha_i) \subset E).$

$W =$ subgroup of $GL(E)$ (or $GL(E^*)$ via ϕ) generated by

$$s_i \quad (i \in I).$$

$R =$ W -orbit of "simple roots" $\{\alpha_i: i \in I\}$

$$= \bigcup_{i \in I} W \cdot \alpha_i \subset E^*.$$

Last time we proved that $|R| < \infty$ and $|W| < \infty$.

We also computed all rank 2 examples:

	$A_1 \times A_1$	A_2	B_2	G_2
R	$\{\pm \alpha_1, \pm \alpha_2\}$	$\pm \{\alpha_1, \alpha_2, \alpha_1 + \alpha_2\}$	$\pm \left\{ \begin{matrix} \alpha_1, \alpha_2, \alpha_1 + \alpha_2, \\ 2\alpha_1 + \alpha_2 \end{matrix} \right\}$	$\pm \left\{ \begin{matrix} \alpha_1, \alpha_2, \alpha_1 + \alpha_2, \\ 2\alpha_1 + \alpha_2, 3\alpha_1 + \alpha_2, \\ 3\alpha_1 + 2\alpha_2 \end{matrix} \right\}$
W	$\langle s_1, s_2 \mid s_1^2 = s_2^2 = e \rangle$	$s_1 s_2 = s_2 s_1$	$s_1 s_2 s_1 = s_2 s_1 s_2$	$s_1 s_2 s_1 s_2 s_1 s_2 = s_2 s_1 s_2 s_1 s_2 s_1$

"braid relⁿs"

§1. Sign-coherence Let \tilde{W} = group gen. by σ_i ($i \in I$) subject to relⁿs $\sigma_i^2 = e$ & $(\sigma_i \sigma_j)^{m_{ij}} = e$

where $m_{ij} = 2, 3, 4$ or 6 ; if $a_{ij} a_{ji} = 0, 1, 2$ or 3 .

Note: $\tilde{W} \twoheadrightarrow W$. We will prove that $\tilde{W} \cong W$ via this identification. (see §4 - pages 7-8 below)
 $\sigma_i \mapsto s_i$

Length function. - For any group G generated by "reflections" - elements of order 2 - $\{g_i\}_{i \in I}$ - define $l: G \rightarrow \mathbb{Z}_{\geq 0}$ by $l(g) = \text{Min} \{k \mid g = g_{i_1} \cdots g_{i_k} \text{ for some } i_1, \dots, i_k \in I\}$.
 (length of g)

Lemma (easy - proof left as an exercise): $l(g \cdot g_i) = l(g) \pm 1$.
 $\forall g \in G, i \in I$.

$$\text{Let } R_+ = R \cap \sum_{i \in I} \mathbb{Z}_{\geq 0} \alpha_i \quad \text{and} \quad R_- = R \cap \sum_{i \in I} \mathbb{Z}_{\leq 0} \alpha_i. \quad (3)$$

Prop. (Sign coherence of roots): -

For every $w \in W$, $i \in I$; if $l(ws_i) > l(w)$, then $w(\alpha_i) \in R_+$.

(if $l(ws_i) < l(w)$, then $w(\alpha_i) \in R_-$).

Hence, $R = R_+ \sqcup R_-$ ($R_- = -R_+$).

Proof. - We will show, by induction on $l(w)$, that:

$$ws_i > w \Rightarrow w(\alpha_i) \in R_+.$$

Base case: $l(w) = 0$ means $w = e$ & the statement is true.

Assume $n = l(w) \geq 1$. Pick a reduced expression $w = s_{i_1} \dots s_{i_n}$
(expression of length $= l(w)$)

For convenience, write $j = i_n$. Note: as $ws_i > w$, $i \neq j$.

We choose the smallest expression of w where s_i, s_j are to the right -

- precisely: let $W_{(i,j)} = \langle s_i, s_j \rangle$ ~~($\neq \langle s_i, s_j \rangle$)~~

pick v_0 to be of smallest length among

$$P = \left\{ u \in W \mid \begin{array}{l} w = u \cdot r \text{ where } r \in W_{(i,j)} \\ \& l(w) = l(u) + l(r) \end{array} \right\}$$

Obs: - • If $u \cdot r$ ^{($r \in W_{(i,j)}$)} is a reduced exp of w ; r cannot end with s_i
(since $ws_i > w$ is given)

$$\bullet w = (s_{i_1} \dots s_{i_{n-1}}) \underset{s_j}{\overset{\parallel}{s_{i_n}}} \Rightarrow l(v_0) \leq l(w) - 1$$

• $v_0 s_i > v_0$ and $v_0 s_j > v_0$. Since, otherwise

$$\begin{aligned}
 w &= v_0 \cdot r_0 = v_0 s_i \cdot s_i r_0 && \text{contradict the} \\
 &= v_0 s_j \cdot s_j r_0 && \text{minimality of } v_0.
 \end{aligned}$$

By induction, we conclude that $v_0(\alpha_i) \in R_+$ and $v_0(\alpha_j) \in R_+$.

Moreover $w = v_0 \cdot r_0$ where $r_0 = \dots s_i s_j$ reduced word in rank 2.

So $r_0(\alpha_i) \in \mathbb{Z}_{\geq 0} \alpha_i + \mathbb{Z}_{\geq 0} \alpha_j$
from rank 2 computation.

& cannot be written as $\dots s_j s_i$ (as $ws_i > w$)

For $R_- = -R_+$, note that $S_{w \cdot \alpha_i} = w s_i w^{-1} \quad \forall w \in W; i \in I$.

hence, if $\alpha = w(\alpha_i) \in R_-; S_\alpha \in W$ and

$$S_\alpha(\alpha) = -\alpha \in R_-.$$

Cor. - $S_i(R_+) \subset R_+ \cup \{-\alpha_i\}$ (i.e. $\forall \gamma \in R_+, \text{ either } S_i \gamma \in R_+ \text{ or } \gamma = \alpha_i$). □

§2. Properties of Weyl groups.

- Geometric action. hyperplane arrangement: $\left\{ H_\alpha = \text{Ker}(\alpha) \subset E \right\}_{\alpha \in R}$

(Note: $H_\alpha = H_{-\alpha}$).

$$\text{let } E^0 = E \setminus \bigcup_{\alpha \in R} H_\alpha$$

- a disconnected space

$\pi_0(E^0) :=$ set of connected components of E^0 .

(5)

As $w \cdot H_\alpha = H_{w(\alpha)}$; W acts on E° and permutes the connected components.

Cor of Sign-Coherence: $C_0 := \{x \in E : \alpha_i(x) > 0 \forall i \in I\}$ (Fundamental Chamber)
is a connected component of E° .

Prop. (i) For every $y \in E^\circ$, $\exists w \in W$ s.t. $w(y) \in C_0$.

(ii) If $a \in C_0$ and $w(a) = b \in C_0$, then $w = e$ and $a = b$.

[Hence $W \rightarrow \pi_0(E^\circ)$ is a bijection.] and $\bar{C}_0 =$ fundamental domain for $W \curvearrowright E$
 $w \mapsto w(C_0)$

Proof. (i) Choose $a \in C_0$ and pick $y_0 \in \underbrace{W \cdot y}_{\text{finite set}}$ closest to a .

Claim: $y_0 \in C_0$. If not, then there exists $i \in I$ s.t. $\alpha_i(y_0) < 0$.

We can check that this implies $\text{dist}(a, s_i(y_0)) < \text{dist}(a, y_0)$
Contradicting the choice of y_0 .

As $s_i(y_0) = y_0 - \alpha_i(y_0) h_i$; we get: $(a - s_i(y_0), a - s_i(y_0))$
and $(s_i(y_0), s_i(y_0)) = (y_0, y_0)$ $= (a, a) + (y_0, y_0)$
 $- 2(a, y_0 - \alpha_i(y_0) h_i)$
 $= (a - y_0, a - y_0) + 2\alpha_i(y_0)(a, h_i)$

As $(a, h_i) = \frac{1}{d_i} \alpha_i(a) > 0$, and $\alpha_i(y_0) < 0$; we get
 $\|a - s_i(y_0)\|^2 < \|a - y_0\|^2$

(ii) We will argue by induction on $l(w)$ again; to show that

$$\begin{aligned} a \in \overline{C_0} \\ b = w(a) \in \overline{C_0} \end{aligned} \Rightarrow w = e. \text{ and hence } a = b.$$

If $l(w) = l \geq 1$, write $w = s_{i_1} \dots s_{i_l}$ (a reduced expression)
 $= v \cdot s_{i_l}$ ($v = s_{i_1} \dots s_{i_{l-1}}$)

Note: $l(ws_{i_l}) < l(w) \Rightarrow w(\alpha_{i_l}) \in R_-$ (by Prop §1 above)

as $b \in \overline{C_0}$, $(w(\alpha_{i_l}))(b) \leq 0$.

But $(w(\alpha_{i_l}))(b) = (w(\alpha_{i_l}))(w(a)) = \alpha_{i_l}(a) \geq 0$.

$\Rightarrow \alpha_{i_l}(a) = 0$ and hence $s_{i_l}(a) = a$ proving that

$b = s_{i_1} \dots s_{i_{l-1}}(s_{i_l}(a)) = v(a)$ and $l(v) < l(w)$. \square

§3. Exchange property.

Theorem. - Let $w \in W$ and $i \in I$. Then the following are equivalent:

(a) $l(ws_i) < l(w)$ (b) $w(\alpha_i) \in R_-$

(c) For every expression $w = s_{i_1} \dots s_{i_n}$, $\exists j$ ($1 \leq j \leq n$).

s.t. $s_{i_{j+1}} \dots s_{i_n} s_i = s_{i_j} \dots s_{i_n}$ [Exchange property]

Proof. - (a) \Rightarrow (b) already proved in Prop §1.

(c) \Rightarrow (a) Pick a reduced exp. of $w = s_{i_1} \dots s_{i_r}$. By [Exchange Prop]

$w = s_{i_1} \dots s_{i_{j-1}} s_{i_{j+1}} \dots s_{i_r} s_i$ is another reduced exp. of w .

$\Rightarrow ws_i = s_{i_1} \dots s_{i_{j-1}} s_{i_{j+1}} \dots s_{i_r}$ has smaller length than w .

(b) ⇒ (c) : $W = s_{i_1} \dots s_{i_n} \quad w(\alpha_i) \in R_-$

Define $\beta_k = s_{i_k} \dots s_{i_n}(\alpha_i)$. $\beta_{n+1} = \alpha_i \in R_+$
 $\beta_1 = w(\alpha_i) \in R_-$

⇒ ∃! r s.t. $\beta_{r+1}, \beta_{r+2}, \dots, \beta_{n+1} \in R_+$ but $\beta_r \in R_-$.
(1 ≤ r ≤ n)

⇒ $\beta_r = s_{i_r}(\beta_{r+1}) \in R_- \Rightarrow \beta_{r+1} = \alpha_{i_r}$.
 $\beta_{r+1} \in R_+$ (see Cor on page 4)

That is, $s_{i_{r+1}} \dots s_{i_n}(\alpha_i) = \alpha_{i_r}$. [Recall $w(\alpha) = \beta \Rightarrow w s_\alpha w^{-1} = s_\beta$]

So $s_{i_{r+1}} \dots s_{i_n} s_{i_r} s_{i_n} \dots s_{i_{r+1}} = s_{i_r}$

⇒ $s_{i_{r+1}} \dots s_{i_n} s_{i_r} = s_{i_r} s_{i_{r+1}} \dots s_{i_n}$ as claimed. □

§4. Presentation of Weyl group. -

(Jacques Tits)

Theorem. - Let M be a semigroup / monoid and $\{g_i\}_{i \in I} \subset M$

be such that $\forall i \neq j$, $\underbrace{g_i g_j g_i \dots}_{m_{ij} \text{ terms}} = \underbrace{g_j g_i g_i \dots}_{m_{ij} \text{ terms}}$
[$m_{ij} \in \{2, 3, 4, 6\}$ - see page 2]

Let $w \in W$ and choose a reduced expression $w = s_{i_1} \dots s_{i_n}$.

Then $g(w) := g_{i_1} \dots g_{i_n}$ is independent of the choice of the reduced expression.

Cor. (see §1 - page 2) : $\tilde{W} \rightarrow W$ is an isomorphism.

(so W has a presentation: $\langle s_i (i \in I) \mid \begin{matrix} s_i^2 = (s_i s_j)^{m_{ij}} = e \\ (i) & (i \neq j) \end{matrix} \rangle$).

[Proof of the cor - $\tilde{W} = \langle \sigma_i (i \in I) \mid \begin{matrix} \sigma_i^2 = e \ \forall i & \& \\ (\sigma_i \sigma_j)^{m_{ij}} = e \ \forall i \neq j \end{matrix} \rangle$
 $\begin{matrix} \sigma_i \\ \downarrow \\ s_i \end{matrix} \quad \begin{matrix} \pi \\ \downarrow \\ W \end{matrix} \text{ (surjective)}$

Define $g: W \rightarrow \tilde{W}$ by $\begin{cases} g(s_i) = \sigma_i \neq \sigma_i \text{ ~~to } g \end{cases}~~$ and
By theorem above, g is well-defined. $\begin{cases} \text{for } w \in W; \text{ pick a red. exp.} \\ w = s_{i_1} \dots s_{i_n} \text{ and set} \\ g(w) = \sigma_{i_1} \dots \sigma_{i_n} \end{cases}$

Exercise. - Verify that g is a group hom. and a two-sided inverse to $\tilde{W} \xrightarrow{\pi} W$. \square

Proof of the theorem. - Let $R(w) =$ set of all red. exp.'s of w
 $= \{ (i_1, \dots, i_n) \in I^n : w = s_{i_1} \dots s_{i_n} \}$
(here $n = \ell(w)$).

To show - (using induction on length) : $\underline{i} = (i_1, \dots, i_n) \in R(w)$
 $\underline{j} = (j_1, \dots, j_n) \in R(w) \Rightarrow g(\underline{i}) = g(\underline{j})$

where $g(\underline{i}) = g_{i_1} \dots g_{i_n}$.

Note: The claim is true for $n=0$ and $n=1$ $R(s_i) = \{i\}$ is singleton.
($w=e$)

Assume $n \geq 2$. We will ~~show~~ construct new reduced expressions of w , using exchange property, as follows.

- use $w = s_{j_1} \cdots s_{j_n}$ to get $\underline{i}^{(1)} = (j_1, \dots, \overset{\text{skipped}}{\hat{j}_r}, \dots, j_n, i_n) \in R(w)$.
 $l(ws_{i_n}) < l(w)$ (Thm 3)

\rightarrow note $g(\underline{i}^{(1)}) = \underbrace{g_{j_1} \cdots g_{\hat{j}_r} \cdots g_{j_n}}_{\text{red. exp. of } ws_{i_n} \text{ of length } < n} \cdot g_{i_n}$

by induction $g_{j_1} \cdots g_{\hat{j}_r} \cdots g_{j_n} = g_{i_1} \cdots g_{i_{n-1}}$

$\Rightarrow g(\underline{i}^{(1)}) = g(\underline{i})$

\rightarrow If $r \neq 1$, then the same induction argument will show that $g(\underline{i}^{(1)}) = g(\underline{j})$ and we are done.

So, either we will be done by induction, or Exchange property always skips the 1st term

$\underline{i} = (i_1, \dots, i_n) \rightarrow \underline{i}^{(1)} = (j_2, \dots, j_{n-1}, i_n) \quad \underline{i}^{(2)} = (i_3, \dots, i_n, j_n, i_n)$
 $\underline{j} = (j_1, \dots, j_n) \rightarrow \underline{j}^{(1)} = (i_2, \dots, i_n, j_n) \quad \underline{j}^{(2)} = (j_3, \dots, j_n, i_n, j_n)$

\dots
 $\underline{i}^{(n)} = (\dots, j_n, i_n, j_n, i_n) \in R(w)$

$\underline{j}^{(n)} = (\dots, i_n, j_n, i_n, j_n) \in R(w)$

either we stop at some point - due to induction - or the process continues

But $g(\underline{i}^{(n)}) = g(\underline{j}^{(n)})$ by our hypothesis.

□