Warm-Up: For each sentence, draw a number line and indicate all x-values making the sentence true.

- (a) (x>2) $\Lambda (x^2>4)$
- (b) $(x>2) \vee (x^2>4)$
- $(c) \quad (x>2) \Rightarrow (x^2>4)$
- $(d) \quad (x>2) \iff (x^2>4)$

Last time: \ and \ \ \ and \ \ \ \.

Ex: Which statements are true?

- (1) (3 x & R) (x + 4 = 9) True: x = 5.
- (2) $(\exists x \in \mathbb{R})[(x+4=9) \land (x \neq 5)]$ False: $x+4=9 \implies x=9-4=5$
- ③ (∀x ∈ R) (x+4 =9)

 False: Try x=0.

$$(\forall x \in \mathbb{R})(x^2 + 6x + 8 \ge 0)$$

Can guess and check, or complete the square:

$$x^{2} + 6x + 8 = x^{2} + 6x + 9 - 1$$

= $(x + 3)^{2} - 1$,

False: Try x = -3.

6
$$(\forall x \in \mathbb{R})(x^2 + 6x + 10 \ge 0)$$

True: $x^2 + 6x + 10 = (x + 3)^2 + 1 \ge 1 > 0$
for all real numbers x.

Observe: · A single example proves a 3 statement.

- ·A single counterexample disproves a V statement.
- · To prove a V statement or disprove a 3 statement, we need an argument that works for all values.

Free + Bonnd Variables

Let P(x) = "x2 + 6x +8 > 0."

· Is P(x) true? It depends on x.

We say that x is a free variable in the Sentence P(x).

Think: The sentence P(x) is a function of x.

· Is (Vx&R) P(x) tme? No!

This sentence does NOT depend on x, because of the quantifier V.

In this case, we say x is a bound variable in the sentence (Vx (R)P(x).

The quantifier I can also bound variables: (Ix)P(x) does not depend on x.

Analogy: $f(x) = x^2$ vs. $\int_0^1 x^2 dx$

Note: When we use a quantifier (V or 3), a bound variable is ranging over a universe of possibilities.

Usnally, re should be explicit about this.

Common choices:

The universe matters!

 $(3x)(x^2-2)$

$$(\exists \times \in \mathbb{R})(x^2=2)$$

 $(\forall x \in \mathbb{R})(x^2 \geqslant 0)$

$$(\forall x \in \mathbb{C})(x^2 \ge 0)$$

Note: Over a finite set (universe),

• V is an "and" statement
• I is an "or" statement

Ex: If $A = \{-3, 1, 4\}$, then $(\forall x \in A)(x^2 < 20) = ((-3)^2 < 20) \land (1^2 < 20) \land (4^2 < 20)$ $(\exists x \in A)(x > 0) = (-3 > 0) \lor (1 > 0) \lor (4 > 0)$ (Both time)

For this reason, we can think of Y as "generalized and" and

I as "generalized or."

Thm (Generalized DeMorgan's Laws)

(a) $\neg [(\forall x \in A) P(x)] \equiv (\exists x \in A) (\neg P(x))$ (b) $\neg [(\exists x \in A) P(x)] \equiv (\forall x \in A) (\neg P(x))$ Proof: (a) Shopping $\neg [(\forall x \in A) P(x)]$ is time.

Proof: (a) Suppose $\neg [(\forall x \in A) P(x)]$ is time. Then $(\forall x \in A) P(x)$ is fulse. So there is some $x_0 \in A$ such that $P(x_0)$ is fulse, i.e. $\neg P(x_0)$ is time.

Hence $(\exists x \in A) (\neg P(x))$; true.

Conversely, suppose $(\exists x \in A) (\neg P(x))$ is time.

Then there is $x_o \in A$ such that $\neg P(x_o)$ is talse.

So (Vx EA) P(x) is fulse. Therefore, - (Vx EA) P(x) is time.

(b) is similar (see book).