

Greatest Common Divisors

Lemma: Let $a, b \in \mathbb{Z}$ not both zero.

There exists a unique $d \in \mathbb{N}$ such that

- ① $d|a$ and $d|b$ (d is a common divisor)
- ② For all $d' \in \mathbb{N}$, if $d'|a$ and $d'|b$, then $d' \leq d$.

We say d is the greatest common divisor of a and b , and write $d = \gcd(a, b)$.

Proof: Consider the set of all positive integers which are common divisors of a and b .

This set is non-empty (1 is a common divisor) and finite (every common divisor d satisfies $d \leq |a|$ or $d \leq |b|$), so it has a largest element. \blacksquare

Warm-Up: Compute

$$\gcd(10, 24)$$

$$\gcd(45, 15)$$

$$\gcd(1, 37)$$

$$\gcd(0, 37)$$

Note: The book uses a slightly different name (highest common factors) and definition.

Ex: If $a \in \mathbb{N}$, then $\gcd(a, 0) = a$.

Ex: Why do we not allow $a = b = 0$?
Every integer divides 0.

Lemma: Let $a, b \in \mathbb{Z}$ not both zero.

$$(a) \gcd(a, b) = \gcd(b, a)$$

$$(b) \gcd(a, b) = \gcd(a, -b)$$

Proof: (a) The definition is symmetric in a and b .

(b) Divisors of $-b$ are precisely divisors of b . □

How to compute $\gcd(a, b)$?

• If a and b are small, can list divisors.

• $\gcd(270, 192)$? Larger numbers?

The Euclidean Algorithm

Lemma: Let $a, b, q, r \in \mathbb{Z}$ such that

$$a = bq + r.$$

Then for all $d \in \mathbb{N}$, d is a common divisor of a and b if and only if d is a common divisor of b and r .

In particular, $\gcd(a, b) = \gcd(b, r)$.

Proof: HW 12.

Algorithm (Euclidean): **INPUT**: $a, b \in \mathbb{N}$ with $a \geq b$.
OUTPUT: $\gcd(a, b)$.

Set $r_{-1} = a$ and $n = 0$.
 $r_0 = b$

While $r_n \neq 0$:

- Divide r_{n-1} by r_n to get

$$r_{n-1} = r_n q_{n+1} + r_{n+1}$$

- If $r_{n+1} = 0$, output r_n and STOP.

- Else, increment $n \rightarrow n+1$.

Ex: $a=270$, $b=192$

$$\begin{pmatrix} r_1 = 270 \\ r_0 = 192 \end{pmatrix}$$

$$270 = 192(1) + 78$$

$$q_1 = 1, r_1 = 78$$

$$192 = 78(2) + 36$$

$$q_2 = 2, r_2 = 36$$

$$78 = 36(2) + 6$$

$$q_3 = 2, r_3 = 6$$

$$36 = 6(6) + 0$$

$$q_4 = 6, r_4 = 0$$

STOP and output 6.

So $\gcd(270, 192) = 6$.

Why does this work?

We must show

① The algorithm terminates.

② The output is correct.

Proof of termination: By the division algorithm,

$$\underline{r_{-1} \geq r_0} > r_1 > r_2 > \dots \geq 0$$

$a \geq b$ is given

Since the remainder decreases at every step, some remainder must eventually be zero. ✓

Proof of correctness: We have

$$r_{-1} = r_0 q_1 + r_1$$

$$r_0 = r_1 q_2 + r_2$$

⋮

$$r_{n-2} = r_{n-1} q_n + r_n$$

← Last non-zero remainder

$$r_{n-1} = r_n q_{n+1} + 0$$

So

$$\begin{aligned}\gcd(a, b) &= \gcd(r_{-1}, r_0) \\ &= \gcd(r_0, r_1) \\ &= \gcd(r_1, r_2) \\ &\vdots \\ &= \gcd(r_{n-1}, r_n) \\ &= \gcd(r_n, 0) = r_n\end{aligned}$$

by the Lemma. ✓

Soon, we'll prove

Thm: Let $a, b \in \mathbb{Z}$, not both zero.

Set $d = \gcd(a, b)$. Then there exist $x, y \in \mathbb{Z}$ such that

$$ax + by = d.$$

Ex: $a = 270$, $b = 192$ (so $d = 6$ by above)

From the Euclidean algorithm, we get

$$6 = 78 - 36(2)$$

$$\begin{aligned} &= 78 - [192 - 78(2)] \cdot 2 = 78(5) + 192(-2) \\ &= [270 - 192] \cdot 5 + 192(-2) \\ &= 270(5) + 192(-7). \end{aligned}$$

So $x = 5$, $y = -7$ solves

$$270x + 192y = 6.$$