The Enclidean Algorithm  
Lemma: Let a, b, q, r eZ such that  

$$a = bq + r$$
.  
Then for all de M, d is a common divisor of a and b  
if and only if d is a common divisor of b and r.  
In particular,  $gcd(a,b) = gcd(b,r)$ .  
Proof: HW 12.

Algorithm (Euclidean): INPUT: 
$$a, b \in N$$
 with  $a \ge b$ .  
OUTPUT:  $gcd(a, b)$ .  
Set  $r_{-1} = a$  and  $n = 0$ .  
 $r_0 = b$   
While  $r_n \neq 0$ :  
• Divide  $r_{n-1}$  by  $r_n$  to get  
 $r_{n-1} = r_n q_{n+1} + r_{n+1}$   
• If  $r_{n+1} = 0$ , ontput  $r_n$  and STOP.  
• Else, increment  $n \rightarrow n+1$ .

 $E_x: a = 270, b = 192$ 

$$270 = 192(1) + 78$$
  

$$192 = 78(2) + 36$$
  

$$78 = 36(2) + 6$$
  

$$36 = 6(6) + 0$$

 $(r_{-1} = 270)$   $r_{0} = 192)$   $q_{1} = 1, r_{1} = 78$   $q_{2} = 2, r_{2} = 36$   $q_{3} = 2, r_{3} = 6$  $q_{4} = 6, r_{4} = 0$ 

STOP and ontput 6. So gcd(270, 192) = 6.

Proof of termination: By the division algorithm, r-1 2 ro, > ri > rz > ··· > 0 a > b is given

$$\frac{Proof of correctness}{r_{-1} = r_{0}q_{1} + r_{1}}$$

$$r_{0} = r_{1}q_{2} + r_{2}$$

$$\vdots$$

$$r_{n-2} = r_{n-1}q_{n} + r_{n} + C$$

$$r_{n-1} = r_{n}q_{n+1} + O$$

$$gcd(a,b) = gcd(r_{-1}, r_{0})$$

$$= gcd(r_{0}, r_{1})$$

$$= gcd(r_{1}, r_{2})$$

$$\vdots$$

$$= gcd(r_{n-1}, r_{n})$$

$$= gcd(r_{n}, 0) = r_{n}$$
by the Lemma.

Soon, we'll prove  
Thm: Let 
$$a, b \in \mathbb{Z}$$
, not both zero.  
Set  $d = \gcd(a, b)$ . Then there exist  $x, y \in \mathbb{Z}$   
such that  
 $ax + by = d$ .

 $\checkmark$ 

Ex: 
$$a = 270$$
,  $b = 192$  (so  $d = 6$  by above)  
From He Euclidean algorithm, we get  
 $6 = 78 - 36(2)$   
 $= 78 - [192 - 78(2)] \cdot 2 = 78(5) + 192(-2)$   
 $= [270 - 192] \cdot 5 + 192(-2)$   
 $= 270(5) + 192(-7)$ .

So 
$$x=5$$
,  $y=-7$  solves  
270x + 192y = 6.