Warm-Up: Prove that n! >2" for all integers n>3 without induction.

Suppose not. Then there is a smallest integer a > 3 such that a! $\leq 2^{\alpha}$. (m/y?)

What can you say about a?

Thm: Let n & N. If n>1, then there is a prime p such that pln.

Proof: Suppose, to get a contradiction, that the theorem is filse.

That is, there is a natural number greater than I which is not divisible by any prime.

By the Well-Ordering Principle, those is a smallest such number. (My?)
Call it a.

So · a ? I · no prime divides a · If I< n < a, then n is dissible by Some prime.

Now, a is prime or composite.

· If a is prime, then ala, so a is divisible by a prime, which is a contradiction.

• If a is composite, then it has a positive divisor d with $d \neq 1$ and $d \neq a$. Since dla, $d \neq a$. But $d \neq a$, so d < a. Also, $d \neq 1$, so d > 1.

Thus, d must have a prime divisor p.

Since pld and dla, he have pla. (HW 10)

This is a contradiction.

Thus, the theorem holds for all natural numbers $n \ge 2$.

The infinitude of primes

Thm: There are infinitely many prime numbers.

Proof: Suppose, for the sake of contradiction, that there are only finitely many primes, say

P1, P2, ..., Pn.

Let m = p.pz...pn be the product of all of these primes.

Now, by the previous theorem, there is a prime 2 such that 21 (m+1).

Since q must be one of the primes $p_1, ..., p_n$ (because these are the only primes), so q_1m . Thus, q divides (m+1) - m = 1.

But this is a contradiction, since $q \ge 2$.