

Example: Checking for what p a quadratic form / \mathbb{Q} has a non-zero solution in \mathbb{Q}_p .

$$\text{Let } f(x, y, z) = x^2 - 13y^2 + 17z^2 \rightsquigarrow C_f \hookrightarrow \mathbb{P}^2$$

$C_f(\mathbb{R}) \neq \emptyset$: $[\sqrt{13}, 1, 0]$. For primes p , we'll use Hensel's Lemma + analysis over \mathbb{F}_p .

For $p \neq 2, 13, 17$, $C_{f \bmod p}$ is non-singular, so it suffices to find a primitive solution mod p (for $p \neq 2, 13, 17$) — then Hensel forces a p -adic solution.

We've already seen (pigeonhole) that these primitive solutions exist — so $C_f(\mathbb{Q}_p) \neq \emptyset$ for $p \neq 2, 13, 17$.

$p=13$: (note $[0, 1, 0]$ is a singular point of $C_{f \bmod 13}$)

$f \bmod 13 = x^2 + 4z^2$, so $[3, 0, 1]$ is a solution to $f \bmod 13$, and $\frac{\partial f}{\partial x}$ or $\frac{\partial f}{\partial z}$ at this point is $\neq 0$, so Hensel $\Rightarrow C_f(\mathbb{Q}_{13})$.

$p=17$: similarly $[8, 1, 0]$ lifts to a 17-adic solution

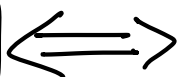
$p=2$: $f \bmod 8$ $f(2, 1, 1) \equiv 0 \bmod 8$ and

$v_2\left(\frac{\partial f}{\partial y}(2, 1, 1)\right) = 1$, so we win again. Conclusion:
 $\forall p, C_f(\mathbb{Q}_p) \neq \emptyset$ and $C_f(\mathbb{R}) \neq \emptyset$.

Question: For this f , is $C_f(\mathbb{Q}) \neq \emptyset$? The 3-variable case of the Hasse-Minkowski thm says YES:

Thm ^(Legendre) Let $f \in \mathbb{Q}[X_0, \dots, X_n]$ be a non-degenerate quadratic form. Then

$$\exists a \in \mathbb{Q}^{n+1} \setminus \{0\} : f(a) = 0$$



$$\forall p, \exists a_p \in \mathbb{Q}_p^{n+1} \setminus \{0\} : f(a_p) = 0$$

and

$$\exists a_\infty \in \mathbb{R}^{n+1} \setminus \{0\} : f(a_\infty) = 0$$

Alternatively, write $C_f \hookrightarrow \mathbb{P}^n$ for the associated nonsingular projective hypersurface. Then

$$C_f(\mathbb{Q}) \neq \emptyset \iff \forall p, C_f(\mathbb{Q}_p) \neq \emptyset \text{ and } C_f(\mathbb{R}) \neq \emptyset$$

(We are only going to prove the $n=2$ case.)

In particular, $\exists (a_0, a_1, a_2) \in \mathbb{Q}^3 \setminus \{0\} : a_0^2 - 13a_1^2 + 17a_2^2 = 0$,

Remark: ① The thm holds \forall number field (- -) this is harder.

② $n=1$ case is an exercise HW

③ We'll prove the $n=2$ case, which is elementary (but tricky)

④ $n=3$ (4 variables) is the hardest case; it uses quadratic reciprocity in a serious way. $n \geq 4$ is deduced by an inductive argument

see Serre, Course in Arithmetic

Now we'll prove the 3-variable thm.

The first key observation:

Lemma: Let k be any field w/ $\text{char}(k) \neq 2$. Let $a, b \in k^\times$.

Consider the field extension $k[\sqrt{b}] / k$ ($k[\sqrt{b}]$ is $\deg 2$
 $\iff b \notin (k^\times)^2$).

The equation $z^2 - ax^2 - by^2 = 0$ has a non-zero solution $(x, y, z) \in k^3 \setminus \{0\}$ if and only if

$$a \in N_{k[\sqrt{b}]/k} (k[\sqrt{b}]^\times) \subseteq k^\times.$$

Pf: Recall $N_{k[\sqrt{b}]/k}$ is the group homomorphism

$k[\sqrt{b}]^\times \rightarrow k^\times$ given by $N(\alpha) = \det(\alpha \text{ acting by multiplication on the } k\text{-vector space}$

Two cases: ① $b \in (k^\times)^2$: then any a is a norm,

and $(\sqrt{b})^2 - a \cdot 0^2 - b \cdot 1^2 = 0$ is a solution in $\mathbb{k}[\sqrt{b}]$

② $b \notin (\mathbb{k}^\times)^2$. Suppose $a \in N(\mathbb{k}[\sqrt{b}]^\times)$ $= \prod \sigma(\alpha)$
 $\sigma \in \text{Hom}_{\mathbb{k}}(\mathbb{k}[\sqrt{b}], \mathbb{k})$
 $= \begin{cases} \alpha & \text{if } \mathbb{k} = \mathbb{k}[\sqrt{b}] \\ a_1^2 - a_2^2 b & \text{if } \mathbb{k}[\sqrt{b}]/\mathbb{k} \\ & \text{is quadratic ext} \\ & \alpha = a_1 + a_2 \sqrt{b} \text{ for } a_i \in \mathbb{k}. \end{cases}$
 so $\exists a_1, a_2 \in \mathbb{k}: a = a_1^2 - b a_2^2$, so
 $(z, x, y) = (a_1, 1, a_2)$ is a solution.
 Conversely, suppose z, x, y is a solution
 in $\mathbb{k}^3 \setminus \{0\}$, $x \neq 0$, because otherwise $b \in (\mathbb{k}^\times)^2$,
 so $a = (\frac{z}{x})^2 - b(\frac{y}{x})^2 \in N(\mathbb{k}[\sqrt{b}]^\times)$. \square

Pf of H-M for 3 variables: Necessity of

having local solutions is clear. We'll show
 it is sufficient. Let $f \in \mathcal{O}[x_0, x_1, x_2]$
 be a non-degenerate quadratic form
 such that $C_f(\mathcal{O}_p) \neq \emptyset \forall p$ and $C_f(\mathbb{R}) \neq \emptyset$.

By linear COV, we may assume f is diagonal. We can rescale
 f so one coefficient is 1. Again by a linear COV we may
 assume $f(x_0, x_1, x_2) = x_0^2 - a x_1^2 - b x_2^2$ where $a, b \neq 0$
 are square-free integers, and $|a| \leq |b|$.

Now carry out proof by induction on $|a| + |b| = m$

Base case: $m=2$, so $|a|=|b|=1$, and $f = x_0^2 \pm x_1^2 \pm x_2^2$.

Any such f has a non-zero solution unless $f = x_0^2 + x_1^2 + x_2^2$,
 but this f has no solution $\in \mathbb{R}^3 \setminus \{0\}$. \checkmark

Suppose $m > 2$, hence $|b| > 2$. Write $b = \pm p_1 \cdots p_r$ for
distinct primes p_i .

Claim: a is a square mod b

Pf: By CRT, $\exists t, p \mid b, a \equiv t^2 \pmod{p}$. Clear if
 $p \mid a$, so assume $a \in \mathbb{Z}_p^\times$. By assumption, $\exists (z, x, y) \in \mathbb{Z}_p^3 \setminus$
 such that $z^2 - a x^2 - b y^2 = 0$, $(p \mathbb{Z}_p)^3$

As $p \mid b$, $z^2 \equiv ax^2 \pmod{p}$, so we win provided $x \not\equiv 0 \pmod{p}$

If $p \mid x$, then $p \mid z$, and $p^2 \mid by^2$, so since b is \square -free,
 $p \mid y \rightarrow$ contradicts primitivity of the solution.
 Thus, $a \equiv \square \pmod{p}$.

Since $a \equiv \square \pmod{b}$, $\exists t \in \mathbb{Z} : |t| \leq |b|/2$ such that
 $t^2 - a = b \cdot b'$ for some $b' \in \mathbb{Z}$.

Thus $bb' \in N_{\mathbb{Q}[\sqrt{a}]/\mathbb{Q}}(\mathbb{Q}[\sqrt{a}]^\times) \quad (N(t + \sqrt{a}))$.

for $\mathbb{Q}, \mathbb{Q}_p, \mathbb{R}$.

• For any such \mathbb{K} , f has a non-zero solution $\iff b \in N_{\mathbb{Q}[\sqrt{a}]/\mathbb{K}}(\mathbb{K}[\sqrt{a}]^\times)$ (Lemma)

$\iff b' \in N_{\mathbb{Q}[\sqrt{a}]/\mathbb{K}}(\mathbb{K}[\sqrt{a}]^\times)$
 (N is mult.)

$\iff x_0^2 - ax_1^2 - b'x_2^2 = f'(x_0, x_1, x_2)$ has a non-zero solution in \mathbb{K}^3

In particular, f' has local solutions in all \mathbb{Q}_p, \mathbb{R} .

But $|b'| = \left| \frac{t^2 - a}{b} \right| \leq \frac{|b|/2^2}{|b|} + \frac{|a|}{|b|} \leq \frac{1}{4} + 1 < |b|$.

b' may not be square-free, but no matter: (|b| > 2)

if $b' = d^2 \cdot b''$ ($d \in \mathbb{Z}, b'' \in \mathbb{Z}$), then f' represents 0

$\iff x_0^2 - ax_1^2 - b''x_2^2$ represents 0.

Certainly $|a| + |b''| < |a| + |b|$, so we win by induction. \blacksquare

The proof is effective: Suppose we want

$$\text{to solve } x^2 - 13y^2 + 17z^2 = 0$$

($a=13$, $b=-17$ square-free integers with $|a| \leq |b|$). The proof tells us:

• Solve $t^2 - a = b \cdot b'$ with $|t| \leq \frac{|b|}{2} = \frac{17}{2}$

$$\begin{array}{ccc} 8^2 - 13 = (-17) \cdot (-3) \\ \uparrow & & \uparrow \\ t & & b' \end{array}$$

$$\text{so } (-17) \cdot (-3) = N_{\mathbb{Q}(\sqrt{13})/\mathbb{Q}}(8 + \sqrt{13})$$

Now, $-17 \in \text{Norms} \Leftrightarrow -3 \in \text{Norms}$, so suffices to solve

$$x^2 + 3y^2 - 13z^2 = 0$$

(solving this will realize

$$-3 = N_{\mathbb{Q}(\sqrt{13})/\mathbb{Q}}(\alpha)$$

Etc — carrying this out is one of the PSet questions.

$$\text{so } -17 = N_{\mathbb{Q}(\sqrt{13})/\mathbb{Q}}\left(\frac{8 + \sqrt{13}}{2}\right)$$

• We now understand how to determine when a nonsingular conic $C_f \hookrightarrow \mathbb{P}^2$ over k ($f \in k[x_0, x_1, x_2]$ homogeneous degree 2) has $C_f(k) \neq \emptyset$ for $k = \mathbb{R}, \mathbb{F}_p, \mathbb{Q}_p, \mathbb{Q}$.

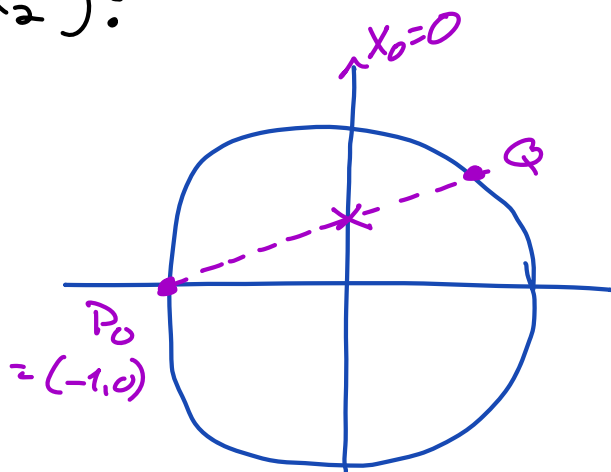
• What about describing the full set $C_f(k)$? Recall the parametrization of "Pythagorean triples" ($f = x_0^2 + x_1^2 - x_2^2$):

$$C_f(\mathbb{Q}) \ni P_0 = [-1, 0, 1].$$

try to extend this to a parametrization of the full projective curve C_f (or $C_f(\mathbb{Q})$)

Where should P_0 go?

P_0 should map to the intersection of the tangent line to C_f at P_0 with our fixed line ($x=0$).



$C_f \longrightarrow \{x_0=0\}$
 $(P_0 \neq Q) \longmapsto$ the point of intersection of $\{x_0=0\}$ with the line $\overline{P_0 Q}$.

$$\left(\frac{1-t^2}{1+t^2}, \frac{2t}{1+t^2} \right) \longmapsto (0, t)$$

Generalization (PSet):

Prop: Let k be any field, and let $f \in k[x_0, x_1, x_2]$ be a homogeneous polynomial of degree 2. Assume $C_f \hookrightarrow \mathbb{P}^2$ is non-singular, and that $C_f(k) \neq \emptyset$. Then $X \cong \mathbb{P}^1$ (isomorphism of algebraic varieties).

More precisely, fix $P_0 \in C_f(k)$, and fix a linear form $L(x_0, x_1, x_2) \in k[x_0, x_1, x_2]$ such that $C_L \not\ni P_0$. Then the projection map

$$\begin{aligned} \pi: C_f &\longrightarrow C_L \\ P_0 \notin Q &\longmapsto C_L \cap \overline{P_0 Q} \\ P_0 &\longmapsto C_L \cap \text{Tan}_{C_f, P_0} \end{aligned}$$

is an isomorphism of varieties over k .

In particular, \forall fields $k' \supset k$, π induces a bijection $C_f(k') \xrightarrow{\sim} C_L(k')$

Moreover $C_L \cong \mathbb{P}^1$, so $C_f \cong \mathbb{P}^1$

("any nonsingular conic/ k with a k -point is \cong to \mathbb{P}^1 ").

(PSet).

Do this exercise!

Beyond Conics

Recall our 4-step analysis for conics/ \mathbb{Q}

- ① Solve over \mathbb{F}_p (by a counting argument)
- ② Use ① + Hensel's Lemma to solve over \mathbb{Q}_p (and \mathbb{R} easy)
- ③ Apply Hasse-Minkowski theorem to determine whether solutions exist over \mathbb{Q}
- ④ If $C_f(\mathbb{Q}) \neq \emptyset$, use a linear parametrization ($C_f \simeq \mathbb{P}^1$) to describe $C_f(\mathbb{Q})$.

In higher degree, every step except ② becomes much more difficult.

Example: ① Consider $3X_0^3 + 4X_1^3 + 5X_2^3 = f(X_0, X_1, X_2)$
defining $C_f \hookrightarrow \mathbb{P}^2$. For $p \equiv 1 \pmod{3}$,

~~$$C_f(\mathbb{F}_p) = \bigcup_{i=0}^2 (C_f(\mathbb{F}_p) \cap \{X_i \neq 0\}) = \{3x^3 + 4y^3 + 5z^3 = 0\} \cup \{3x^3 + 4 + 5z^3 = 0\} \cup \{3 + 4y^3 + 5z^3 = 0\}$$~~

An equation like $3x^3 = 5 - 4y^3$ does not have solutions by a simple counting argument, since $|\mathbb{F}_p^\times|^3 = \frac{p-1}{3}$.

There is in this case an elementary argument to show \exists solution, but

- how many solutions?
- higher degree, more complicated polys?

③ The analogue of the H-M theorem fails ("most of the time," in fact!). We'll do an example next.

④ Even if we know $C_f(\mathbb{Q}) \neq \emptyset$, describing $C_f(\mathbb{Q})$ is extremely difficult and even when $\deg(f) = 3$ the subject of deep open problems. We'll spend much of the course starting to study this $\deg(f) = 3$ case.

Goal:

Failure of the local-global principle for genus 1 curves.

Example: $f(x_0, x_1, x_2) = 3x_0^3 + 4x_1^3 + 5x_2^3 \in \mathbb{Q}(x_0, x_1, x_2)$

$C_f(\mathbb{R}) \neq \emptyset$, and $\forall p C_f(\mathbb{Q}_p) \neq \emptyset$, but $C_f(\mathbb{Q}) = \emptyset$

↑
checking these statements is a good exercise

↑
tricky. (see Keith Conrad's note on this for details).

We'll do the details of an easier example.

Let $p \equiv 1 \pmod{8}$ such that $\left(\frac{2}{p}\right)_4 \neq 1$ ($x^4 = 2 \pmod{p}$ not solvable in \mathbb{F}_p)
(eg $p = 17$)

Consider the affine curve $C^0 \hookrightarrow \mathbb{A}^2$ over \mathbb{Q} .

$$\{w^2 = 2 - 2pz^4\}.$$

C^0 is nonsingular (check this), but the projective closure $\overline{C^0} = \{W^2T^2 = 2T^4 - 2pZ^4\} \hookrightarrow \mathbb{P}^2$ is singular on $\overline{C^0} \cap \{T=0\} = \{[0, 1, 0]\}$.
(check)

Fact: Let C^0 be any nonsingular affine curve. Then there is a nonsingular projective curve (maybe in higher-dim \mathbb{P}^n) C such that $C^0 \cong C \setminus \{\text{finite set}\}$.

In our example, we can see this explicitly:

Define $C \hookrightarrow \mathbb{P}^3$ as the intersection of two quadrics

$$C^0 \xrightarrow{f} U_0 = \{x_0 \neq 0\} \xrightarrow{\cong} \mathbb{A}^3$$

$$x_0x_3 - x_1^2 = 0 \text{ and } x_2^2 - 2x_0^2 + 2px_3^2 = 0.$$

$$f: (w, z) \mapsto [1, z, w, z^2]$$

so we see $f: C^0 \rightarrow C \cap U_0$, and this map

has inverse $C \cap U_0 \rightarrow C^0$
 $= \{(y_1, y_2, y_3) \in \mathbb{A}^3 \mid y_3 = y_1^2, y_2^2 - 2 + 2py_3^2 = 0\}$

given by $g(y_1, y_2, y_3) = (y_2, y_1)$

So we find $C^0 \xrightarrow{\cong} C \cap U_0$ (iso of varieties).

Note: $C \setminus C^0 = C \cap \{x_0 = 0\} = \{[0, 0, \pm\sqrt{-2p}, 1]\}$

→ Check that C is nonsingular also at these points

Theorem: For this $C \hookrightarrow \mathbb{P}^3$,

$C(\mathbb{Q}_p) \neq \emptyset \forall p$, but $C(\mathbb{Q}) = \emptyset$.

$$C(\mathbb{R}) \neq \emptyset$$

Next time: we'll start by defining nonsingularity for general varieties, and then we'll prove the theorem. ^{very quickly} - see Gaurav's talk ^{video}

PSet post tomorrow am.