

4/14/08

- (1) **Simplicial Resolutions.** Let  $F: \mathbf{C} \rightarrow \mathbf{D}$  be a left adjoint to  $U: \mathbf{D} \rightarrow \mathbf{C}$ , and denote by  $\sigma$  (respectively  $\delta$ ) the unit (respectively counit) of the adjunction. So  $\sigma(d) \in \mathbf{D}(d, UF(d))$  corresponds to  $\text{id}_{F(d)}$  and  $\delta(c) \in \mathbf{C}(FU(c), c)$  to  $\text{id}_{U(c)}$ .
- (a) For any object  $c \in \mathbf{C}$ , show that the following formulas define a simplicial object  $\perp_\bullet(c)$  in  $\mathbf{C}$ : Let  $\perp_n(c) = (FU)^n(c)$  and let
- $$d_i = (FU)^i \delta((FU)^{n-i}(c)): (FU)^{n+1}(c) \longrightarrow (FU)^n(c)$$
- $$s_i = (FU)^i F\sigma(U(FU)^{n-i}(c)): (FU)^{n+2}(c) \longrightarrow (FU)^{n+1}(c).$$
- (b) Show that  $\delta(c)$  induces a simplicial map  $\epsilon(c): \perp_\bullet(c) \rightarrow \mathbf{c}_\bullet$ , where the right hand side denotes the constant simplicial  $\mathbf{C}$ -object.
- (c) Assume that there is a functor  $K: \mathbf{D} \rightarrow \underline{\text{Set}}$  such that  $KU(\perp_\bullet(c))$  is Kan. Show that the augmentation  $\epsilon(c)$  gives a weak equivalence  $KU(\epsilon(c))$ , i.e. it induces an isomorphism on all homotopy groups (that vanish above dimension zero for the constant functor).
- (d) Apply this to your favorite pair of adjoint functors and see what you get. For example, you could use  $\mathbf{C} = R\text{-Mod}$  or  $\text{Ring}$  and  $\mathbf{D} = \text{Ab}$  or  $\underline{\text{Set}}$ . If  $\mathbf{C}$  happens to be an abelian category, one can apply the (normalized) chain complex to  $\epsilon(c)$  and get all resolutions we have studied in class so far!

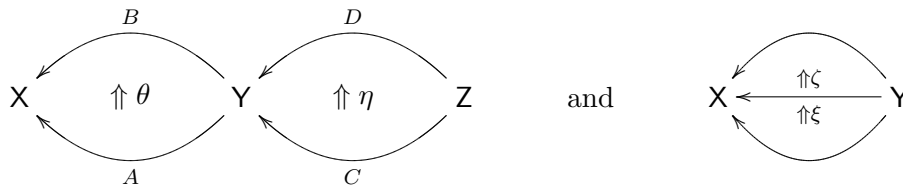
**Solution.**

(a) We will prove a more general result which will imply the desired result. Our proof will rely on the following facts:

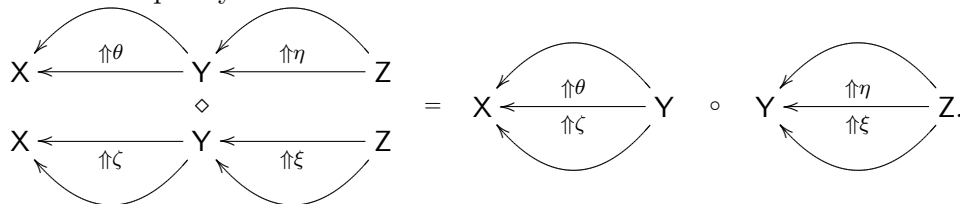
Facts: (i) If  $\mathbf{X}, \mathbf{Y}$ , and  $\mathbf{Z}$  are categories, we have an equivalence

$$\text{Fun}(\mathbf{X}, \text{Fun}(\mathbf{Y}, \mathbf{Z})) \cong \text{Fun}(\mathbf{X} \times \mathbf{Y}, \mathbf{Z}) \cong \text{Fun}(\mathbf{Y}, \text{Fun}(\mathbf{X}, \mathbf{Z})).$$

(ii) There is a 2-category  $\text{Cat}$  whose objects are small categories, 1-morphisms are functors, and 2-morphisms are natural transformations. Recall that there are two compositions  $\circ$  and  $\diamond$  of 2-morphisms in  $\text{Cat}$ , given respectively by:



for natural transformations  $\theta: A \Rightarrow B$ ,  $\eta: C \Rightarrow D$ , etc. Furthermore,  $\underline{\mathfrak{C}}$  is strict in the sense that we have equality



We will construct a simplicial functor  $\perp_\bullet \in \text{Fun}(s\Delta, \text{Fun}(\mathbf{C}, \mathbf{C})) \cong \text{Fun}(\mathbf{C}, \text{Fun}(s\Delta, \mathbf{C}))$ , i.e.,  $\perp_\bullet(c)$  will be a simplicial object in  $\mathbf{C}$ . To define this simplicial functor, we will look at a 2-subcategory  $\mathfrak{C}$  of

Cat whose objects are the categories  $\mathcal{C}$  and  $\mathcal{D}$ ; 1-morphisms are composites of  $F$ ,  $U$ ,  $\text{id}_{\mathcal{C}}$ , and  $\text{id}_{\mathcal{D}}$ ; and 2-morphisms are composites of identity natural transformations,  $\sigma: \text{id}_{\mathcal{D}} \Rightarrow UF$ , and  $\delta: FU \Rightarrow \text{id}_{\mathcal{C}}$ . We will use the following suggestive notation: to say a 1-morphism  $A$  is in  $\mathfrak{C}$ , we will write  $A \in \mathfrak{C}$ , and to say there is a natural transformation  $\eta: A \rightarrow B$  for  $A, B \in \mathfrak{C}$ , we will write  $\eta \in \mathfrak{C}(A, B)$ . Using the definitions of composition of 2-morphisms in  $\mathfrak{C}$ , one can easily prove the following lemma:

**Lemma 1:** (i) Consider  $(FU)^n \in \mathfrak{C}$ . Then

$$\begin{array}{c} \text{id}_{\mathcal{C}} \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{C} \xleftarrow{\quad} \mathcal{C} \\ \curvearrowleft \\ FU \end{array} \quad \begin{array}{c} (FU)^n \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{C} \xleftarrow{\quad} \mathcal{C} \\ \curvearrowleft \\ (FU)^n \end{array} = \begin{array}{c} (FU)^n \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{C} \xleftarrow{\quad} \mathcal{C} \\ \curvearrowleft \\ (FU)^{n+1} \end{array},$$

i.e.  $\delta \circ \text{id}_{(FU)^n} = \delta(FU)^n$ . Similarly,  $\text{id}_{(FU)^n} \circ \delta = (FU)^n \delta$ ,  $\sigma \circ \text{id}_{(UF)^n} = \sigma(UF)^n$ , and  $\text{id}_{(UF)^n} \circ \sigma = (UF)^n \sigma$ .

(ii) Consider  $FU \in \mathfrak{C}$ . Then

$$\begin{array}{c} \text{id}_{\mathcal{C}} \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{C} \xleftarrow{\quad} \mathcal{C} \\ \curvearrowleft \\ FU \end{array} = \begin{array}{c} \text{id}_{\mathcal{C}} \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{C} \\ \curvearrowleft \\ FU \end{array},$$

i.e.  $\delta \circ \text{id}_{FU} = \delta$ . Similarly,  $\sigma \circ \text{id}_{UF} = \sigma$ .

(iii)  $\text{id}_A \circ \text{id}_B = \text{id}_{AB}$  and  $\text{id}_A \diamond \text{id}_A = \text{id}_A$  for composable  $A, B \in \mathfrak{C}$ .

Also, since we have an adjunction, we have the following lemma:

**Lemma 2:** We have the following relations among  $F, U, \sigma, \delta, \text{id}_F, \text{id}_U$ :

(i)  $\delta F \diamond F \sigma = \text{id}_F$ :

$$\begin{array}{c} F \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{D} \xleftarrow{\quad} \mathcal{D} \\ \curvearrowleft \\ F \end{array} = \begin{array}{c} F \\ \curvearrowright \\ \mathcal{C} \xleftarrow{\quad} \mathcal{D} \\ \curvearrowleft \\ F \end{array},$$

(ii)  $U \delta \diamond \sigma U = \text{id}_U$ . The diagram is similar.

We now introduce a powerful tool: a graphical calculus for working in  $\mathfrak{C}$ . Usually,  $A \in \mathfrak{C}$  is written as an arrow from its source to its target. One could instead write  $A$  as a point and its source and target as arrows going in and out of  $A$ . This diagram is the dual diagram:

$$\mathcal{X} \xleftarrow{A} \mathcal{Y} = \xleftarrow{\mathcal{X}} A \xleftarrow{\mathcal{Y}} = \xleftarrow{\mathcal{X}} A \xrightarrow{\mathcal{Y}}$$

where one often suppresses the directions on the arrows when the convention is understood (all arrows will point left in these dual diagrams).

One usually writes natural transformations as 2-cells. Dually, we can write them as pictures from dual diagrams to dual diagrams. This gives the added benefit that we can denote an identity

2-morphisms as a “string” going from the bottom to the top:

$$\begin{array}{c}
 \begin{array}{ccc}
 & A & \\
 \swarrow & & \searrow \\
 X & \uparrow \text{id}_A & Y \\
 \nwarrow & & \nearrow \\
 & A & 
 \end{array}
 = 
 \begin{array}{c}
 \begin{array}{ccc}
 X & A & Y \\
 \hline
 & \uparrow \text{id}_A & \\
 \hline
 X & A & Y
 \end{array}
 \end{array}
 \end{array}$$

and we can denote the unit or counit by a “cap,” once more reading bottom to top:

$$\begin{array}{c}
 \begin{array}{ccc}
 & UF & \\
 \swarrow & & \searrow \\
 D & \uparrow \sigma & D \\
 \nwarrow & & \nearrow \\
 & \text{id}_D & 
 \end{array}
 = 
 \begin{array}{c}
 \begin{array}{ccccc}
 & D & & C & & F & & D \\
 & \swarrow & & \searrow & & \swarrow & & \searrow \\
 & U & & & & F & & \\
 & \searrow & & \swarrow & & & & \\
 & & & & & \text{id}_D & & 
 \end{array}
 \end{array}$$
  

$$\begin{array}{c}
 \begin{array}{ccc}
 & \text{id}_C & \\
 \swarrow & & \searrow \\
 C & \uparrow \delta & C \\
 \nwarrow & & \nearrow \\
 & FU & 
 \end{array}
 = 
 \begin{array}{c}
 \begin{array}{ccccc}
 & & & \text{id}_C & & \\
 & & & \swarrow & & \searrow \\
 & & & & & & & \\
 & & & & & F & & D & & U & & C \\
 & & & & & \swarrow & & \searrow & & & & \\
 & & & & & & & & & & & 
 \end{array}
 \end{array}$$

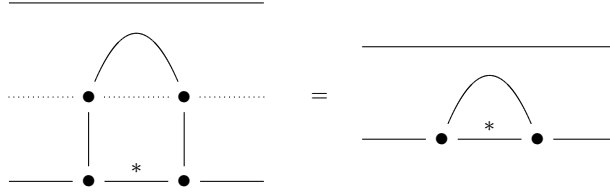
Now  $\circ$ -composition of 2-morphisms corresponds in the dual diagram language to splicing pictures together sideways, and  $\diamond$ -composition corresponds to stacking. Often we do not label the categories and 1-morphisms as the 2-morphisms encode this data.

$$\begin{array}{c}
 \begin{array}{ccc}
 & A & \\
 \swarrow & & \searrow \\
 X & \uparrow \text{id}_A & Y \\
 \nwarrow & & \nearrow \\
 & A & 
 \end{array}
 \begin{array}{ccc}
 & B & \\
 \swarrow & & \searrow \\
 Y & \uparrow \text{id}_B & Z \\
 \nwarrow & & \nearrow \\
 & B & 
 \end{array}
 = 
 \begin{array}{c}
 \begin{array}{ccccc}
 & \bullet & & \bullet & \\
 & \swarrow & & \searrow & \\
 & \text{id}_A & & \text{id}_B & \\
 & \swarrow & & \searrow & \\
 & \bullet & & \bullet & 
 \end{array}
 \end{array}
 \end{array}$$

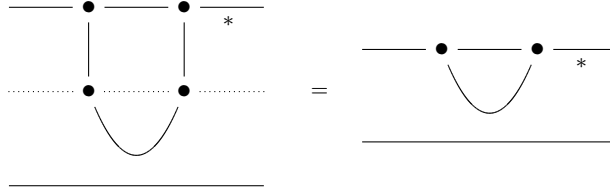
When a cap appears, we sometimes do not label the 2-morphisms as they are completely determined. Furthermore, we omit the bullet representing the identity functor:

$$\begin{array}{c}
 \begin{array}{ccc}
 & \text{id}_C & \\
 \swarrow & & \searrow \\
 C & \uparrow \delta & C \\
 \nwarrow & \uparrow \text{id}_{FU} & \nearrow \\
 & FU & 
 \end{array}
 = 
 \begin{array}{c}
 \begin{array}{ccccc}
 & & & & \\
 & & & \text{cap} & \\
 & & & \swarrow & \searrow \\
 & & & \bullet & \bullet \\
 & & & \swarrow & \searrow \\
 & & & \bullet & \bullet \\
 & & & \swarrow & \searrow \\
 & & & & 
 \end{array}
 \end{array}
 \end{array}$$

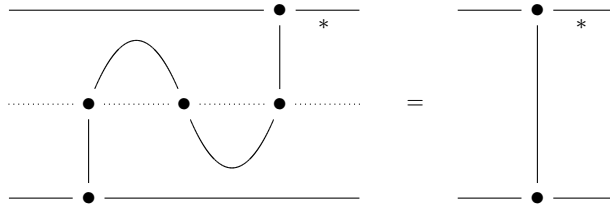
Using the convention of shading the regions of the diagram which have D along the outer boundary, we get string diagrams in which the lemmas above correspond exactly to isotopy of the strings. In the following diagrams, shaded regions have a \* in them:



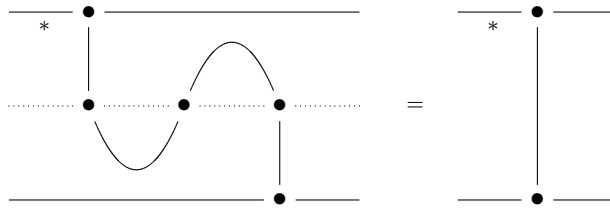
**Figure 1:** Lemma 1.ii,  $\delta \diamond \text{id}_{FU} = \delta$ .



**Figure 2:** Lemma 1.ii,  $\sigma \diamond \text{id}_{UF} = \sigma$ .



**Figure 3:** Lemma 2.i,  $\delta F \diamond F \sigma = \text{id}_F$ .



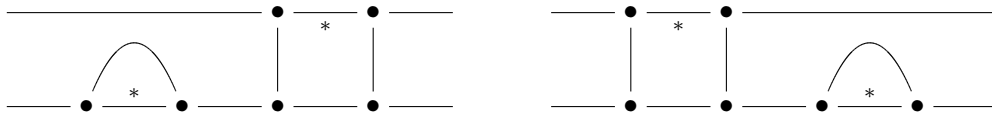
**Figure 4:** Lemma 2.ii,  $U \delta \diamond \sigma U = \text{id}_U$ .

Now that we have this graphical calculus at our disposal, we can build our simplicial functor  $\perp_\bullet$ . Set  $\perp_n = (FU)^{n+1} \in \mathfrak{C}$  and

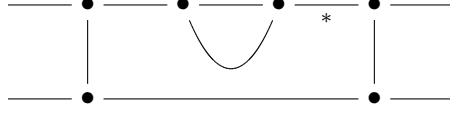
$$d_i = (FU)^i \delta (FU)^{n-i} : (FU)^{n+1} \longrightarrow (FU)^n$$

$$s_i = (FU)^i F \sigma U (FU)^{n-i} : (FU)^{n+2} \longrightarrow (FU)^{n+1}.$$

We can represent these natural transformations diagrammatically:



**Figure 5:**  $d_0, d_1 \in \mathfrak{C}((FU)^2, FU)$ .



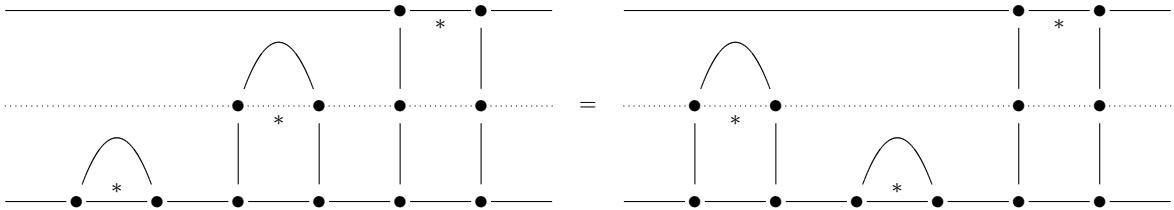
**Figure 6:**  $s_0 \in \mathfrak{C}(FU, (FU)^2)$ .

To prove that this defines a simplicial functor, we must show the following relations:

$$\begin{aligned}
 d_i d_j &= d_{j-1} d_i \text{ for } i < j \\
 s_i s_j &= s_{j+1} s_i \text{ for } i \leq j \\
 d_i s_j &= \begin{cases} s_j d_{i-1} & \text{if } i < j \\ \text{id} & \text{if } i = j, j + 1 \\ s_{j-1} d_i & \text{if } i > j + 1. \end{cases}
 \end{aligned}$$

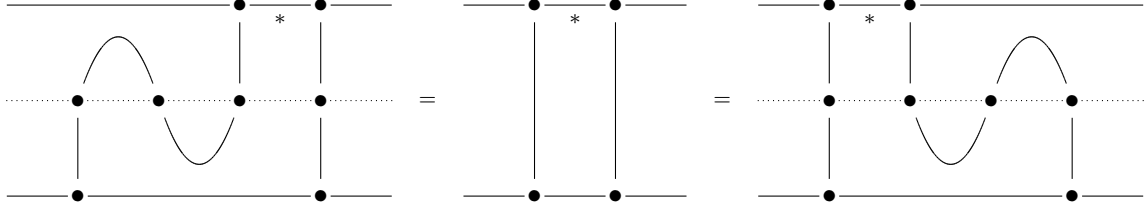
However, these are straightforward from the earlier lemmas by drawing the appropriate diagrams. We will prove  $d_0 d_1 = d_0 d_0 : (FU)^3 \rightarrow (FU)^1$  first using the usual diagrams in  $\mathfrak{C}$ , and secondly using the graphical calculus, i.e. the dual diagrams.

$$\begin{aligned}
 d_0 d_1 &= \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta FU \\ \uparrow FU \delta FU \\ \text{FU} \end{array} \mathfrak{C} = \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta \\ \uparrow \text{id}_{FU} \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \text{id}_{\text{id}_{\mathfrak{C}}} \\ \uparrow \delta \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \text{id}_{FU} \\ \uparrow \text{id}_{FU} \\ \text{FU} \end{array} \mathfrak{C} \\
 &= \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{FU} \\ \uparrow \text{id}_{FU} \\ \text{FU} \end{array} \mathfrak{C} \\
 &= \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \text{id}_{\text{id}_{\mathfrak{C}}} \\ \uparrow \delta \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta \\ \uparrow \text{id}_{FU} \\ \text{FU} \end{array} \mathfrak{C} \begin{array}{c} \text{FU} \\ \uparrow \text{id}_{FU} \\ \uparrow \text{id}_{FU} \\ \text{FU} \end{array} \mathfrak{C} \\
 &= \mathfrak{C} \begin{array}{c} \text{id}_{\mathfrak{C}} \\ \uparrow \delta FU \\ \uparrow \delta (FU)^2 \\ \text{FU} \end{array} \mathfrak{C} = d_0 d_0
 \end{aligned}$$



**Figure 7:**  $d_0d_1 = d_0d_0 \in \mathfrak{C}((FU)^3, FU)$ .

The only non-obvious relation is  $d_i s_i = id = d_{i+1} s_i$ , but these follow directly from Lemma 2. We will prove  $d_1 s_1 = id_{FU} = d_2 s_1 \in \mathfrak{C}((FU)^2, (FU)^2)$  using the dual diagrams.



**Figure 8:**  $d_0 s_0 = id_{FU} = d_1 s_0 \in \mathfrak{C}(FU, FU)$ .

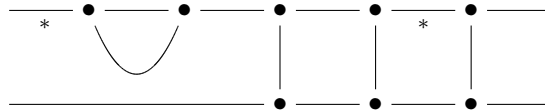
Thus,  $\perp_\bullet$  is a simplicial functor in  $\text{Fun}(\mathbf{C}, \mathbf{C})$ , and evaluation at  $c \in \mathbf{C}$  gives a simplicial object  $\perp_\bullet(c)$  in  $\mathbf{C}$ .

(b) Define  $\perp_{-1} = id_{\mathbf{C}}$  and  $\epsilon: (FU)^n \rightarrow id_{\mathbf{C}}$  by ‘‘capping everything off,’’ i.e.  $\delta(d_0)^{n-1}$ . Note that this is equivalent to  $\delta d_{i_1} \cdots d_{i_{n-1}}$  by the dual diagrams. Next, define  $id_\bullet$  as the constant simplicial functor, i.e.  $id_n = id_{\mathbf{C}}$  and  $d_i, s_j$  are all  $id_{id_{\mathbf{C}}}$ . It is obvious that  $\epsilon$  induces a simplicial map from  $\perp_\bullet \rightarrow id_\bullet$ , i.e.  $d_i \epsilon = \epsilon d_i$ , and  $s_j \epsilon = \epsilon s_j$ , since  $\epsilon$  is capping everything off. Once more, evaluation at  $c \in \mathbf{C}$  gives the simplicial map  $\perp_\bullet \rightarrow c_\bullet$ .

(c) Composition with  $U$  gives a simplicial functor  $U \perp_\bullet$  in  $\text{Fun}(\mathbf{C}, \mathbf{D})$ . The dual diagrams (for the  $d_i$ 's and  $s_j$ 's) are altered by adding one more string to the left. One immediately observes the existence of an extra degeneracy  $s_{-1}: U(FU)^n \rightarrow U(FU)^{n+1}$  by

$$s_{-1} = \sigma U(FU)^{n+1} \in \mathfrak{C}(U(FU)^{n+1}, U(FU)^{n+1}).$$

Dual diagrammatically, we have:



**Figure 9:**  $s_{-1} \in \mathfrak{C}(UFU, U(FU)^2)$ .

Now when we apply the functor  $K$ , we have a simplicial functor  $KU \perp_\bullet$  in  $\text{Fun}(\mathbf{C}, \underline{\text{Set}})$  with an extra degeneracy.

At this point, we must evaluate  $KU \perp_\bullet$  at an object  $c \in \mathbf{C}$  to get a simplicial set for which we can describe the homotopy groups. We show  $\epsilon(c): KU \perp_\bullet(c) \rightarrow KU id_\bullet(c) = KU c_\bullet$  is a weak equivalence, i.e. all homotopy groups  $\pi_n(KU \perp_\bullet(c), *) \cong 0$  for  $n > 0$  and  $\pi_0(KU \perp_\bullet(c), *) \cong KU(c)$  for a basepoint  $* \in KU FU(c)$ . Suppose  $x, x' \in Z_n(KU \perp_\bullet(c), *)$  for  $n > 0$ . Recall that a horn  $h$  is given by  $(n+1)$   $n$ -simplices  $y_1, \dots, y_{i-1}, \hat{y}_i, y_{i+1}, \dots, y_{n+2}$  such that  $d_i y_j = d_{j-1} y_i$  if  $i < j$ . Note that we have a horn

$$h = (\hat{y}_0, y_1 = s_{-1}(x), y_2 = s_{-1}(x'), *, *, \dots, *) \in KU \perp_{n+1}(c)$$

if  $n > 0$  since for all  $i > 0$ ,

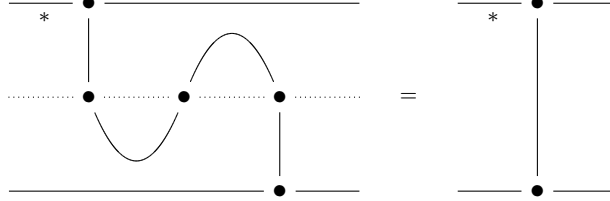
$$d_i s_{-1}(x) = s_{-1} d_{i-1}(x) = * = s_{-1} d_{i-1}(x') = d_i s_{-1}(x').$$

Since  $KU \perp_{\bullet}(c)$  is Kan, we can fill this horn, so there is a  $y_0 \in KU \perp_{n+1}(c)$  such that  $d_i y_0 = *$  for  $i > 1$ ,

$$\begin{aligned} d_0 y_0 &= d_0 y_1 = d_0 s_{-1}(x) = x \text{ and} \\ d_1 y_0 &= d_0 y_2 = d_0 s_{-1}(x') = x'. \end{aligned}$$

Hence  $y_0$  is a homotopy  $x \sim x'$ , and  $\pi_n(KU \perp_{\bullet}(c), *) = 0$  for all  $n > 0$ .

We show  $\pi = \pi_0(KU \perp_{\bullet}(c), *) \cong KU(c)$ . We have an extra degeneracy  $s_{-1}: KU(c) \rightarrow KU \perp_0(c)$  which satisfies  $\epsilon s_{-1} = \text{id}_{KU(c)}$ :

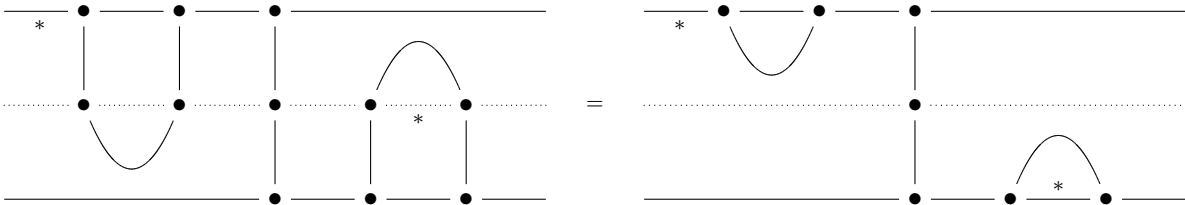


**Figure 10:**  $\epsilon s_{-1} = \text{id}_U \in \mathfrak{C}(U, U)$ .

Hence  $s_{-1}$  is injective and  $\epsilon$  is surjective. Now  $Z = Z_0(KU \perp_{\bullet}(c), *) = KU \perp_{\bullet}(c)$  and  $\epsilon(Z) = KU(c)$ . Note that if  $x \sim x'$  for  $x, x' \in Z$ , then there is a  $y \in KU \perp_1(c)$  such that  $d_0(y) = x$  and  $d_1(y) = x'$ . Then  $\epsilon(x) = \epsilon(x')$  as  $\epsilon d_1 = \epsilon d_0$  (this is the “capping off” trick discussed in (b)). Thus,  $\epsilon$  induces a map  $\tilde{\epsilon}: \pi \rightarrow KU(c)$  which must be surjective as  $\epsilon = \tilde{\epsilon}q$  where  $q: Z \rightarrow \pi$  is the canonical epimorphism.

$$\begin{array}{ccc} Z & \xrightarrow{\epsilon} & KU(C) \\ & \xleftarrow{s_{-1}} & \\ \downarrow q & \nearrow \tilde{\epsilon} & \\ \pi & & \end{array}$$

Moreover  $\tilde{\epsilon}q s_{-1} = \epsilon s_{-1} = \text{id}$ , so  $q s_{-1}$  is injective. Let  $x \in Z$  and  $y = s_{-1}(x)$ . Then  $d_0(y) = d_0 s_{-1}(x) = x$  and  $d_1(y) = d_1 s_{-1}(x) = s_{-1}\epsilon(x)$ .



**Figure 11:**  $d_1 s_{-1} = s_{-1} \epsilon \in \mathfrak{C}(UFU, UFU)$ .

Hence  $x \sim s_{-1}\epsilon(x)$ . Thus,  $q s_{-1}$  is bijective with inverse  $\tilde{\epsilon}$ , so  $\epsilon$  induces a bijection  $\pi \cong KU(c)$ , and we are finished.

(d) We will illustrate two examples. First, consider the example  $F = \text{Free}: \mathbf{Set} \rightarrow \mathbf{Ab}$  and  $U = \text{Forget}: \mathbf{Ab} \rightarrow \mathbf{Set}$ . Then  $\delta(c): \mathbb{Z}\langle A \rangle \rightarrow A$  is evaluation of a formal finite linear combination of elements of  $A \in \mathbf{Ab}$ . We have a simplicial group

$$\mathbb{Z}\langle A \rangle \begin{array}{c} \xrightarrow{s_0} \\ \xleftarrow{d_0, d_1} \end{array} \mathbb{Z}\langle \mathbb{Z}\langle A \rangle \rangle \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} \dots$$

The maps  $d_i$  are evaluation of a formal linear combination done at the  $i^{\text{th}}$  step. For example, if  $x \in \mathbb{Z}\langle \mathbb{Z}\langle A \rangle \rangle$ , then we have

$$x = \sum_{i=1}^{N_1} n_i \left( \sum_{j=1}^{N_2} m_j a_j \right),$$

where  $n_i, m_j \in \mathbb{Z}$  and  $a_j \in A$ , and we may not distribute the  $n_i$ 's into the sum over  $j$ . Then we have

$$d_0(x) = \sum_{i=1}^{N_1} \sum_{j=1}^{N_2} n_i m_j a_j \text{ and } d_1(x) = \sum_{i=1}^{N_1} n_i b_i \text{ where } b_j = \sum_{j=1}^{N_2} m_j a_j \in A.$$

Furthermore,  $\sigma(S): d \rightarrow U(\mathbb{Z}\langle S \rangle)$  is inclusion by  $s \mapsto 1 \cdot s$ , so if  $y \in \mathbb{Z}\langle A \rangle$ , then we have

$$y = \sum_{i=1}^{N_3} k_i a_i \text{ and } s_0(y) = \sum_{i=1}^{N_3} k_i (1 \cdot a_i)$$

where  $1 \cdot a \in \mathbb{Z}\langle A \rangle$ . It is clear that  $d_0 s_0 = d_1 s_0 = \text{id}$  in this case.

Consider  $F = \mathbb{Z}-: \mathbf{Group} \rightarrow \mathbf{Ring}$ , i.e. taking the group ring, and  $U = -^\times: \mathbf{Ring} \rightarrow \mathbf{Group}$ , i.e. taking the group of units. Then  $\delta(R): \mathbb{Z}(R^\times) \rightarrow R$  is once again evaluation and  $\sigma(G): G \rightarrow (\mathbb{Z}G)^\times$  is inclusion. We have a simplicial ring

$$\mathbb{Z}(R^\times) \begin{array}{c} \xrightarrow{s_0} \\ \xleftarrow{d_0, d_1} \end{array} \mathbb{Z}(\mathbb{Z}(R^\times))^\times \begin{array}{c} \xrightarrow{\quad} \\ \xleftarrow{\quad} \end{array} \dots$$

and maps  $d_i$  and  $s_j$  are defined similarly as before. □