Categorified trace for module tensor categories over braided tensor categories

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AMS Special Session on Fusion categories and topological phases of matter

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## Categorified traces

A trace on a k-algebra is a linear map  $\operatorname{tr}:A\to k$  such that

•  $\operatorname{tr}(ab) = \operatorname{tr}(ba)$  for all  $a, b \in A$ .

A categorified trace on a tensor category  $\mathcal M$  takes values in another category  $\mathcal C.$  It's a functor  $\mathrm{Tr}_{\mathcal C}:\mathcal M\to \mathcal C$  such that

- ▶ we have isomorphisms  $\tau_{x,y} : \operatorname{Tr}_{\mathcal{C}}(x \otimes y) \xrightarrow{\simeq} \operatorname{Tr}_{\mathcal{C}}(y \otimes x)$ natural in x and y called the *traciators*
- there is a hexagon compatibility axiom:



### An example of a categorified trace

Let  $\mathcal M$  be a pivotal tensor category. Define  $\mathrm{Tr}_{\mathsf{Vec}}:\mathcal M\to\mathsf{Vec}$  by

$$\operatorname{Tr}_{\mathsf{Vec}}(x) := \operatorname{Hom}_{\mathcal{M}}(1_{\mathcal{M}}, x).$$

The pivotal structure gives us our natural isomorphisms:

$$\operatorname{Tr}_{\mathsf{Vec}}(x \otimes y) = \operatorname{Hom}_{\mathcal{M}}(1_{\mathcal{M}}, x \otimes y)$$
$$\cong \operatorname{Hom}_{\mathcal{M}}(y^*, x)$$
$$\cong \operatorname{Hom}_{\mathcal{M}}(1_{\mathcal{M}}, y^{**} \otimes x)$$
$$\cong \operatorname{Hom}_{\mathcal{M}}(1_{\mathcal{M}}, y \otimes x)$$
$$= \operatorname{Tr}_{\mathsf{Vec}}(y \otimes x)$$

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#### Diagrammatic calculus for a categorified trace

Some algebraic equations are better represented on different topological spaces. Traces on algebras can be represented on  $S^1$ :

$$\operatorname{tr}(xy) = \begin{array}{c} x \\ y \end{array} = \begin{array}{c} x \\ y \end{array} = \begin{array}{c} y \\ x \end{array} = \operatorname{tr}(yx)$$

Similarly, we can denote the objects  $\operatorname{Tr}_{\mathcal{C}}(x \otimes y) \cong \operatorname{Tr}_{\mathcal{C}}(y \otimes x)$  by dots on circles, but we must remember the isomorphism  $\tau_{x,y}$ , which we denote by a cylinder with strings:

$$\tau_{x,y} = \underbrace{\prod_{x,y}}_{xy} : \operatorname{Tr}_{\mathcal{C}}(x \otimes y) \xrightarrow{\sim} \operatorname{Tr}_{\mathcal{C}}(y \otimes x)$$

### Categorified traces from central functors

When  $\mathcal{M}$  is pivotal, Bezrukavnikov-Finkelberg-Ostrik showed that if  $\Phi^{\mathcal{Z}} : \mathcal{C} \to \mathcal{Z}(\mathcal{M})$  is a functor such that  $\Phi := F \circ \Phi^{\mathcal{Z}}$  admits a right adjoint  $\operatorname{Tr}_{\mathcal{C}}$ , then  $\operatorname{Tr}_{\mathcal{C}}$  is a categorified trace.

The adjunction lets us construct the traciator  $\tau_{x,y}$ :

$$\begin{split} & \left[ \bigoplus \in \mathcal{C}(\operatorname{Tr}_{\mathcal{C}}(x \otimes y), \operatorname{Tr}_{\mathcal{C}}(x \otimes y)) \cong \mathcal{M}(\Phi(\operatorname{Tr}_{\mathcal{C}}(x \otimes y)), x \otimes y) \\ & \cong \mathcal{M}(\Phi(\operatorname{Tr}_{\mathcal{C}}(x \otimes y)) \otimes y^*, x) \\ & \cong \mathcal{M}(y^* \otimes \Phi(\operatorname{Tr}_{\mathcal{C}}(x \otimes y)), x) \\ & \cong \mathcal{M}(\Phi(\operatorname{Tr}_{\mathcal{C}}(x \otimes y)), y \otimes x) \\ & \cong \mathcal{C}(\operatorname{Tr}_{\mathcal{C}}(x \otimes y), \operatorname{Tr}_{\mathcal{C}}(y \otimes x)) \ni \end{split} \right]$$

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Suppose  $\mathcal{C}$  is a tensor category. A module category for  $\mathcal{C}$  is a category  $\mathcal{M}$  together with a bifunctor  $\otimes : \mathcal{C} \boxtimes \mathcal{M} \to \mathcal{M}$  satisfying certain axioms.

A braided tensor category C is a 3-category with one object and one 1-morphism, so it is one categorical dimension higher. Modules for C are also one categorical dimension higher.

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Module tensor categories for braided tensor categories

#### Definition

A module tensor category for the braided tensor category  $\mathcal{C}$  is a tensor category  $\mathcal{M}$  and a braided functor  $\Phi^{\mathcal{Z}} : \mathcal{C} \to \mathcal{Z}(\mathcal{M})$ .

• We define  $\Phi = F \circ \Phi^{\mathcal{Z}}$  where  $F : \mathcal{Z}(\mathcal{M}) \to \mathcal{M}$  is the forgetful functor.



# Example: $Mod_{\mathcal{C}}(a)$

Let C be a braided tensor category with  $a \in C$  a commutative algebra object. Let  $\mathcal{M} = Mod_{\mathcal{C}}(a)$ , the category of left *a*-module objects in C.

#### Definition

We define the free module functor  $\Phi^{\mathcal{Z}}:\mathcal{C}\to\mathcal{M}$  by  $\Phi^{\mathcal{Z}}(c)=(\Phi(c),e_{\Phi(c)})$  where

• 
$$\Phi(c) = a \otimes c$$

• 
$$e_{\Phi(c),x}: \Phi(c) \otimes_a x \cong c \otimes x \xrightarrow{\beta_{c,x}} x \otimes c \cong x \otimes_a \Phi(c).$$

$$\bullet \ \Phi(f: c \to d) = \mathrm{id}_a \otimes f: a \otimes c \to a \otimes d$$

#### Example

Let  $C = SU(2)_{10}$  with objects  $1, \ldots, 11$ . Let  $a = 1 \oplus 7$ . Then  $\mathcal{M} = Mod_{\mathcal{C}}(a)$  is the  $E_6$  module tensor category.

### Categorified trace for module tensor categories

We now assume C is a pivotal braided tensor category,  $\mathcal{M}$  is a pivotal module tensor category, and  $\Phi^{\mathcal{Z}} : C \to \mathcal{Z}(\mathcal{M})$  is a braided pivotal functor.

Under mild assumptions,  $\Phi:\mathcal{C}\to\mathcal{M}$  has a right adjoint  $\mathrm{Tr}_{\mathcal{C}}$ , which is a categorified trace.



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## Adjoints of tensor functors are lax monoidal

It is well known that the right adjoint of a tensor functor is lax monoidal, so we get canonical natural multiplication maps

$$\mu_{x,y} = \prod_{\substack{x,y \\ x \ y}} : \operatorname{Tr}_{\mathcal{C}}(x) \otimes \operatorname{Tr}_{\mathcal{C}}(y) \to \operatorname{Tr}_{\mathcal{C}}(x \otimes y)$$

#### Theorem

The traciators and multiplication maps are compatible, giving us a graphical caluclus of srings on tubes which can branch and braid.

# Graphical calculus



 $\begin{array}{l} {\rm Tr}_{\mathcal{C}}: \mathcal{M} \rightarrow \mathcal{C} \\ {\rm is \quad a \quad categorified} \\ {\rm trace \ [BF009, \ Ost14]} \end{array}$ 





braided pivotal functor

 $\Phi^{\mathcal{Z}}$ 

С

 $\mathcal{Z}(\mathcal{M})$ 

 $\operatorname{Tr}_{\mathcal{C}}$ 

 $\mathcal{M}$ 



 $\begin{array}{l} {\rm Tr}_{\mathcal{C}}: \mathcal{M} \rightarrow \mathcal{C} \\ {\rm is \ a \ lax \ monoidal} \\ {\rm functor \ [Kel74].} \end{array}$ 



Application: constructing algebras in  $\ensuremath{\mathcal{C}}$ 

Since tensor functors are lax monoidal, given an algebra  $a \in \mathcal{M}$ ,  $\operatorname{Tr}_{\mathcal{C}}(a)$  is an algebra in  $\mathcal{C}$ .

#### Theorem

If  $a, b \in \mathcal{M}$  are algebras, then  $\operatorname{Tr}_{\mathcal{C}}(a \otimes_{\mathcal{M}} b)$  is an algebra in  $\mathcal{C}$ .



## Application: anchored planar algebras

We can use the internal trace to construct planar algebras *internal* to a braided tensor category.

#### Definition

A planar algebra in Vec is a sequence of objects  $\mathcal{P}_n \in$  Vec and an action of planar tangles.

An anchored planar algebra in C is a sequence of objects  $\mathcal{P}_n \in C$ and an action of anchored planar tangles.



### Thank you for listening!

Slides available at: http://www.math.ucla.edu/~dpenneys/PenneysAMS2016.pdf

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Article available at: http://arxiv.org/abs/1509.02937